# jClarity

### **Hotspot Garbage Collection - The Useful Parts**

### Martijn Verburg (@karianna)

Session Code: 1500

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### Who am I?

### aka "The Diabolical Developer"

- I cause trouble in the Java/JVM and F/OSS worlds

- Especially Agile/Scrum/SC BS

### CTO of jClarity

- Java Performance Tooling start-up
- "Measure don't guess"

### Co-lead London Java Community (LJC)

- Run global programmes to work on OpenJDK & Java EE
- Adopt-a-JSR and Adopt OpenJDK
- Community night tomorrow night!

# What I'm going to cover

- Part I Diving into the Dark (~30 min)
  - GC Theory
  - Hotspot memory organisation and collectors
- Break! (2 min)
  - Our brains hurt

#### • Part II - Shining a light into the Darkness (8 min)

- Reading GC Logs
- Tooling and Basic Data

#### • Part III - Real World Scenarios (8 min)

- Likely Memory Leaks
- Premature Promotion
- Healthy App
- High Pausing

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### What I'm not covering

#### G1 Collector

- It's supported in production now
- Not a lot of good independent empirical research on this

#### • JRockit, Azul Zing, IBM J9 etc

- Sorry, these warrant their own talks
- Go see Azul on level 3 though, what they do is... cool.

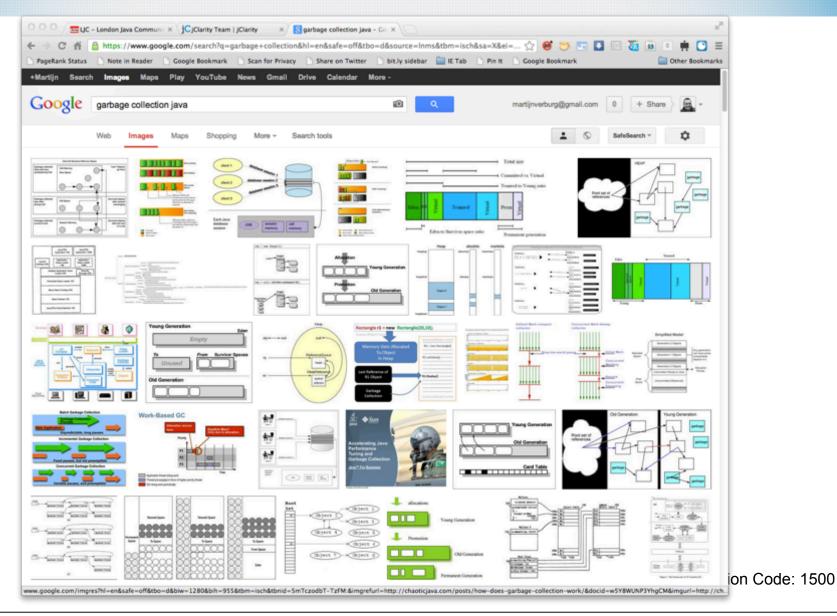
#### PhD level technical explanations

- I want you to have a working understanding
  - Reality: I'm not that smart
- Going for that PhD? See me after

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### **Search for Garbage Collection..**



Thursday, 7 March 13

# 

# **Part I - Diving into the Dark**

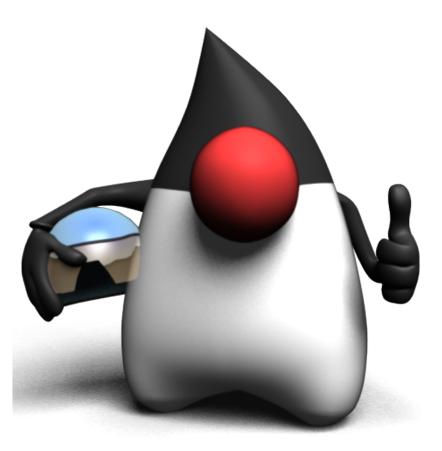
- What is Garbage Collection (GC)?
- Hotspot Memory Organisation
- Collector Types
- Young Collectors
- Old Collectors
- Full GC

# What is Garbage Collection (GC)?

- The freeing of memory that is no longer "live"
  - Otherwise known as "collecting dead objects"
    - Which is a misnomer
- GC is typically executed by a managed runtime
- Javascript, Python, Ruby, .NET CLR all have GC

### And so does Java!

One of the main 'selling' points in its early life



### Why should I care?

- Hotspot just sorts this out doesn't it?
- Just set -xms and -xmx to be == right?
  - Stab myself in the eye with a fork
- A poorly tuned GC can lead to:
  - High pause times / high % of time spent pausing
  - OutOfMemoryError
- It's usually worth tuning the GC!
  - "Cheap" performance gain
  - Especially in the short to medium term

### **Hotspot Java Virtual Machine**

#### Hotspot is a C/C++/Assembly app

- Native code for different platforms
- Roughly made up of Stack and Heap spaces

#### The Java Heap

- A Contiguous block of memory
- Entire space is reserved
- Only some space is allocated
- Broken up into different memory pools

#### Object Creation / Removal

- Objects are created by application (mutator) threads
- Objects are removed by Garbage Collection

