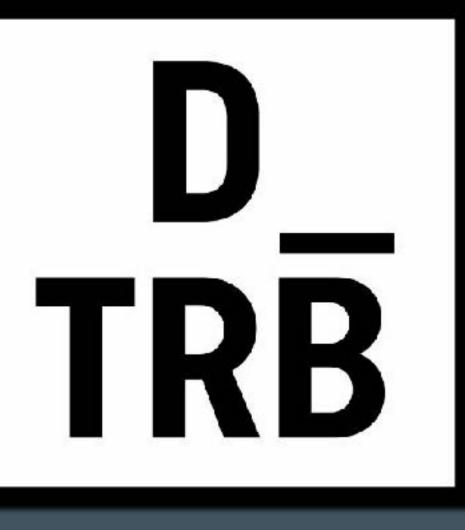
DRIVETRIBE ENGINEERING

A SOCIAL NETWORK ON STREAMS



DRIVETRIBE

The world biggest motoring community.

A social platform for petrolheads.

By Clarkson, Hammond and May.



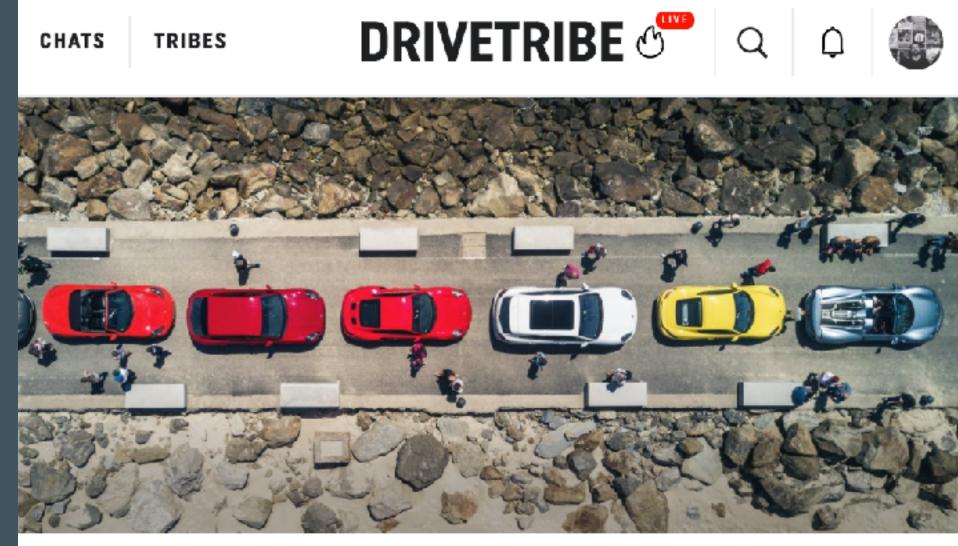






DRIVETRIBE

- A content destination at the core.
- Users consume feeds of content: images, videos, long-form articles.
- Content is organised in homogenous categories called "tribes".
- Different users have different interests and the tribe model allows to mix and match at will.



Porsche

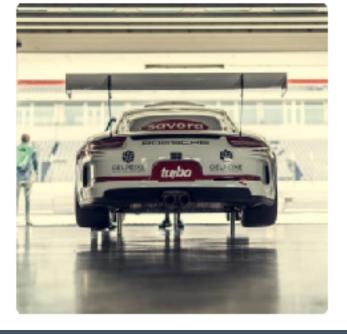
. . .

Official home of Porsche. Written by lucky enthusiasts at Porsche and by Mark

JOINED



CHAT STORIES



Mark Webber • 4 hours ago

"A garage without a Porsche 911 is just a dark hole." - Walter Röhrl

30

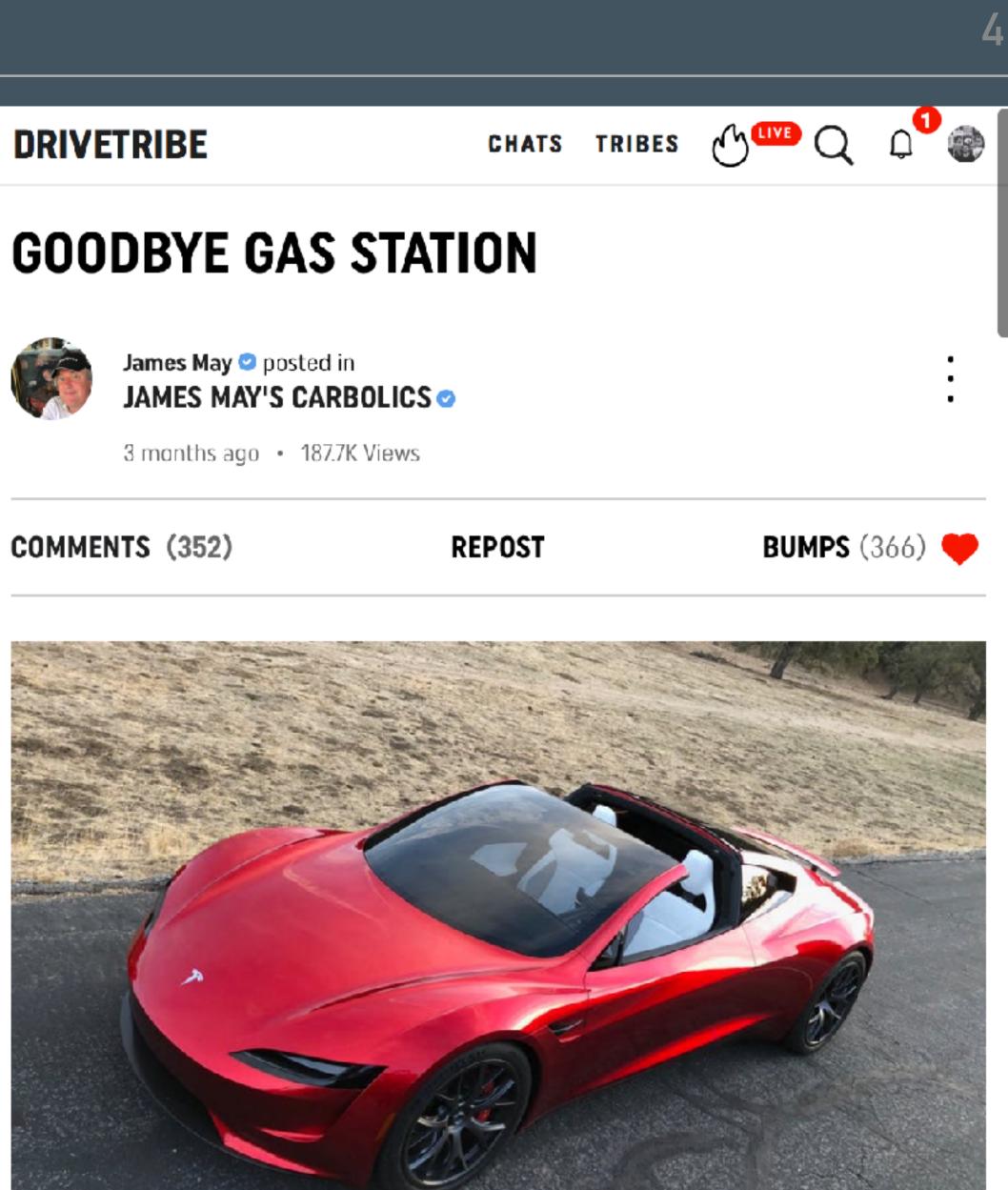




- Single article by James May.
- Contains a plethora of content and engagement information.
- What do we need to compute an aggregate like this?







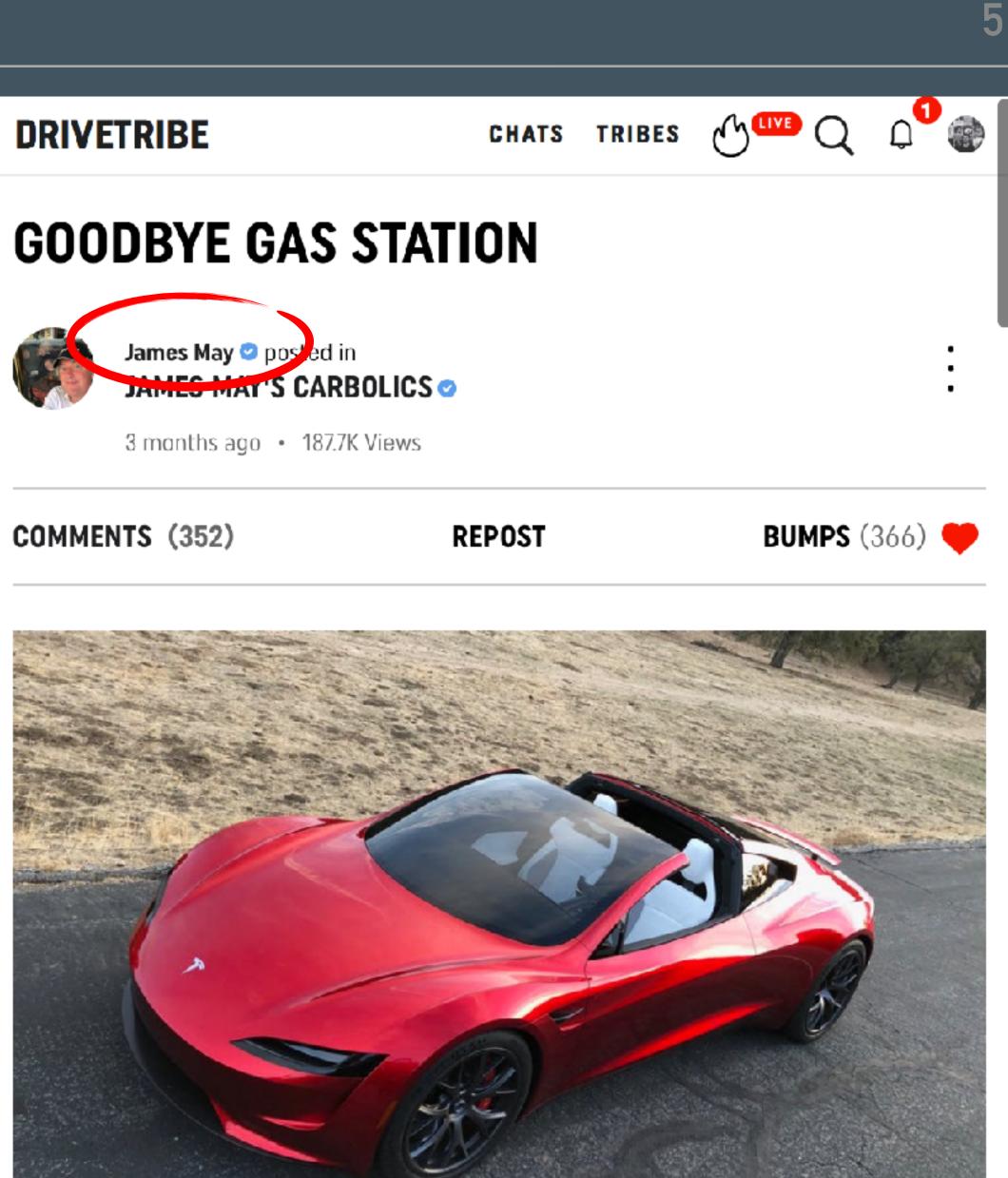
- I'm looking forward to the first scandal, so we can talk about 'Elongate'.

