CRDTS Data Types for EC Systems



Eventual Consistency

Eventual consistency is a consistency model used in distributed computing that informally guarantees that, if no new updates are made to a given data item, eventually all accesses to that item will return the last updated value.

--Wikipedia









\$\$\$Big Iron (still fails)





Commodity Servers CDNs, App servers Expertise











Amazon found every 100ms of latency cost them 1% in sales.



Low Latency



Low Latency





Low Latency

Google found an extra 0.5 seconds in search page generation time dropped traffic by 20%.









Causal RYOW Session Monotonic Read/Write





Pick Your Own

<u>http://basho.com/understanding-riaks-</u> <u>configurable-behaviors-part-1/</u>

http://www.bailis.org/blog/when-is-acid-acidrarely/







Google F1

"We have a lot of experience with eventual consistency systems at Google."

"We find developers spend a significant fraction of their time building extremely complex and error-prone mechanisms to cope with eventual consistency"





Riak Overview Erlang implementation of Dynamo











Keys are namespaced into Buckets



key	value
key	value

Riak Overview Consistent Hashing













Riak Overview Dynamic Membership

Replication factor







Replica



High Availability

Any non-failing node can respond to any request.

--Gilbert & Lynch



Riak Overview Two Writes: {Writer, Value, Time}







Riak Ov Last Writer Wins Allow Mult



Riak Overview Last Writer Wins [{b, v1, t2}]

[{b, v1, t2}]

[{b, v1, t2}]
Allow Mult

Riak Overview

[{a, v1, a1}, {b, v2, b1}]

[{a, v1, a1}, {b, v2, b1}]

[{a, v1, a1}, {b, v2, b1}]

User specified



Semantic Resolution

Fill in Form P17QR-35 File for input



DB IS DOWN

Fill in Form P17QR-35 File for input

OOPS



DB IS DOWN

Dynamo The Shopping Cart





HAIRDRYER

















PENCIL CASE











Set Union of Values Simples, right?









Deterministic **Idempotent**





Deterministic Idempotent Associative





Deterministic Idempotent Associative Commutative



Absence

How can you tell if X is missing from A and not B because A hasn't seen the addition, or if A has removed X?











Convergent Replicated Data Types



Commutative Replicated Data Types



CRDTS Conflict-free Replicated Data Types









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A comprehensive study of **Convergent and Commutative Replicated Data Types**

Marc Shapiro, INRIA & LIP6, Paris, France Nuno Preguiça, CITI, Universidade Nova de Lisboa, Portugal Carlos Baquero, Universidade do Minho, Portugal

Marek Zawirski, INRIA & UPMC, Paris, France



INSTITUT NATIONAL DE RECHERCHE EN INFORMATIQUE ET EN AUTOMATIQUE

An Optimized Conflict-free Replicated Set

Annette Bieniusa, INRIA & UPMC, Paris, France Marek Zawirski, INRIA & UPMC, Paris, France Nuno Preguiça, CITI, Universidade Nova de Lisboa, Portugal Marc Shapiro, INRIA & LIP6, Paris, France Carlos Baquero, HASLab, INESC TEC & Universidade do Minho, Portugal Valter Balegas, CITI, Universidade Nova de Lisboa, Portugal

Sérgio Duarte CITI, Universidade Nova de Lisboa, Portugal



INRIA

Dotted Version Vectors: Logical Clocks for Optimistic Replication

Nuno Preguiça CITI/DI FCT, Universidade Nova de Lisboa Monte da Caparica, Portugal nmp@di.fct.unl.pt

Abstract

The mentioned systems follow a design where the data store is always writable. A consequence is that In cloud computing environments, a large number replicas of the same data item are allowed to diverge, of users access data stored in highly available storage and this divergence should later be repaired. Accurate systems. To provide good performance to geographitracking of concurrent data updates can be achieved cally disperse users and allow operation even in the by a careful use of well established causality tracking presence of failures or network partitions, these sysmechanisms [5], [6], [7], [8]. In particular, for data tems often rely on optimistic replication solutions that storage systems, version vectors [6] enables the system guarantee only eventual consistency. In this scenario, to compare any pair of replica versions and detect if it is important to be able to accurately and efficiently