### **Memory Pools**

#### Young Generation Pools

- Eden
- Survivor 0
- Survivor 1

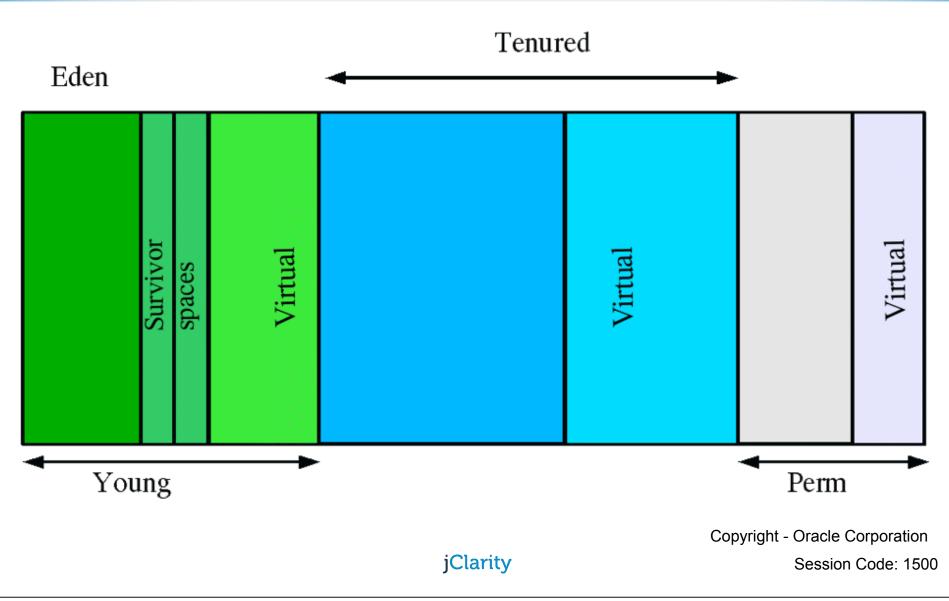
#### Old Generation Pool (aka Tenured)

- Typically much larger than young gen pools combined

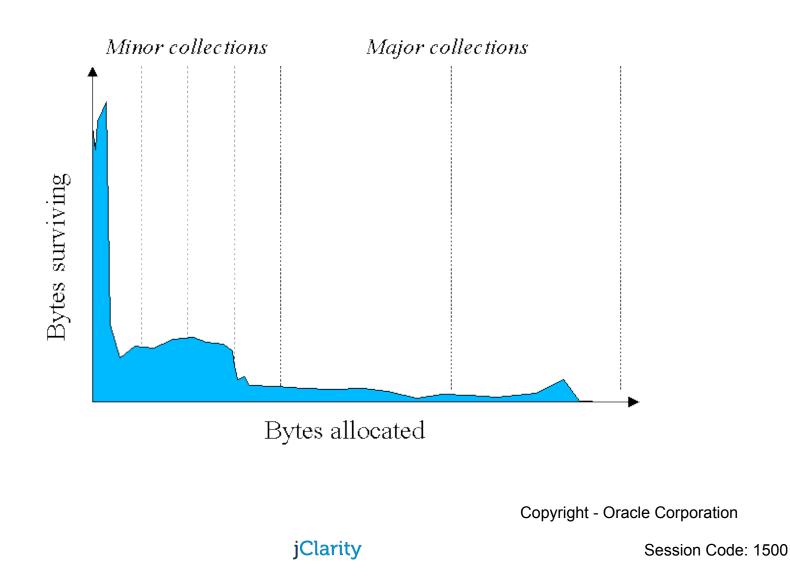
#### PermGen Pool

- Held separately to the rest of the Heap
- Was intended to hold objects that last a JVM lifetime
  - Reloading and recycling of classes occurs here.
- Going away in Java 8

### **Java Heap Layout**



### **Weak Generational Hypothesis**



### Only the good die young...



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# Сору

#### aka "stop-and-copy"

- Some literature talks about "Cheney's algorithm"

#### Used in many managed runtimes

- Including Hotspot
- GC thread(s) trace from root(s) to find live objects

#### Typically involves copying live objects

- From one space to another space in memory
- The result typically looks like a move as opposed to a copy

### Mark and Sweep

#### Used by many modern collectors

- Including Hotspot, usually for old generational collection

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#### • Typically 2 mandatory and 1 optional step(s)

- 1. Find live objects (*mark*)
- 2.'Delete' dead objects (sweep)
- 3. Tidy up optional (*compact*)

### **Mark and Sweep collectors in Hotspot**

#### Several Hotspot collectors use Mark and Sweep

- Concurrent Mark and Sweep (CMS)
- Incremental Concurrent Mark and Sweep (iCMS)
- MarkSweepCompact (aka Serial)
- PS MarkSweep (aka ParallelOld)

#### So it's worth learning the theory

### Java objects

#### Java objects have Ordinary Object Pointers (OOPs)

- That point to an object...
- Which points to the header

#### The header contains a mark bit for GC

- Plus other metadata (hashcodes, locking state etc)