getUser(id: Id[User])









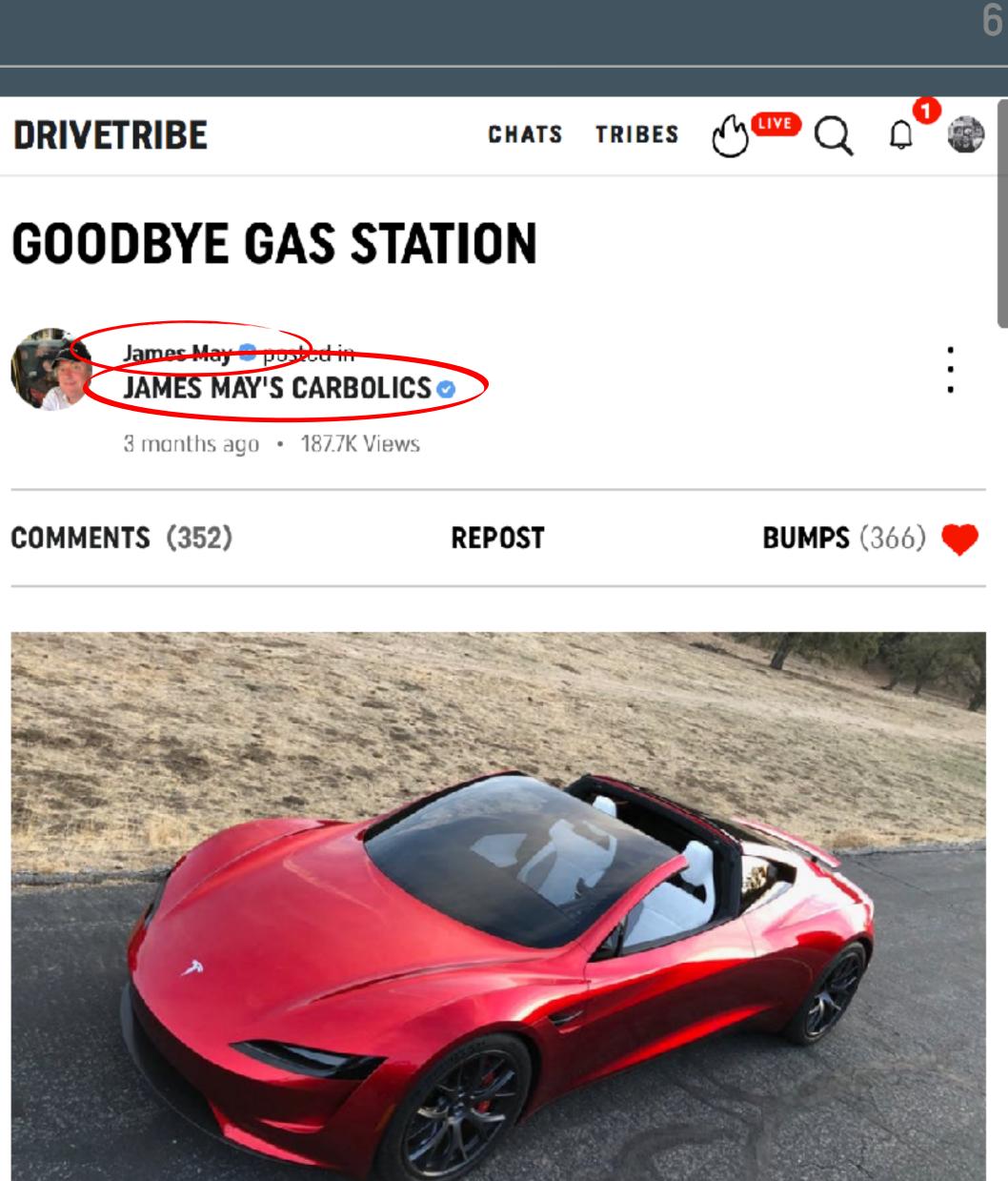
I'm looking forward to the first scandal, so we can talk about 'Elongate'.

getUser(id: Id[User]) getTribe(id: ld[Tribe])