Carlos Baquero, Paulo Sérgio Almeida, Victor Fonte, Ricardo Gonçalves CCTC/DI Universidade do Minho Braga, Portugal {cbm,psa,vff}@di.uminho.pt, rtg@lsd.di.uminho.pt

Bounded Join Semilattices



Bounded Join Semilattices Partially ordered set; least upper bound; ACI.



Bounded Join Semilattices Associativity: (X Y) Z = X (Y Z)



Bounded Join Semilattices Commutativity: X Y = Y X



Bounded Join Semilattices Idempotence: X X = X



Bounded Join Semilattices Objects grow over time; merge computes LUB



Bounded Join Semilattices



Examples



Set; merge function: union.


Increasing natural; merge function: max.



Booleans; merge function: or.





Deterministic Idempotent Associative Commutative



Pick your semantic Add wins **Remove wins** Keep both







Version Vectors Entry Per Actor (Charron-Bost)

Actors?

RIAK 1.4 Counters: non-idempotent; 0 (Actors)

Riak 200 Sets: Add, Remove, Membership; Idempotent

Sets: Add wins 0(Actors + Elements)



Riak 2.0 Maps: Recursive; Associative Array; Nestable

Riak 2

Maps: Update wins; O(Actors + Elements)



Riak 2.0 Maps: LWW-Register, Booleans, Sets and Maps

Riak 2.0 LWW-Register: last writer wins

Boolean: Enabled, Disabled; O(Actors)



git clone git@github.com:basho/riak_dt.git







Evolution of a Set

Causality Version Vectors

[{a, 1}, {b, 3}, {c, 2}]



Causality Version Vectors

[{a, 2}, {b, 3}, {c, 2}]

> [{a, 1}, {b, 3}, {c, 2}]

Causality Version Vectors

[{a, 2}, {b, 3}, {c, 2}]

[{a, 2}, {d, 1}, {c, 2}]

[{a, 1}, {b, 4}, {c, 2}]

[{a, 2}, {b, 4}, {c, 2}]

Causality Version Vectors **'Dots' are Events**

Causality 'Dots' are Events

[{a, 2}, {b, 3}, {c, 2}]

Evolution of a Set

Evolution of a Set

[{a, 1}]

{a, 1} Shelly

[{a, 1}]

{a, 1} Shelly

[{a, 1}]

{a, 1} Shelly

[{a, 1}, {b, 3}]

[{a, 1}, {b,3}]

[{a, 1}, {b, 3}]

[{a, 2}, {b, 3}]

[{a, 1}, {b, 3}]

[{a, 2}, {b, 3}]

[{a, 2}, {b, 3}]







Large-scale computation without synchronisation

Large-scale on-line services including social networks and multiplayer games handle huge quantities of frequently changing shared data. Maintaining its consistency is relatively simple in a centralised cloud, but no longer possible due to increased scalability requirements. Instead, data must replicated across several distributed data centres, requiring new principled approaches to consistency that will be explored by the SyncFree project.

SyncFree is a European research project taking place for 3 years, staring October 2013, and is funded by the European Union, grant agreement n°609551.

SyncFree at Cloudscape

Members of the SyncFree project attended the Cloudscape VI even held in Brussels on February 24th-25th. This series of workshops brings together public and private sector IT and legal professionals to discuss current and future issues and opportunities presented by the cloud.

A position paper about SyncFree published at the workshop can be found at the following link:

http://admin.cloudscapeseries.eu/Repository/document/PositionP



☆



Riak 20 Dev Preview http://docs.basho.com/riak/ 2.0.0pre11/downloads/







Strong Consistency



Security





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<u>http://bashojobs.resumator.com</u>







Questions?

@JoelJacobson

joel@basho.com