#### When you call a constructor

- Space for the object is allocated

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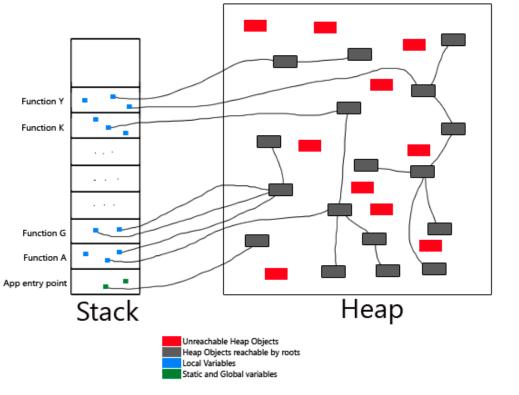
### **Step 1 - Clear the Mark**

- The header contains the boolean mark field
  - If true --> the object is live
- Step 1 set all the mark fields to false
  - We need to start afresh

### **Step 2 - Mark live objects**

GC Roots

- A pointer to data in the heap that you need to keep



Copyright - Michael Triana

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## **Step 2 - Mark live objects**

#### • GC Roots are made up of:

- Live threads
- Objects used for synchronisation
- JNI handles
- The system class loaders
- Possibly other things depending on your JVM
- Plus one more special case...

## **Step 2 - Mark live objects**

- Special case Old Gen refs into Young Gen
  - Treated as roots during a young collection
- Special card table to track these
  - Each card references an area of 512 bytes in old gen
  - If it references young gen it will have been marked as dirty
  - Dirty areas are scanned as part of the young collection
- Conclusion there's a lot to trace!

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Step 3 - Sweep

#### • Sweep

- Mark space that dead objects occupy as deleted

#### • Compact

- Not part of the normal operation of some collectors
- Always attempted before OOME's can be thrown
- 'Defrags' the remaining space
  - Not quite a full defrag
- I'll cover some Java specific collectors shortly

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### **Heap of Fish Demo**

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# **Young Generation Pools**

#### • Eden

- Where new objects should get created
- Objects are added at the end of currently allocated block
- Uses Thread Local Allocation Buffers (TLABs)
  - · Points at end of allocated block of objects

#### Survivor 0 and Survivor 1

- Known as Hemispheric GC
- Only one is active at a time
- The other one is empty, we call it the *target* space

### **Young Generation Collectors**

#### When Eden gets "full"

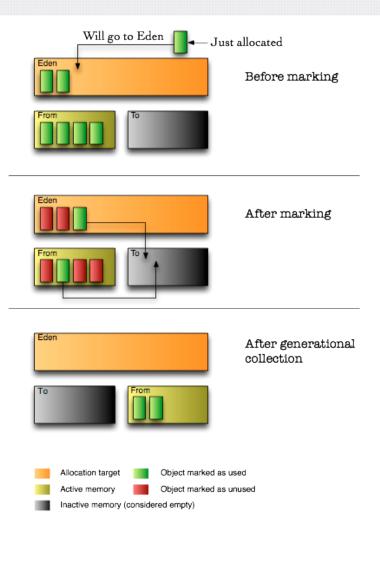
- "Full" is technically passing a threshold
- A collector will run
- Live objects get copied to the *target* Survivor space
  - From Eden and active Survivor space

#### Some Live objects are promoted to Old Gen

- If they've survived > tenuringThreshold collections
- Or if they can't fit in the *target* space

#### When the collector is finished

- A simple pointer swizzle activates the *target* Survivor space
- Dead objects effectively disappear (no longer referenced)



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### **Young Generation Collectors**

#### Most use parallel threads

- i.e. A multi-core machine can make your GC faster

#### • I'll cover the PS Scavenge and ParNew collectors

- They're almost identical
- PS Scavenge works with PS MarkSweep old gen
- ParNew works with ConcurrentMarkSweep (CMS) old gen

#### Other young collectors:

Copy (aka Serial)

– G1

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### **PS Scavenge / ParNew**

- aka "Throughput collectors"
- Number of threads is set as a ratio to # of cores
- They're Stop-The-World (STW) collectors
  - They're monolithic (as opposed to incremental)
- Each thread gets a set of GC roots
  - They do work stealing
- It performs an copy (aka evacuate)
  - Surviving objects move to the newly active survivor pool

### **Age and Premature Promotion**

- Objects have an age
- Every time they survive a collection..
  - age++
- At age > tenuringThreshold
  - Objects get moved (promoted) to old/tenured space
  - Default tenuringThreshold is 4
- Premature Promotion occurs when
  - High memory pressure (high life over death ratio)
    - Eden is too small to deal with rate of new objects
  - Objects are too big to fit in Eden
  - Objects are too big to be promoted to Survivor spaces



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### **Old Generation Collectors**

Most are variations on Mark and Sweep

#### Most use parallel threads

- e.g. A multi-core machine can make your GC faster

#### I'll cover PS MarkSweep & CMS

- CMS is often paired with the ParNew young collector

#### Other old collectors:

- MarkSweepCompact (aka Serial)
- Incremental CMS (iCMS)