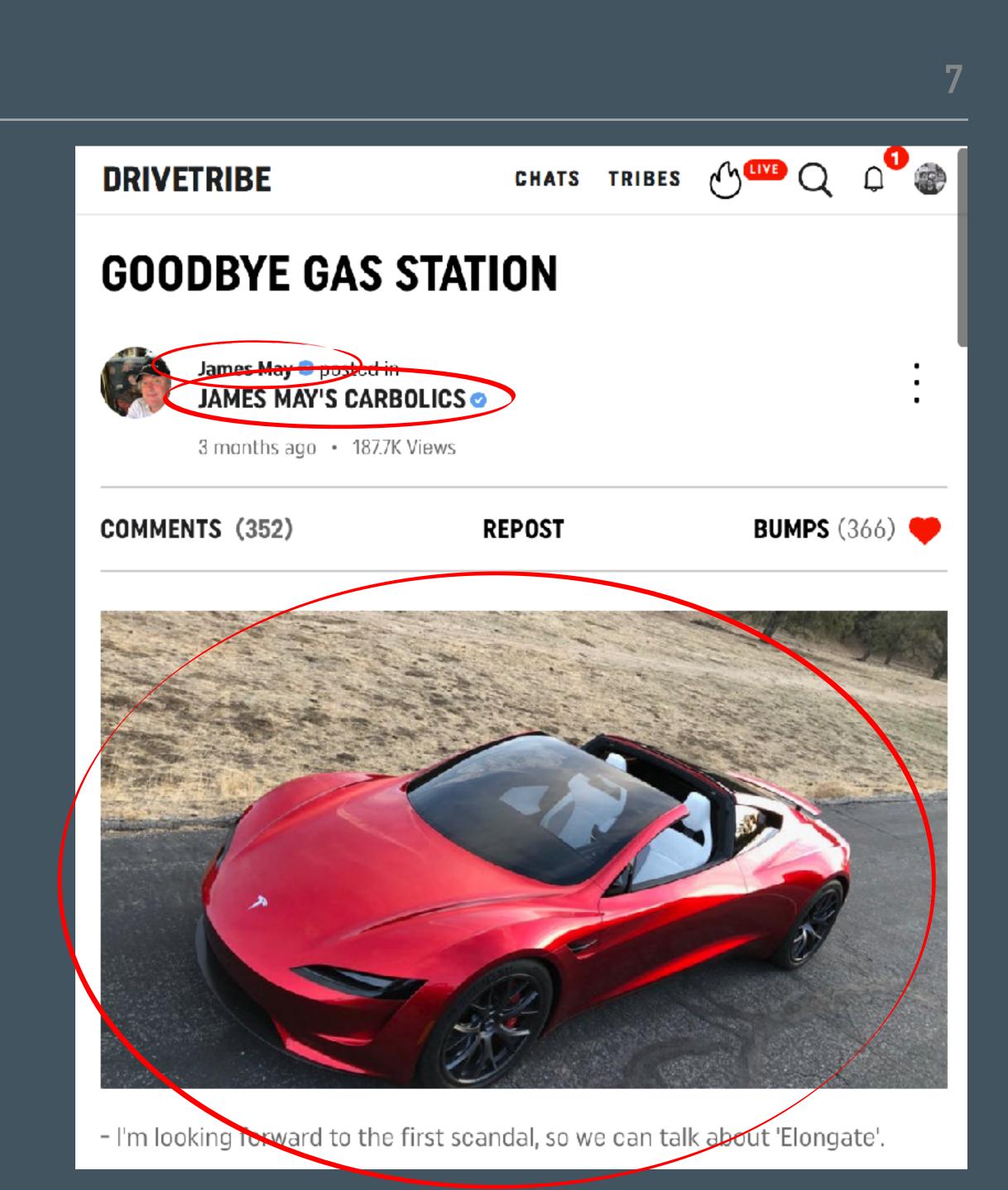




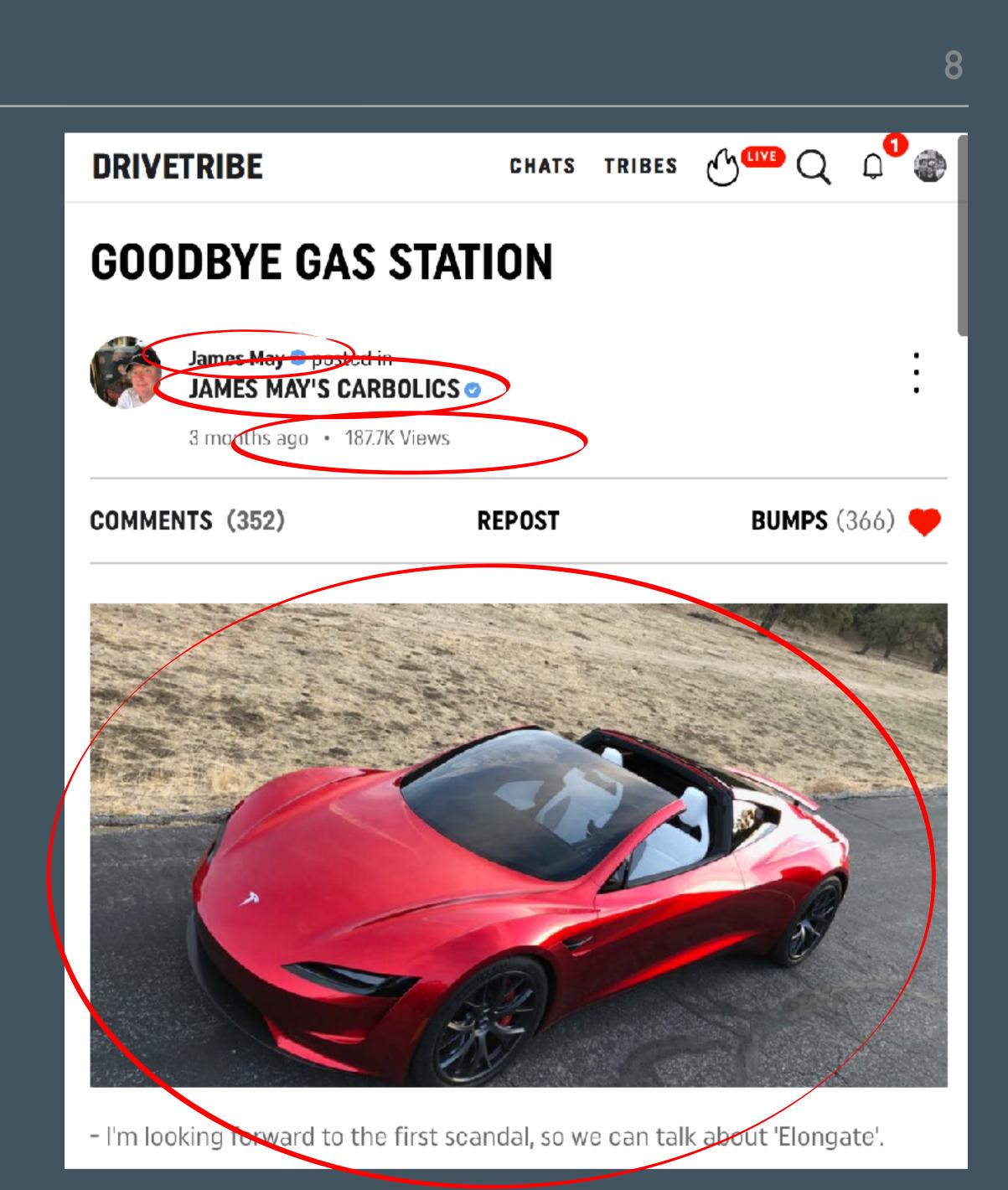


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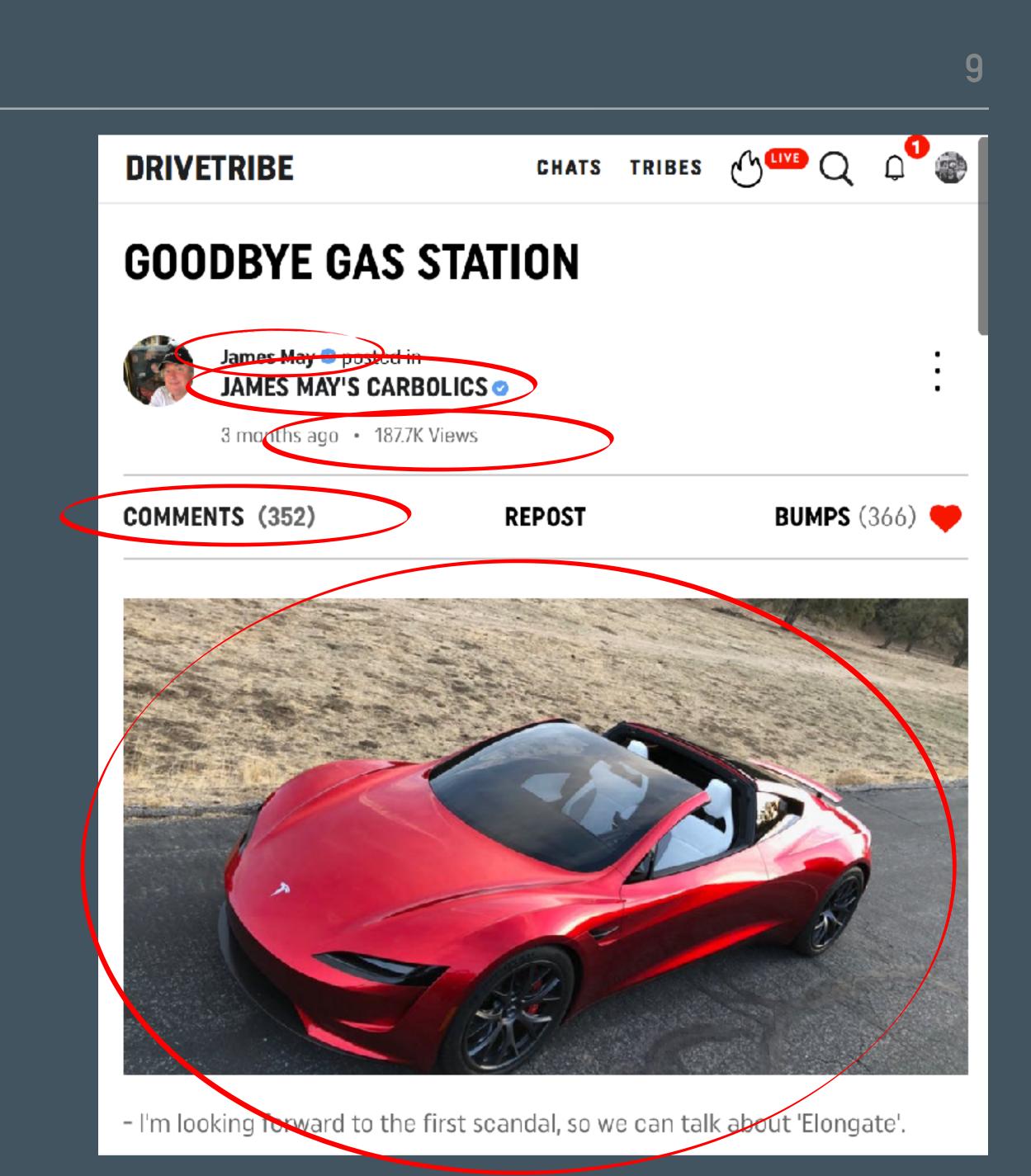
getUser(id: Id[User])
getTribe(id: Id[Tribe])
getArticle(id: Id[Article])



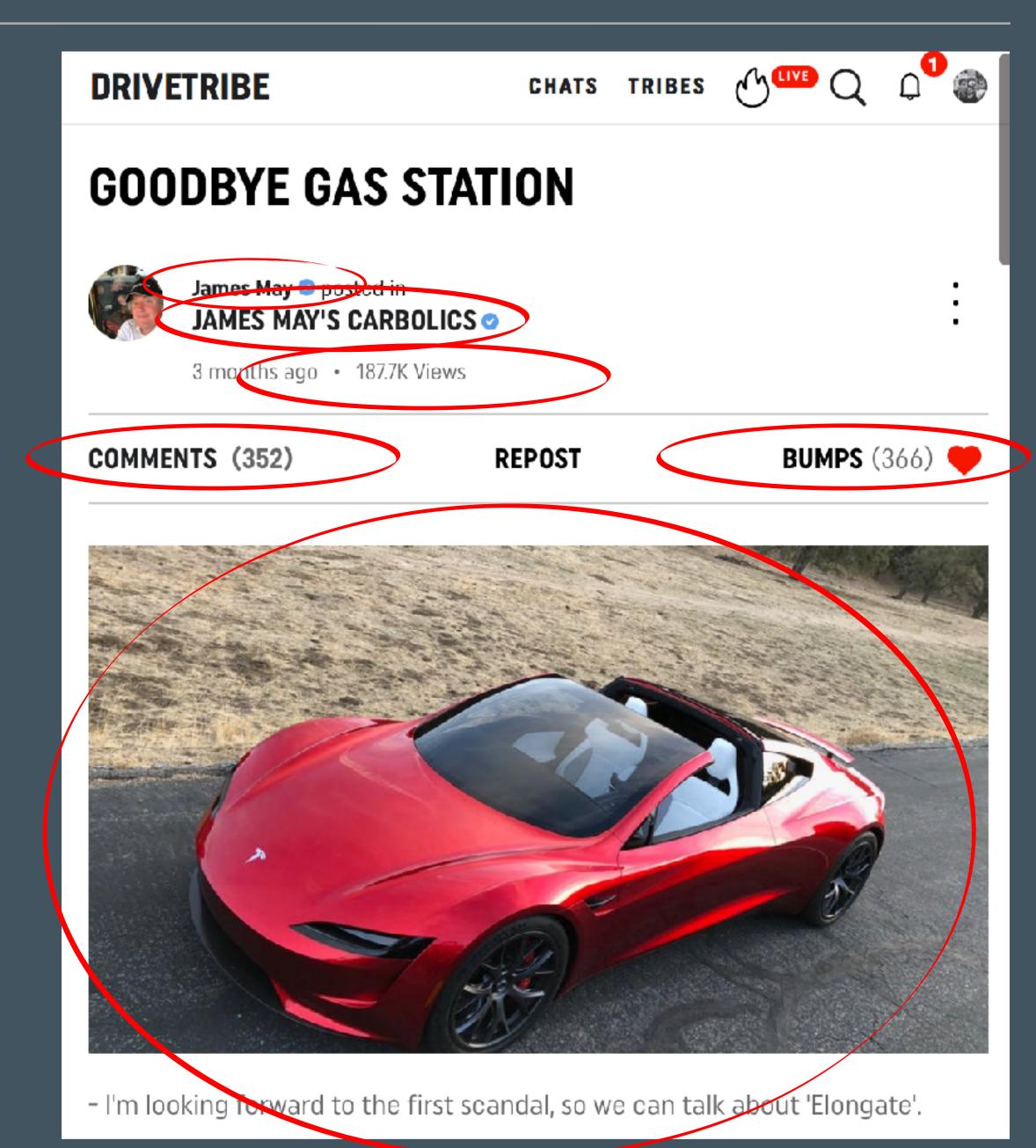
- getUser(id: Id[User])
- getTribe(id: Id[Tribe])
- getArticle(id: Id[Article])
- countViews(id: Id[Article])



- getUser(id: Id[User])
- getTribe(id: Id[Tribe])
- getArticle(id: Id[Article])
- countViews(id: Id[Article])
- countComments(id: Id[Article])



- getUser(id: Id[User])
- getTribe(id: Id[Tribe])
- getArticle(id: Id[Article])
- countViews(id: Id[Article])
- countComments(id: Id[Article])
- countBumps(id: Id[Article])





DRIVETRIBE FEED OF ARTICLES

- rankArticles(forUserId).flatMap { a => ... }
- getUser(id: Id[User])
- getTribe(id: Id[Tribe])
- getArticle(id: ld[Article])
- countViews(id: Id[Article])



Craig Scarborough in Everything Technical (\mathbf{Q})

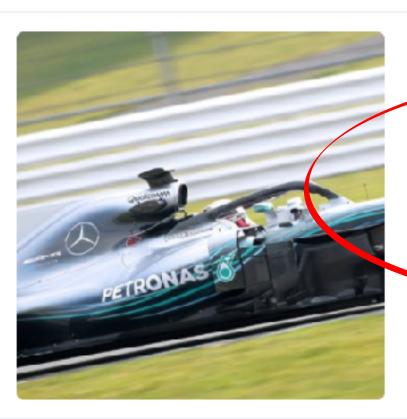
F1: THE FERRARI SF71H IN DETAIL -**PROGRESS ON ALL FRONTS**

FERRARI'S NEWLY RELEASED F1 CAR SHOWS ESSIVE DESIGN IN EVERY AREA IN ORDER TO CHASE MERCEDES FOR THE CROWN IN 2018

James May in James May's Carbolics

THE WEIGHT IS OVER

But it should be under, surely?



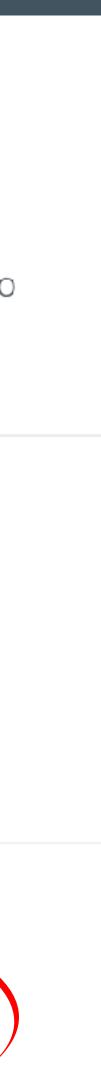
Craig Scarborough in Everything Tech

F1: THE MERCEDES W09 IN DETAIL MAKING THE BEST EVEN BETTER

A deep dive into what's new and what's going to change on Mercedes new F1 challenger

69





QUINTESSENTIAL PREREQUISITES

- it and worry about it later.
- iterate is paramount.
- Maintainable. Spaghetti code works like interest on debt.