- G1

### **PS MarkSweep**

#### • aka "ParallelOld"

- Often paired with PS Scavenge for young gen

#### Parallel GC threads get sections to look after

- Usual Mark and Sweep occur

#### Special Compact phase takes place

- low occupancy sections get merged
- e.g. A compact / defrag operation

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### **CMS Old Gen Collector**

#### Only runs when Tenured is about to get full

- Tunable as to what 'about to get full' means

#### Attempts to share CPU with application

- About a 50/50 ratio as a default
- Application can keep working whilst GC is taking place

#### It's a partial Stop-The-World (STW) collector

- It has 6 phases
  - 2 STW
  - 4 Concurrent

#### • It does not compact unless it fails..

#### **CMS** Phases

- Phase 1 Initial Mark (STW)
  - Marks objects adjacent to GC roots
- Phase 2 Mark (Concurrent)
  - Completes depth first marking

#### • Phase 3 - Pre Clean (Concurrent)

- Retraces the updated objects, finds dirty cards

#### • Phase 4 - Re Mark / Re Scan (STW)

- Hopefully a smaller graph traversal over dirty paths

#### Phase 5/6 - Concurrent Sweep and Reset

- Sweep out dead objects and reset any data structures

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# **Concurrent Mode Failure (CMF)**

- Occurs when CMS can't complete 'in time'
  - 'In time' meaning that tenured has filled up
- GC subsystem reverts to a Full GC at this point
  - Basically ouch

#### **Promotion Failure**

#### Occurs when objects can't be promoted into Tenured

- Often due to the Swiss Cheese nature of Old Gen

- Because CMS does not compact
- This will almost always happen.... eventually

#### • Triggers a Full GC

- Which compacts old space
- No more Swiss Cheese! For a short while...

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# **Full GC**

#### Can be triggered by a number of causes

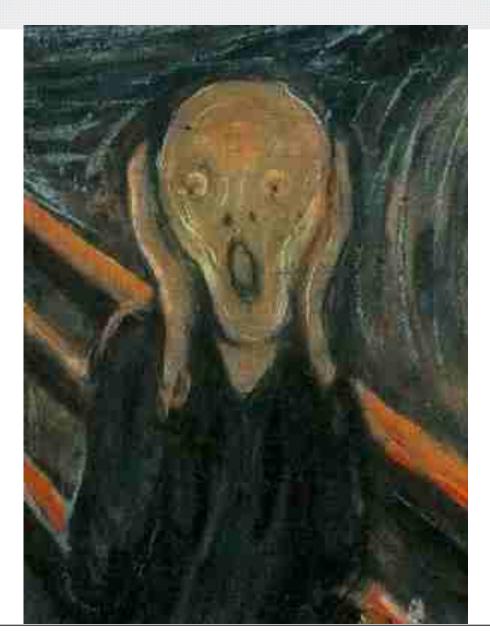
- A CMF from the CMS Collector
- Promotion Failure
- When tenured gets above a threshold
- System.gc()
- Remote System.gc() via RMI

#### Runs a full STW collection

- Over Young and Old generational spaces
- Compacts as well

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#### **Special Case: OOME**



### **Special Case: OOME**

- 98%+ time is spent in GC
- < 2% of Heap is freed in a collection</p>
- Allocating an object larger than heap
- Sometimes when the JVM can't spawn a new Thread

# Part II - Shining a light into the dark

- Collector Flags aboy
- Reading CMS Log records
- Tooling and basic data

# 'Mandatory' Flags

#### • -verbose:gc

- Get me some GC output

#### • -Xloggc:<pathtofile>

- Path to the log output, make sure you've got disk space

#### • -XX:+PrintGCDetails

- Minimum information for tools to help
- Replace -verbose:gc with this

#### • -XX:+PrintTenuringDistribution

Premature promotion information

#### -XX:+PrintGCApplicationStoppedTime

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# **Basic Heap Sizing Flags**

• -Xms<size>

– Set the minimum size reserved for the heap

#### • -Xmx<size>

- Set the maximum size reserved for the heap

#### • -XX:MaxPermSize=<size>

- Set the maximum size of your perm gen
- Good for Spring apps and App servers

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#### **Other Flags**

- -XX:NewRatio=N
- -XX:NewSize=N
- -XX:MaxNewSize=N
- -XX:MaxHeapFreeRatio
- -XX:MinHeapFreeRatio
- -XX:SurvivorRatio=N
- -XX:MaxTenuringThreshold=N
- • • •

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#### More Flags than your Deity



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# Why Log Files?

Log file can be post processed

#### Log files contain more information

- Than runtime MXBeans

#### Runtime MXBeans impact the running application

- Causing it's own GC problems!

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## **Raw GC Log File**

```
Z \odot
                                                                                       par.cms.wd.wt.log - Kate
File Edit View
                 Bookmarks Sessions Tools Settings
                                                    Help
O New
                                         🔚 Save 📝 Save As
                                                              🔀 Close 👘 Undo ( 🕮 Redo
          Open |
                   Back
                            Forwards
    Desired survivor size 1343488 bytes, new threshold 2 (max 4)
Documents
                  1304520 bytes,
           1:
                                     1304520 total

    age

            2:
                     79280 bytes,
                                     1383800 total

    age

    age

           3:
                    55176 bytes,
                                     1438976 total
                   370720 bytes,

    age

           4:
                                     1809696 total
: 16828K->1769K(18624K), 0.0030239 secs] 74972K->60174K(83392K), 0.0031259 secs] [Times: user=0.00 sys=0.00, real=0.00 secs]
    23.570: [GC [1 CMS-initial-mark: 58405K(64768K)] 62606K(83392K), 0.0008419 secs] [Times: user=0.00 sys=0.00, real=0.00 secs]
    23.571: [CMS-concurrent-mark-start]
    23.586: [GC 23.586: [ParNew
    Desired survivor size 1343488 bytes, new threshold 2 (max 4)
                   634264 bytes,
                                      634264 total
    - age 1:
    - age 2:
                  1184776 bytes,
                                     1819040 total
    : 17740K->1779K(18624K), 0.0034827 secs] 76145K->60479K(83392K), 0.0035902 secs] [Times: user=0.01 sys=0.00, real=0.00 secs]
    23.605: [GC 23.605: [ParNew
    Desired survivor size 1343488 bytes, new threshold 4 (max 4)
    - age 1:
                   432832 bytes,
                                      432832 total
                   591944 bytes,
                                     1024776 total
          2:

    age

     : 17771K->1003K(18624K), 0.0020149 secs] 76471K->60386K(83392K), 0.0021186 secs] [Times: user=0.00 sys=0.00, real=0.00 secs]
    23.622: [GC 23.622: [ParNew
    Desired survivor size 1343488 bytes, new threshold 3 (max 4)
                   427360 bytes,
                                      427360 total

    age

           1:

    age

            2:
                   379488 bytes,
                                      806848 total
                   591944 bytes,
                                     1398792 total

    age

           3:
    : 16995K->1368K(18624K), 0.0021869 secs] 76378K->60751K(8392K), 0.0023114 secs] [Times: user=0.00 sys=0.00, real=0.00 secs]
    23.639: [GC 23.639: [ParNew
    Desired survivor size 1343488 bytes, new threshold 4 (max 4)
                   313136 bytes,
                                      313136 total
          1:

    age

    - age 2:
                   370240 bytes,
                                      683376 total
           3:
                   379488 bytes,
                                     1062864 total

    age

    : 17318K->1041K(18624K), 0.0020645 secs] 76701K->61002K(83392K), 0.0021475 secs] [Times: user=0.01 sys=0.00, real=0.00 secs]
    23.652: [CMS-concurrent-mark: 0.068/0.081 secs] [Times: user=0.12 sys=0.01, real=0.09 secs]
    23.652: [CMS-concurrent-preclean-start]
    23.652: [CMS-concurrent-preclean: 0.000/0.000 secs] [Times: user=0.00 sys=0.00, real=0.00 secs]
    23.652: [CMS-concurrent-abortable-preclean-start]
    23.655: [GC 23.655: [ParNew
    Desired survivor size 1343488 bytes, new threshold 4 (max 4)

    age

           1:
                   333304 bytes,
                                      333304 total

    age

            2:
                    252120 bytes,
                                      585424 total

    age

           3:
                   370240 bytes,
                                      955664 total
                    279499 butor
                                      1225152 tota
```

Line: 5,553 Col: 58 INS\_LINE\_UTF-8\_par.cms.wd.wt.log





#### **General Format**

# from->to(total size) i.e: 16963K->884K(18624K)

#### ocupancy(size) i.e: 62606K(83392K)

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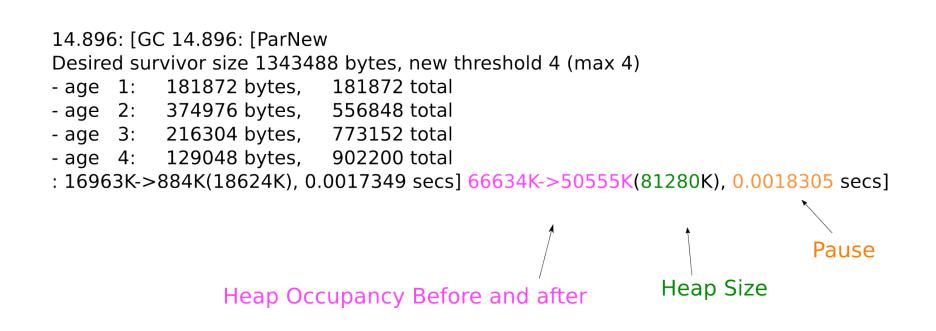
# Young Gen Collection Part I

14.896: [GC 14.896: [ParNew Desired survivor size 1343488 bytes, new threshold 4 (max 4) - age 1: 181872 bytes, 181872 total - age 2: 374976 bytes, 556848 total - age 3: 216304 bytes, 773152 total - age 4: 129048 bytes, 902200 total : 16963K->884K(18624K), 0.0017349 secs] 66634K->50555K(81280K), 0.0018305 secs]