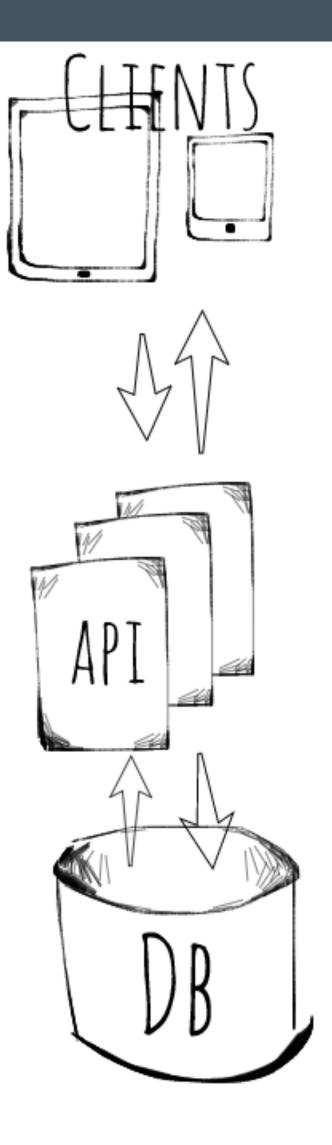
Scalable. Jeremy Clarkson has 7.2M Twitter followers. Cannot really hack

Performant. Low latency is key and mobile networks add quite a bit of it. Flexible. Almost nobody gets it right the first time around. The ability to



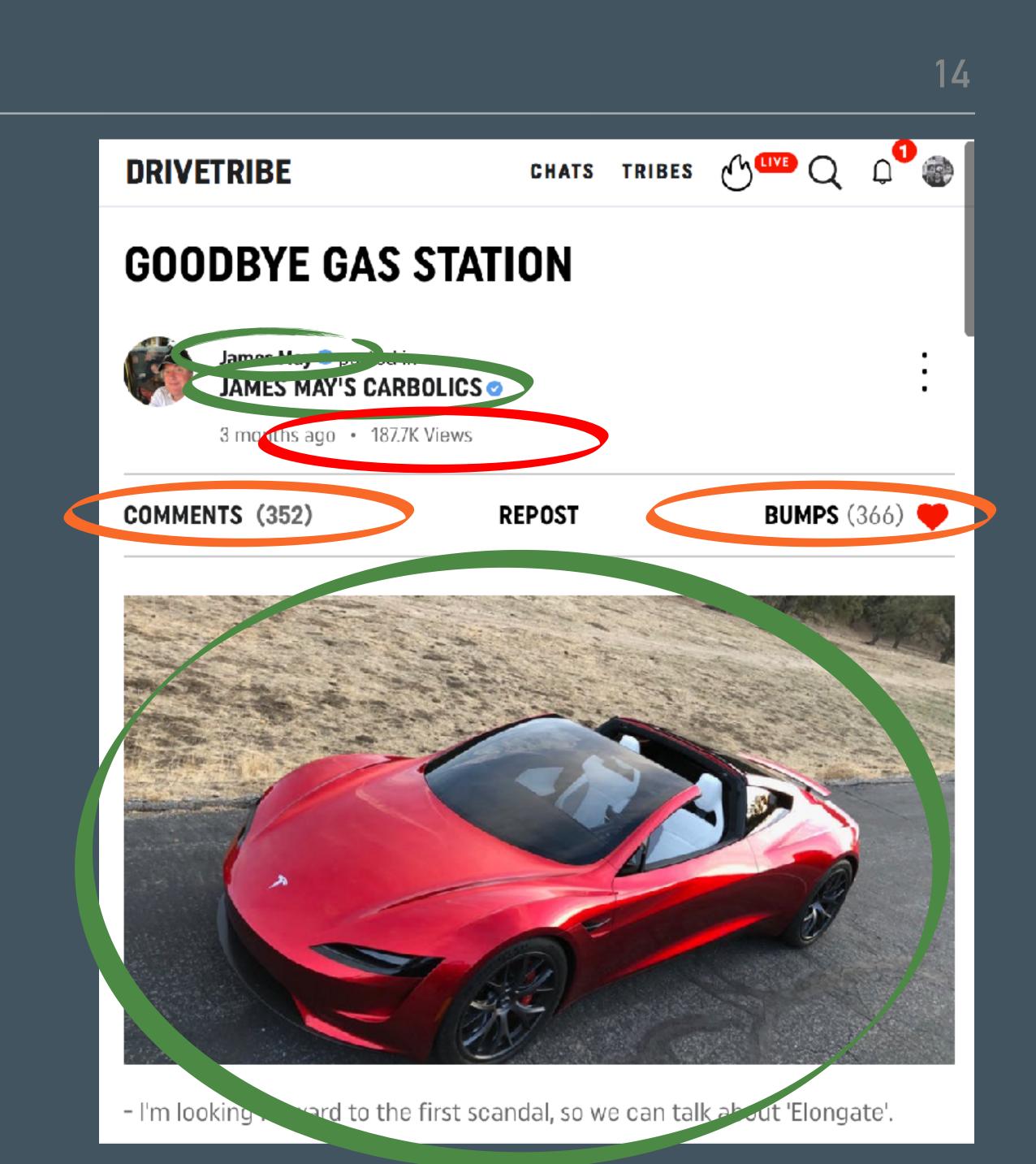
THREE TIER APPROACH

- Clients interact with a fleet of stateless servers (aka "API" servers or "Backend") via HTTP (which is stateless).
- Global shared mutable state (aka the Database).
- Starting simple: Store data in a DB.
- Starting simple: Compute the aggregated views on the fly.





- getUser(id: Id[User])
- getTribe(id: Id[Tribe])
- getArticle(id: Id[Article])
- countComments(id: Id[Article])
- countBumps(id: Id[Article])
- > countViews(id: Id[Article])



READ TIME AGGREGATION

- ► (6 queries per Item) x (Y items per page)
- Cost of ranking and personalisation.
- Quite some work at read time.
- Slow. Not really Performant.



Craig Scarborough in Everything Technical 0

F1: THE FERRARI SF71H IN DETAIL -**PROGRESS ON ALL FRONTS**

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James May in James May's Carbolics 63

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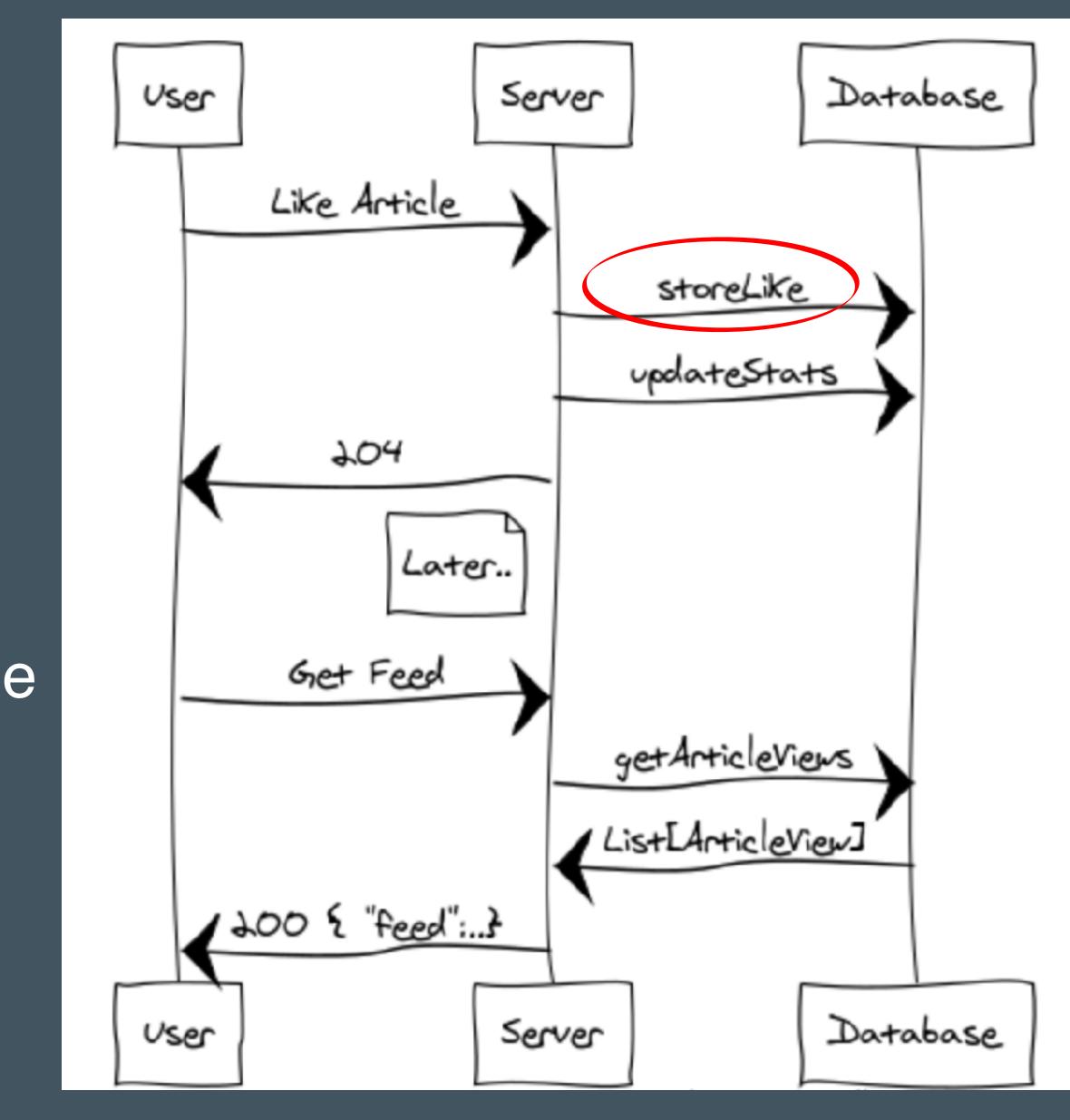
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- Then a single query can fetch all the views at once. That scales.