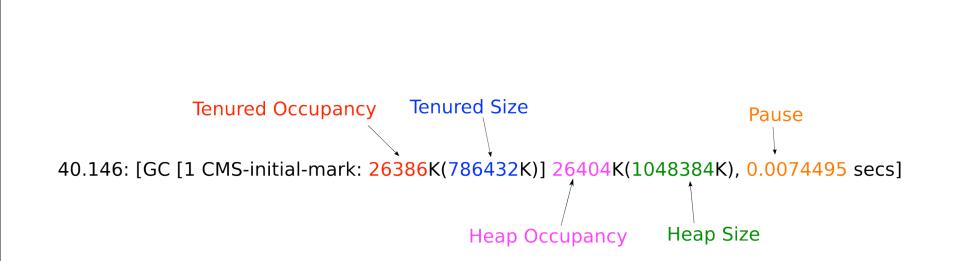
Young Occupancy before and after

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## Young Gen Collection Part II



#### **CMS Initial Mark**



#### **AII CMS**

12.986: [GC [1 CMS-initial-mark: 33532K(62656K)] 49652K(81280K), 0.0014191 secs]

12.987: [CMS-concurrent-mark-start] 13.071: [CMS-concurrent-mark: 0.068/0.084 secs]

13.071: [CMS-concurrent-preclean-start] 13.075: [CMS-concurrent-preclean: 0.001/0.004 secs]

13.077: [GC[YG occupancy: 3081 K (18624 K)]13.077: [Rescan (parallel) , 0.0009121 secs]13.078: [weak refs processing, 0.0000365 secs] [1 CMS-remark: 35949K(62656K)] 39030K(81280K), 0.0010300 secs]

13.078: [CMS-concurrent-sweep-start] 13.097: [CMS-concurrent-sweep: 0.016/0.019 secs]

13.264: [CMS-concurrent-reset-start] 13.266: [CMS-concurrent-reset: 0.001/0.001 secs] Tenured Occupancy Tenured Size Young occupancy Young Size Heap Occupancy Heap Size Pause Time

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# Tooling

- HPJMeter (Google it)
  - Solid, but no longer supported / enhanced
- GCViewer (<u>http://www.tagtraum.com/gcviewer.html</u>)
  - Has rudimentary G1 support
- GarbageCat (<u>http://code.google.com/a/eclipselabs.org/p/garbagecat/</u>)
   Best name
- IBM GCMV (<u>http://www.ibm.com/developerworks/java/jdk/tools/gcmv/</u>)
   J9 support
- jClarity Censum (<u>http://www.jclartity.com/products/censum</u>)
  - The prettiest and most useful, but we're biased!

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#### **HPJMeter - Summary**

<u>F</u> ile ⊻iew <u>H</u> e	lp					
<b>1</b>						
Summary	Heap Usage After GC Du	ration Cumulative Allocat	ion Creation Rate Use	r-Defined   Multiple User-De	efined	
Heap Capacity						
	Ed	en S	urvivor	Old	Perm	Total
Initial Capacity	N/A	N/	Ą	N/A	N/A.	57.375 (MB)
Final Capacity	N/A	N/	A	783.625 (MB)	20.75 (MB)	933.375 (MB)
Peak Capacity	N/A	N/	A	783.625 (MB)	20.75 (MB)	933.375 (MB)
Peak Usage of Capaci	V N/A	N/	Ą	100%	16.049%	100%
GC Activity Summary						
	Last occurre	nce (s) Count	Average in	terval (s) Ave	erage duration (s)	Average rate of collection
CMS	84.54 (s)	27	3.181 (s)	0.81	. (s)	0 (B/s)
Parallel Scavenge	60.839 (s)	226	0.266 (s)	0.075	5 (s)	163.802 (MB/s)
Other full GC	94.216 (s)	2	9.797 (s)	7.183	3 (s)	13.349 (MB/s)
Overall Statistics						
Name		Value	Name	2	v	/alue
Duration of the measu	rement	96.08 (s)	96.08 (s) Time :		53	).165 (s)
Total bytes allocated		3.798 (GB)	3.798 (GB) Perce		55	5.335%
Number of GC events		255	255 Time		14	1.366 (s)
Average bytes allocate	ed per GC	15.251 (MB)	Percer	centage of time in Full GC		4.952%
Avg. ideal allocation rate		90.625 (MB/s)	Avg. al	allocation rate		0.478 (MB/s)
Residual bytes		850.562 (MB)				
1						



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## **HPJMeter - Heap Usage After GC**



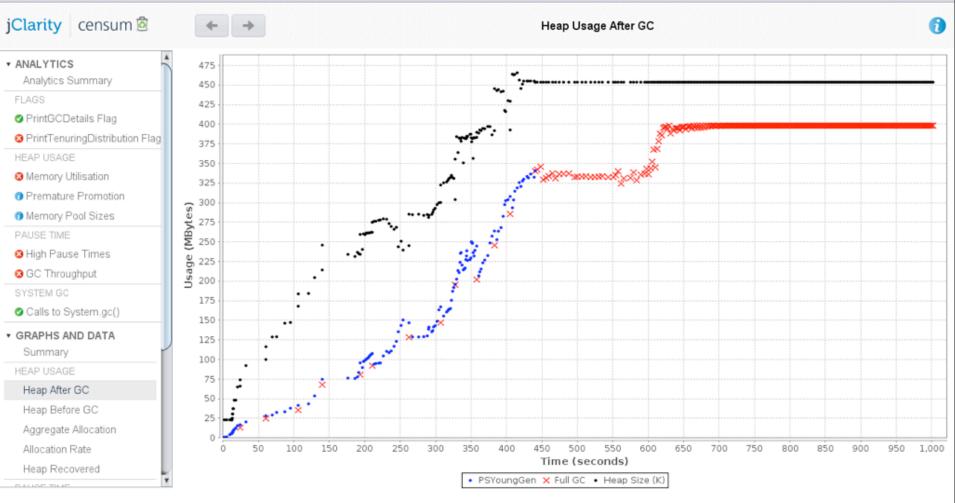
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## **Part III - Scenarios**

- Possible Memory Leak(s)
- Premature Promotion
- Healthy Application
- High percentage of time spent pausing

# **A Memory Leak**

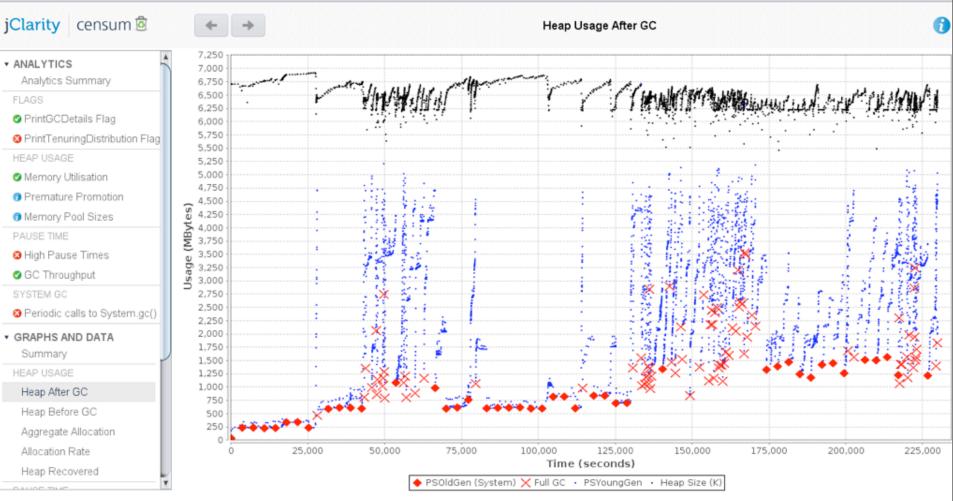
File Help



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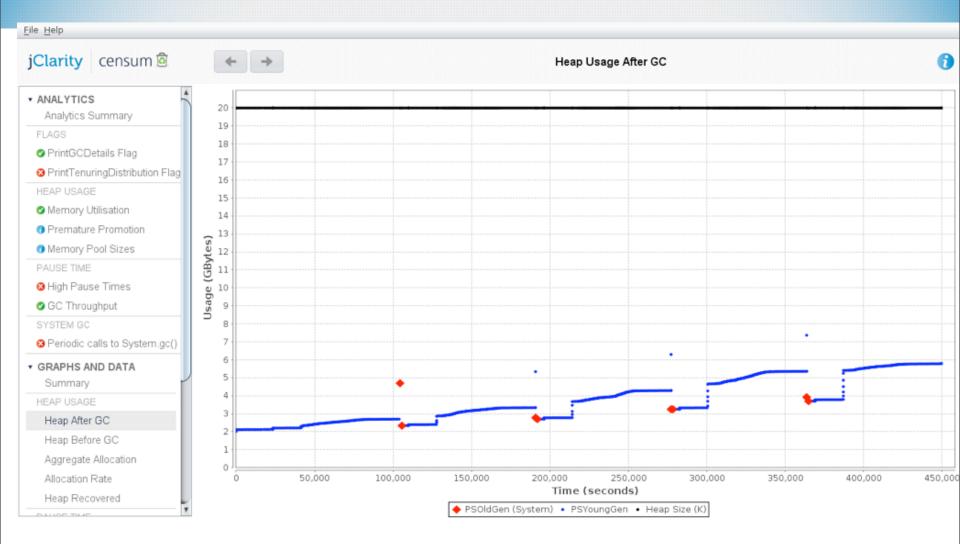
# A Possible Memory Leak - I