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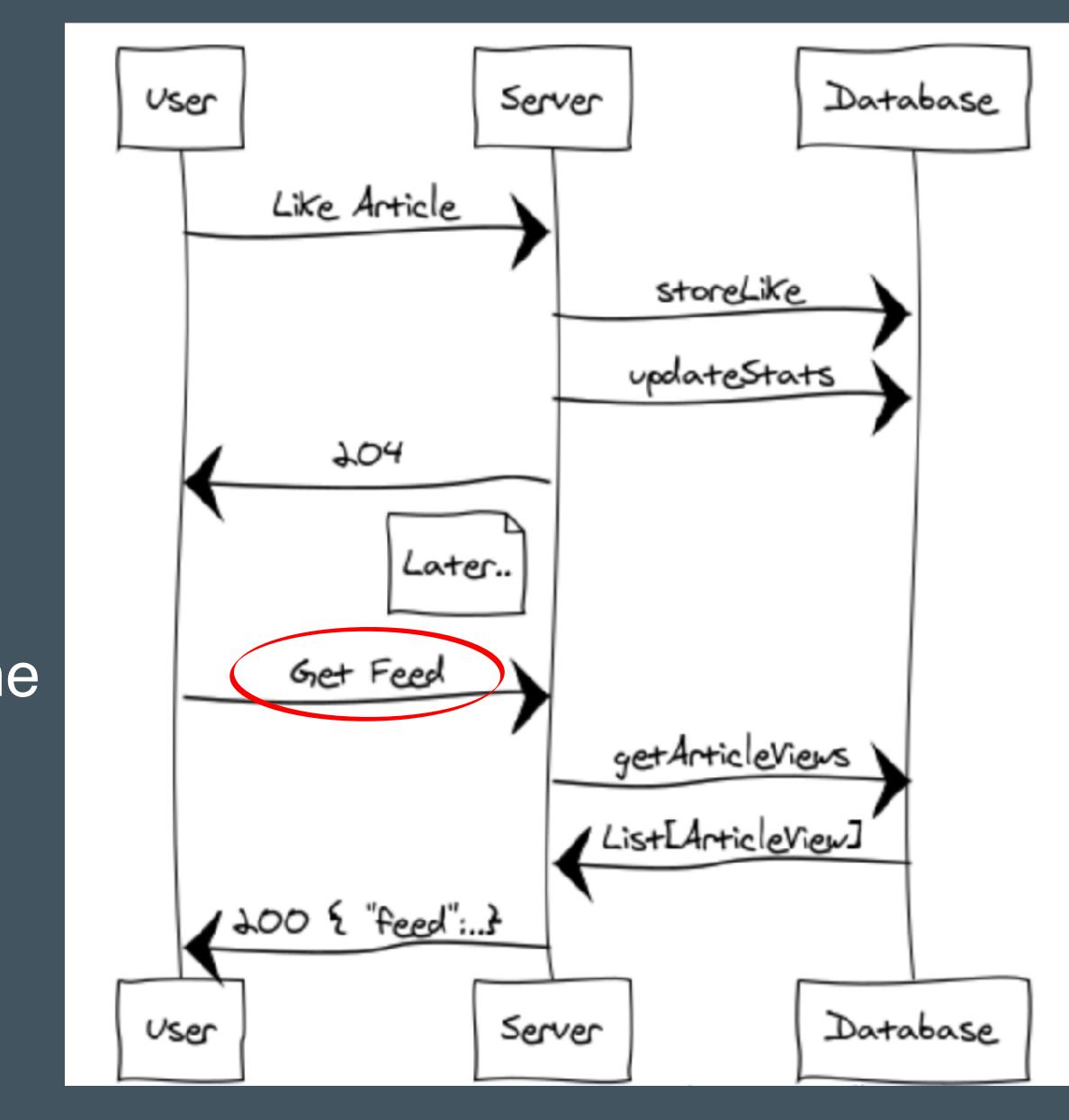
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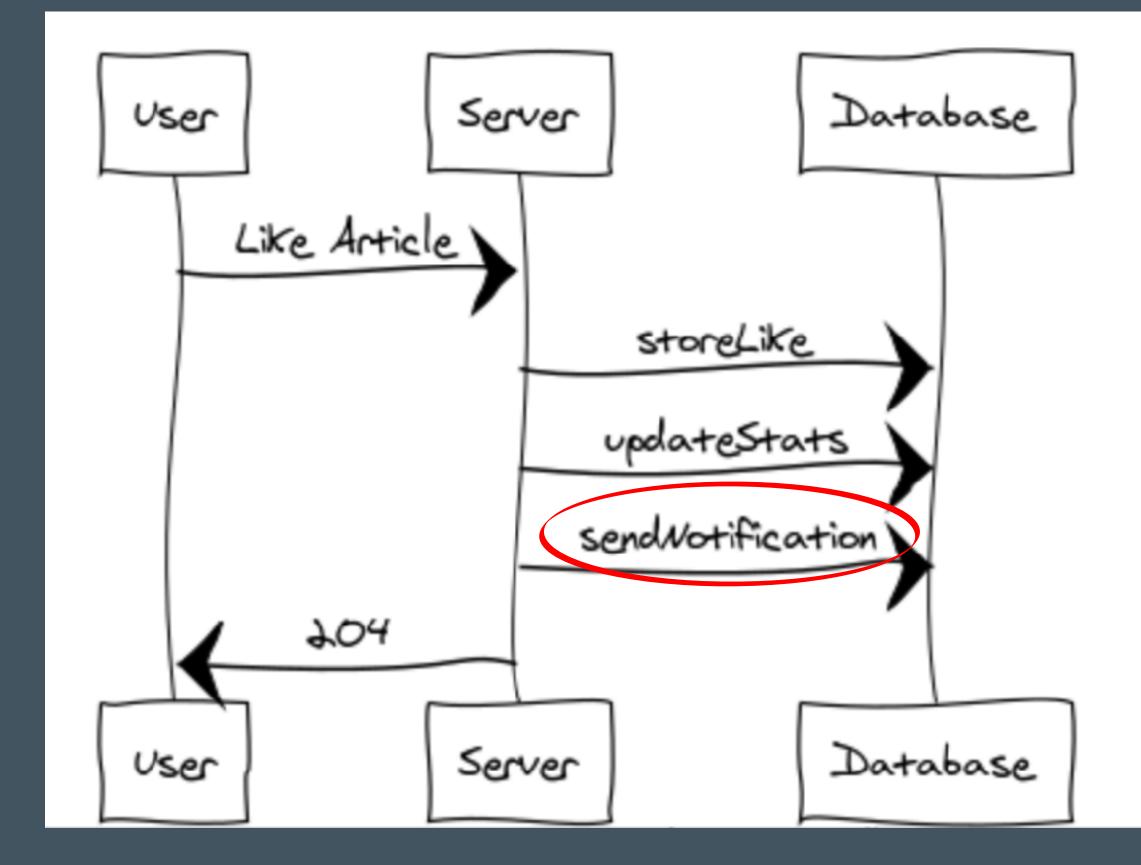






WRITE TIME AGGREGATION - EVOLUTION

sendNotification

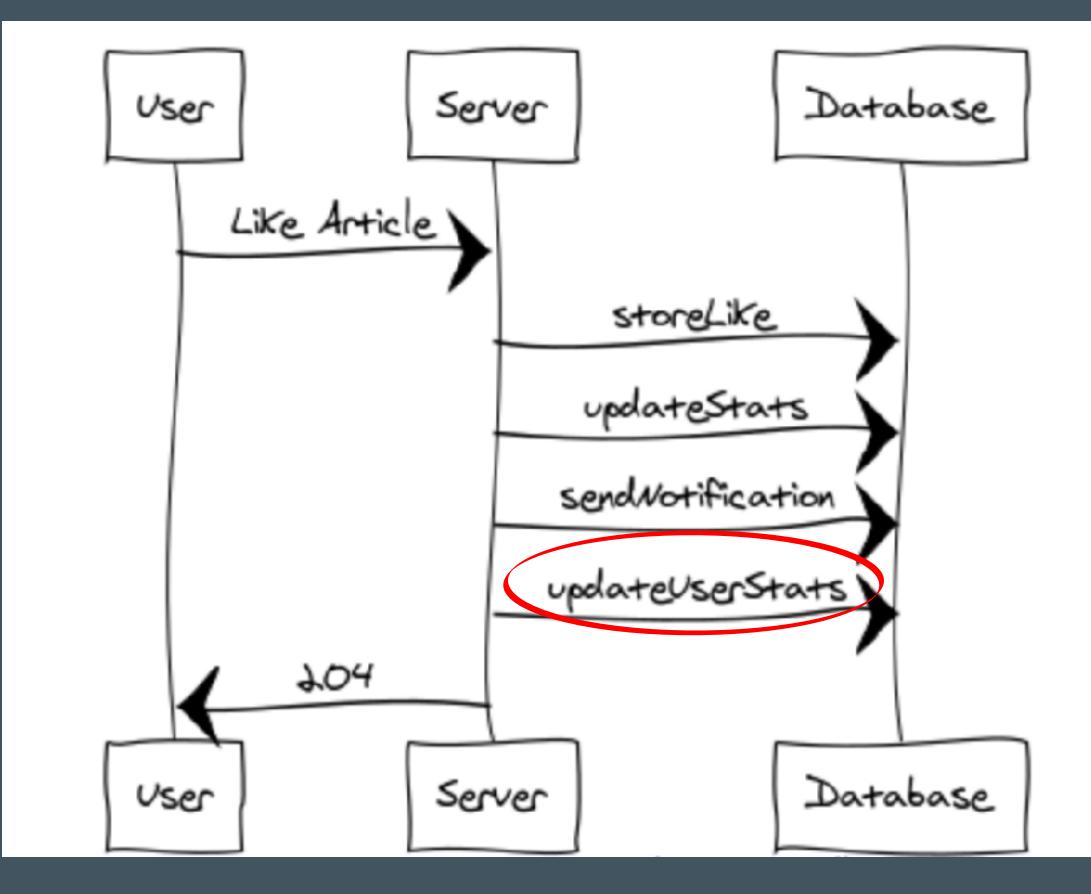






WRITE TIME AGGREGATION - EVOLUTION

sendNotification updateUserStats

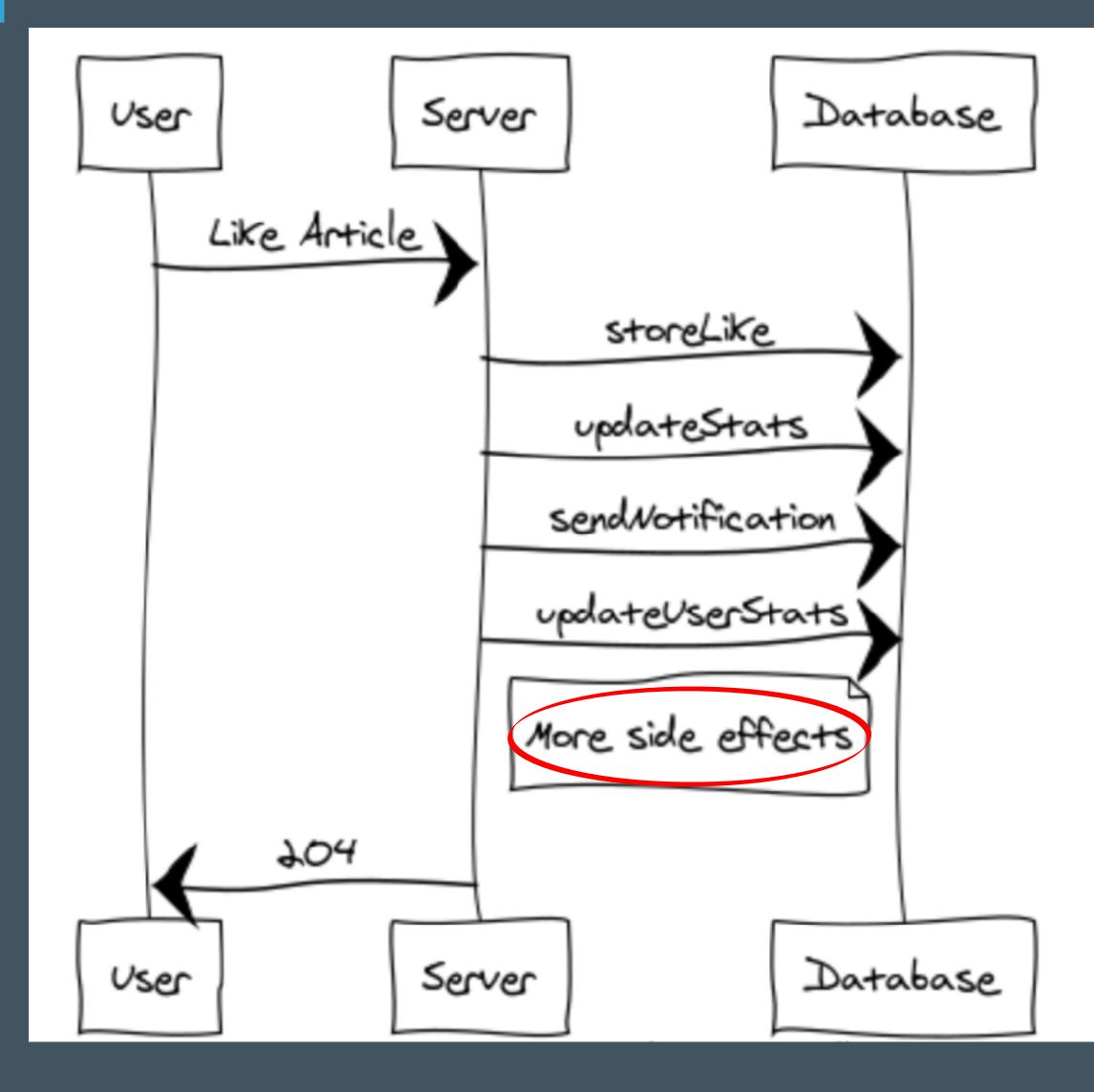






WRITE TIME AGGREGATION - EVOLUTION

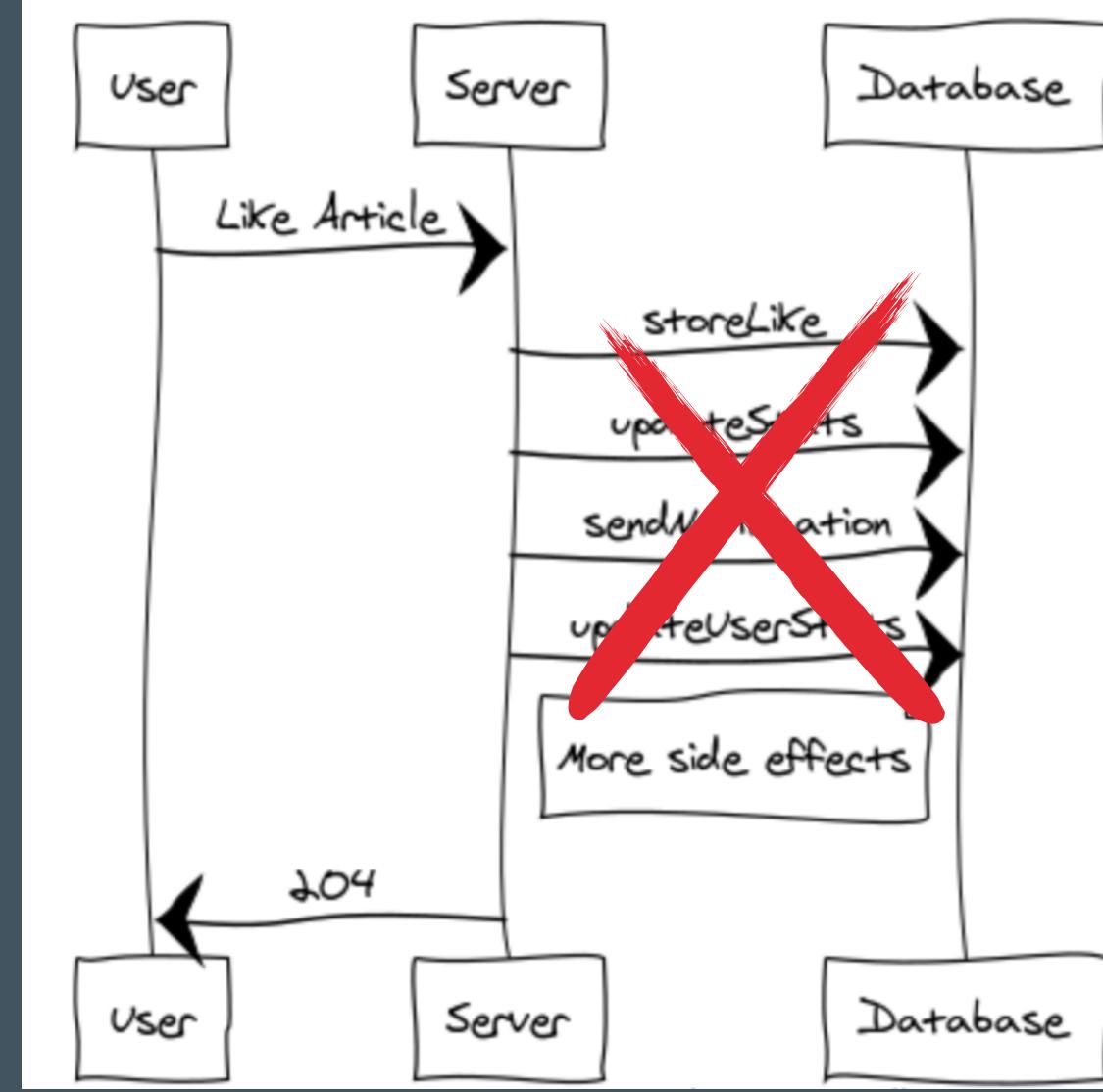
- sendNotification
- updateUserStats.
- What if we have a cache?
- Or a different database for search?







- A simple user action is triggering a potentially endless sequence of side effects.
- Most of which need network IO.
- Many of which can fail.







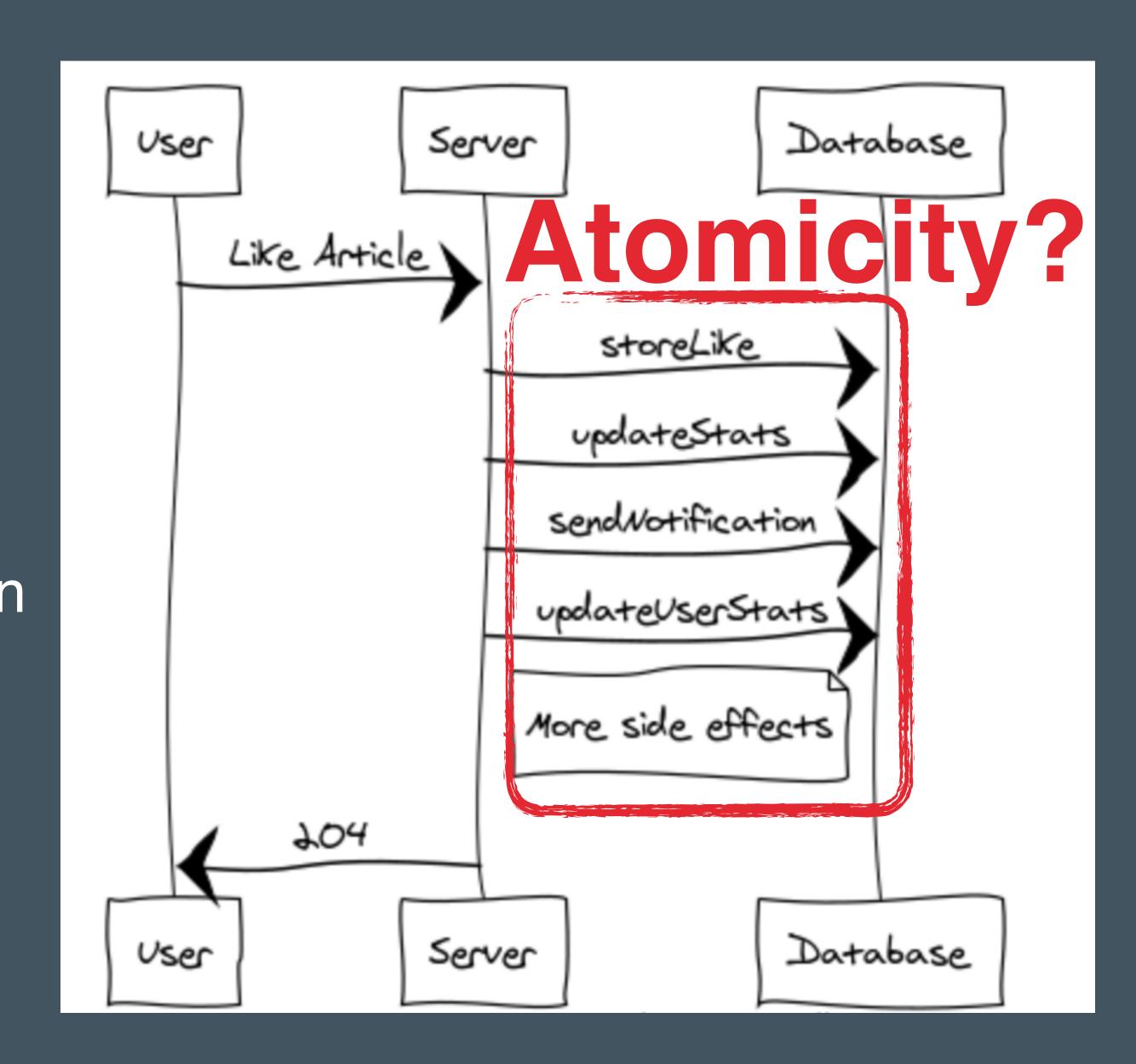


Ατομιζιτγ

What happens if one of them fails? What happens if the server fails in the middle?

We may have transaction support in the DB, but what about external systems?

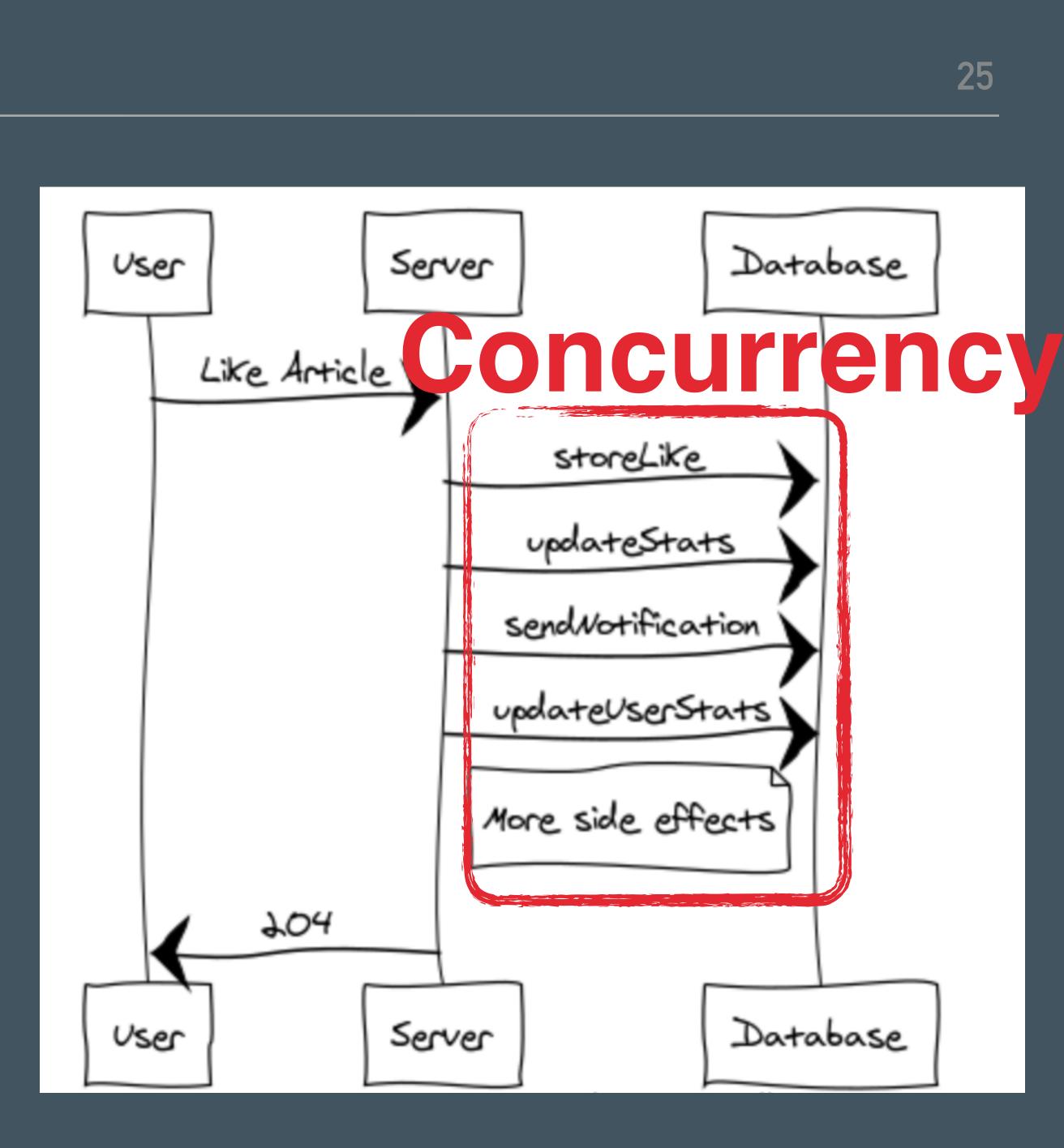
Inconsistent.