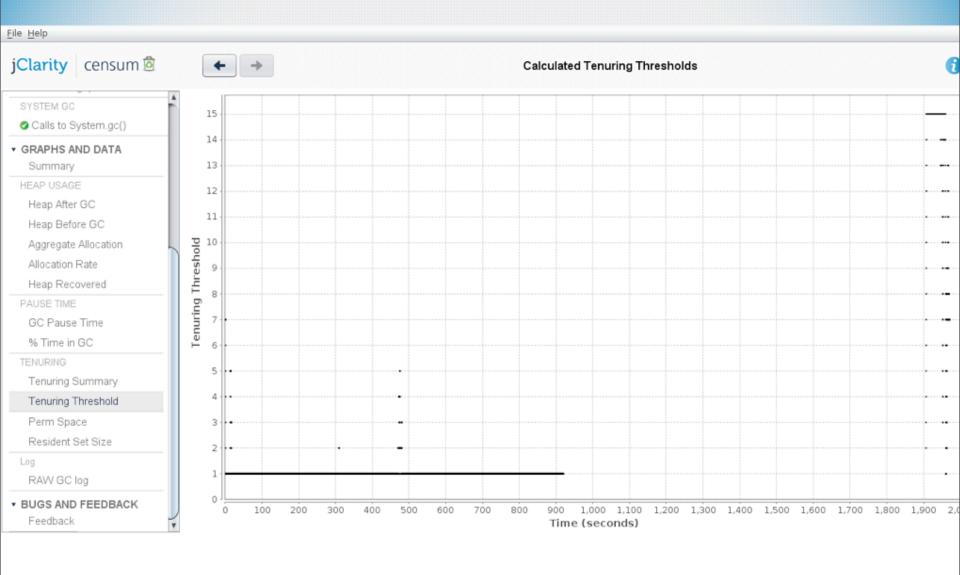
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## A Possible Memory Leak - II



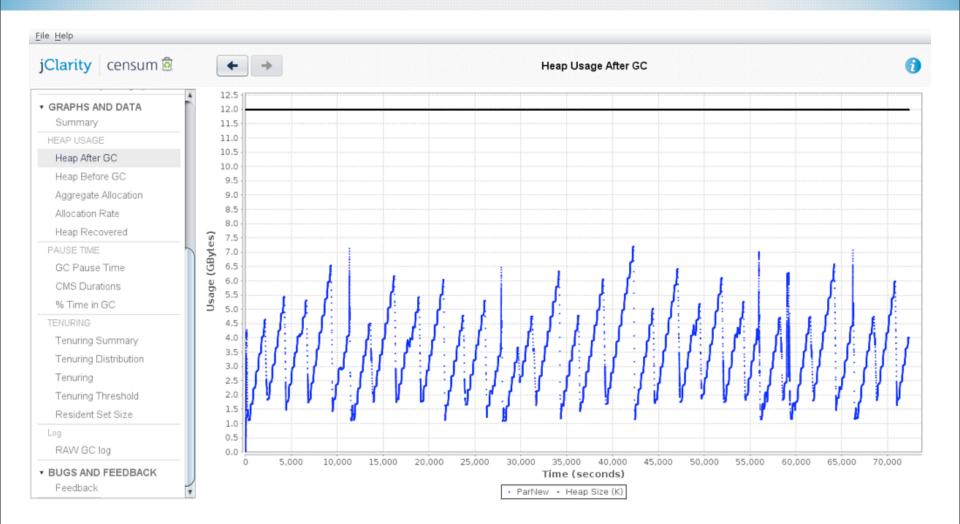
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## **Premature Promotion**



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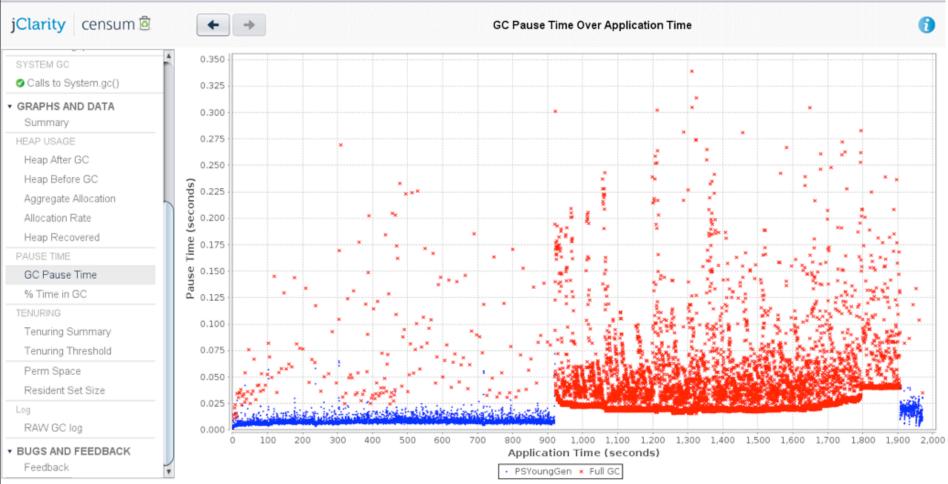
# **Healthy Application**



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## **High Percentage of time Paused**

#### <u>File</u> <u>H</u>elp



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- You need to understand some basic GC theory
- You want most objects to die young, in young gen

#### Turn on GC logging!

- Reading raw log files is hard
- Use tooling!

#### Use tools to help you tweak

- "Measure, don't guess"

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# Join our performance community

# http://www.jclarity.com

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