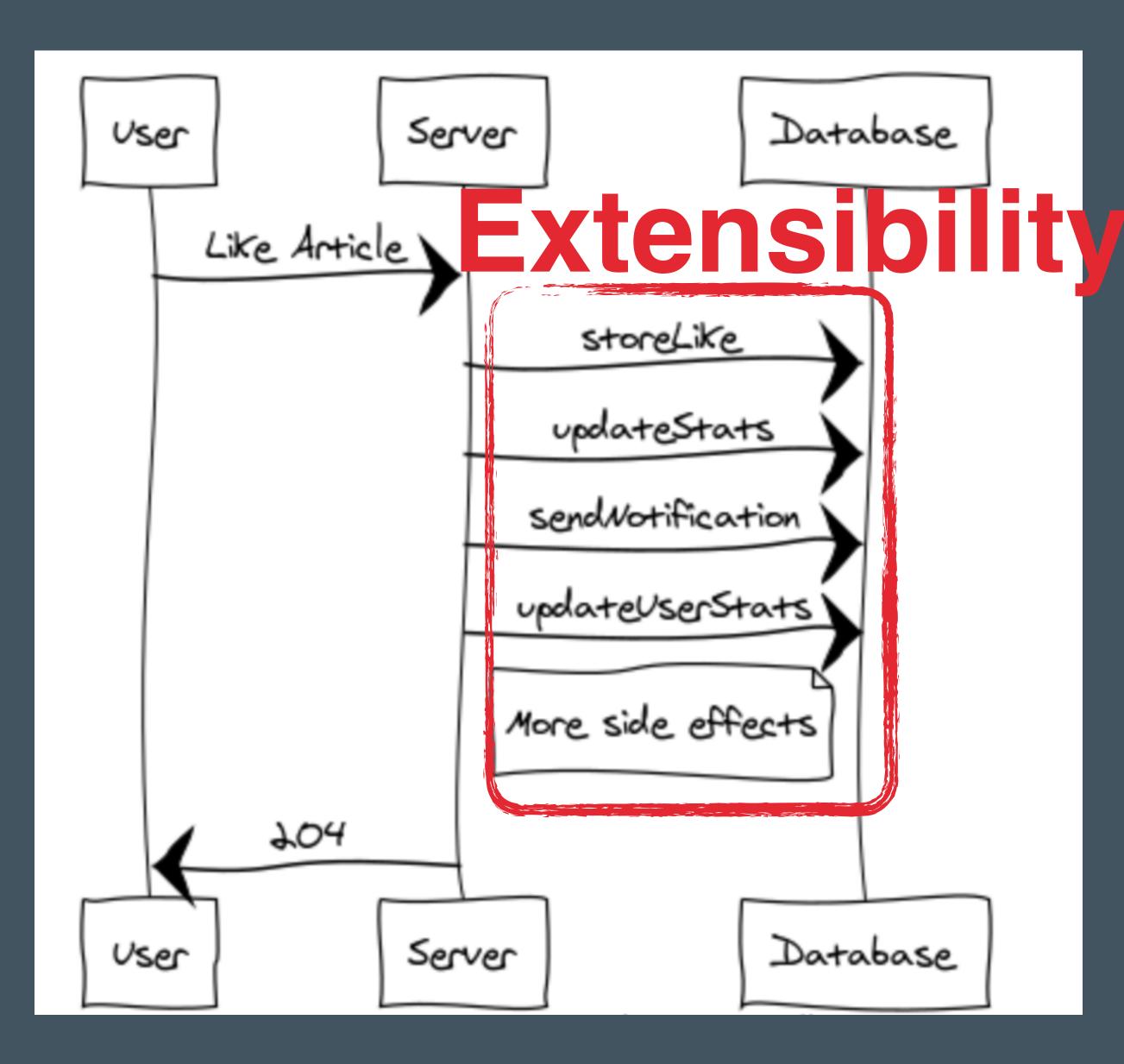
CONCURRENCY

- Concurrent mutations on a global shared state entail race conditions.
- State mutations are destructive and can not be (easily) undone.
- A bug can corrupt the data permanently.



ITALIAN PASTA

- Model evolution becomes difficult. Reads and writes are tightly coupled.
- Migrations are scary.
- This is neither Extensible nor Maintainable.





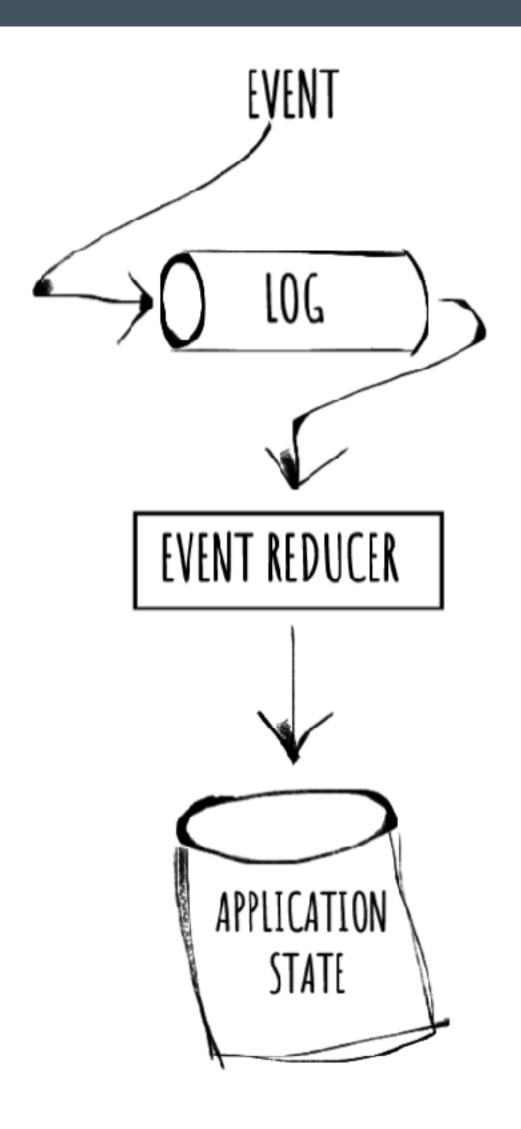
DIFFERENT APPROACH

- Let's take a step back and try to decouple things.
- Clients send events to the API: "John liked Jeremy's post", "Maria updated her profile"
- Events are immutable. They capture a user action at some point in time.
 Every application state instance can be modelled as a projection of those
- Every application state instance ca events.



- Persisting those yields an appendonly log of events.
- An event reducer can then produce application state instances.
- Even retroactively. The log is immutable.
- > This is event sourcing.

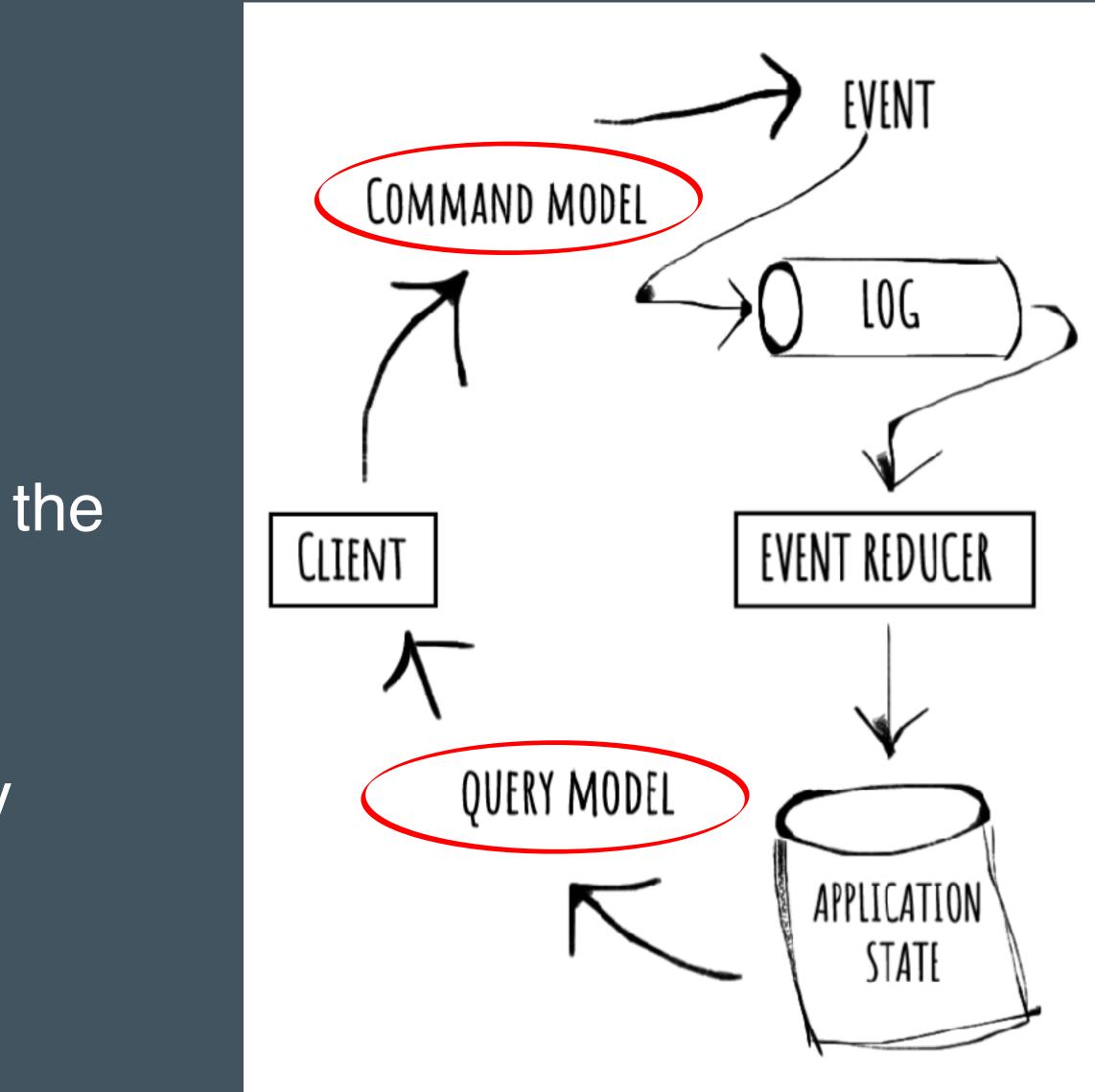






- The write-time model (command model) and the read time model (query model) can be separated.
- Decoupling the two models opens the door to more efficient, custom implementations.

This is known as Command Query Responsibility Segregation aka CQRS.





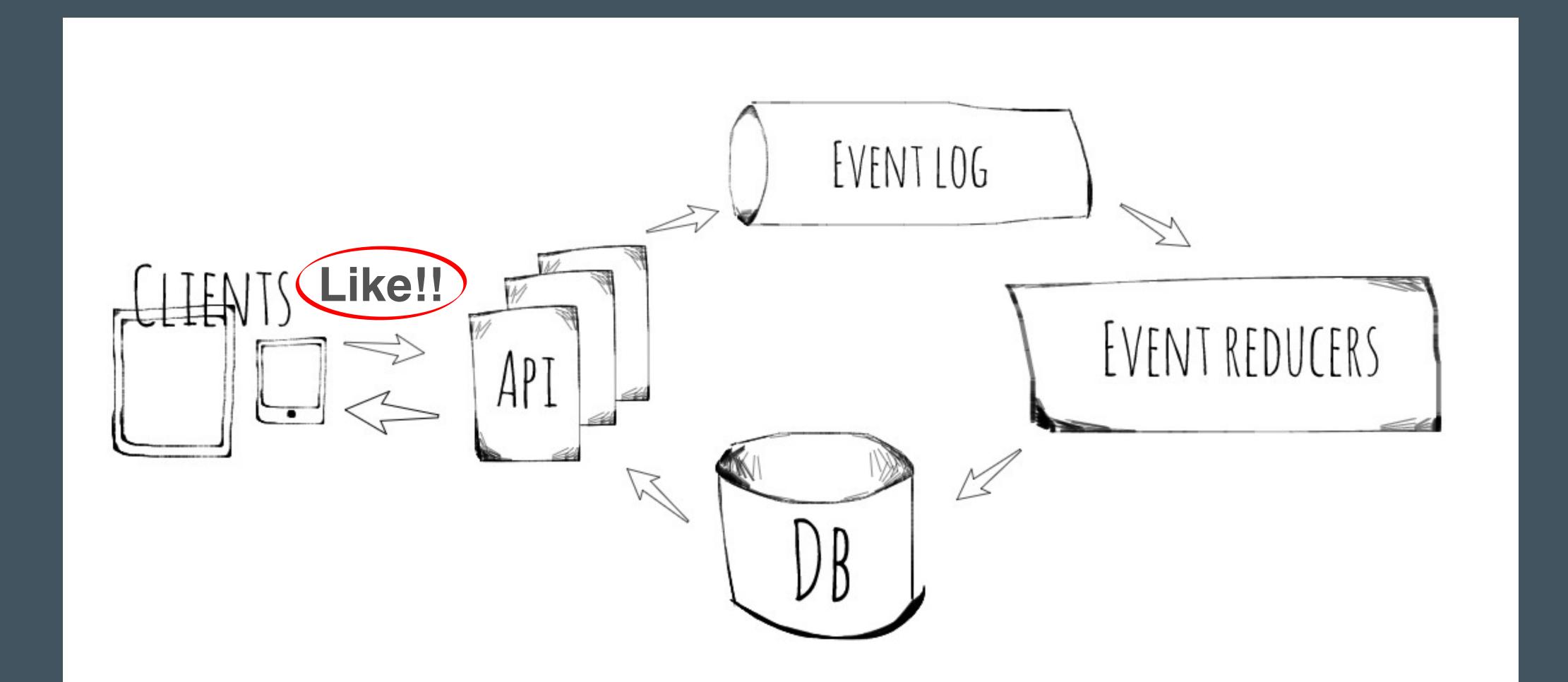








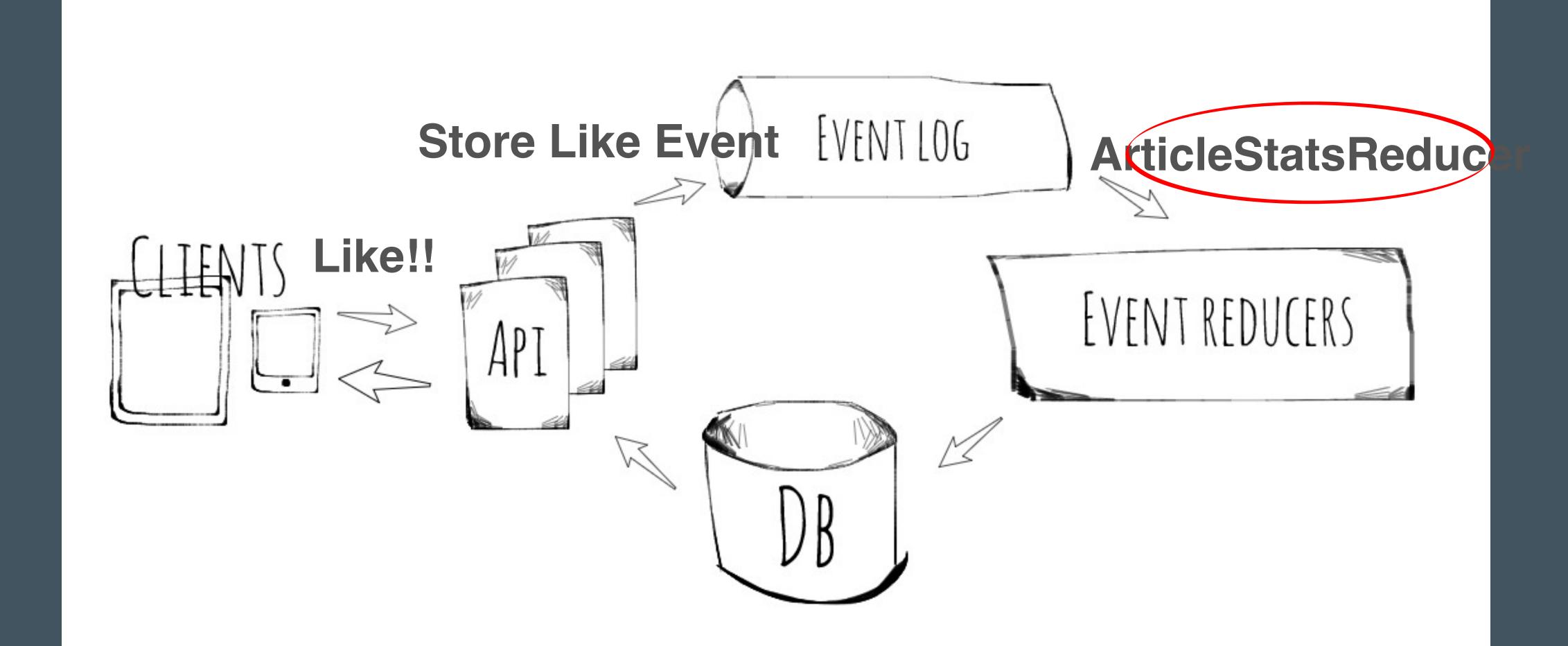




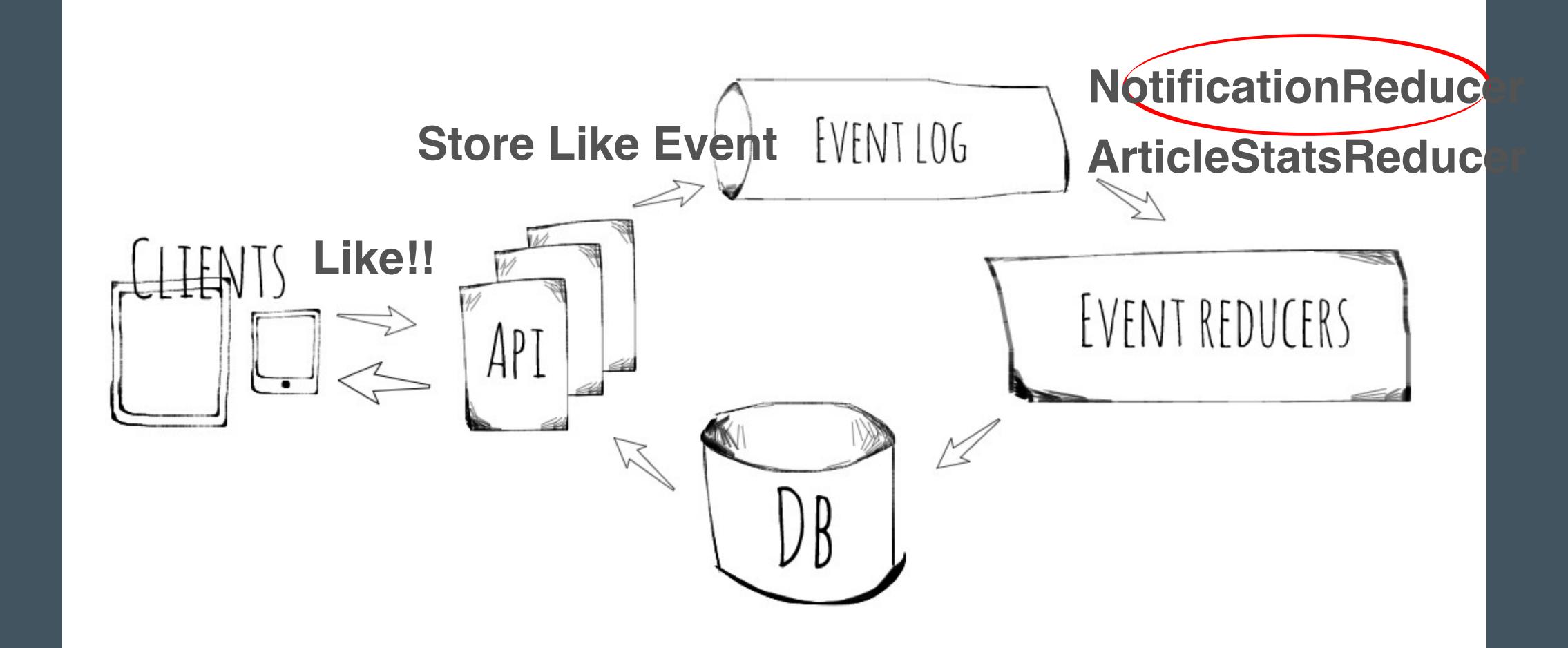




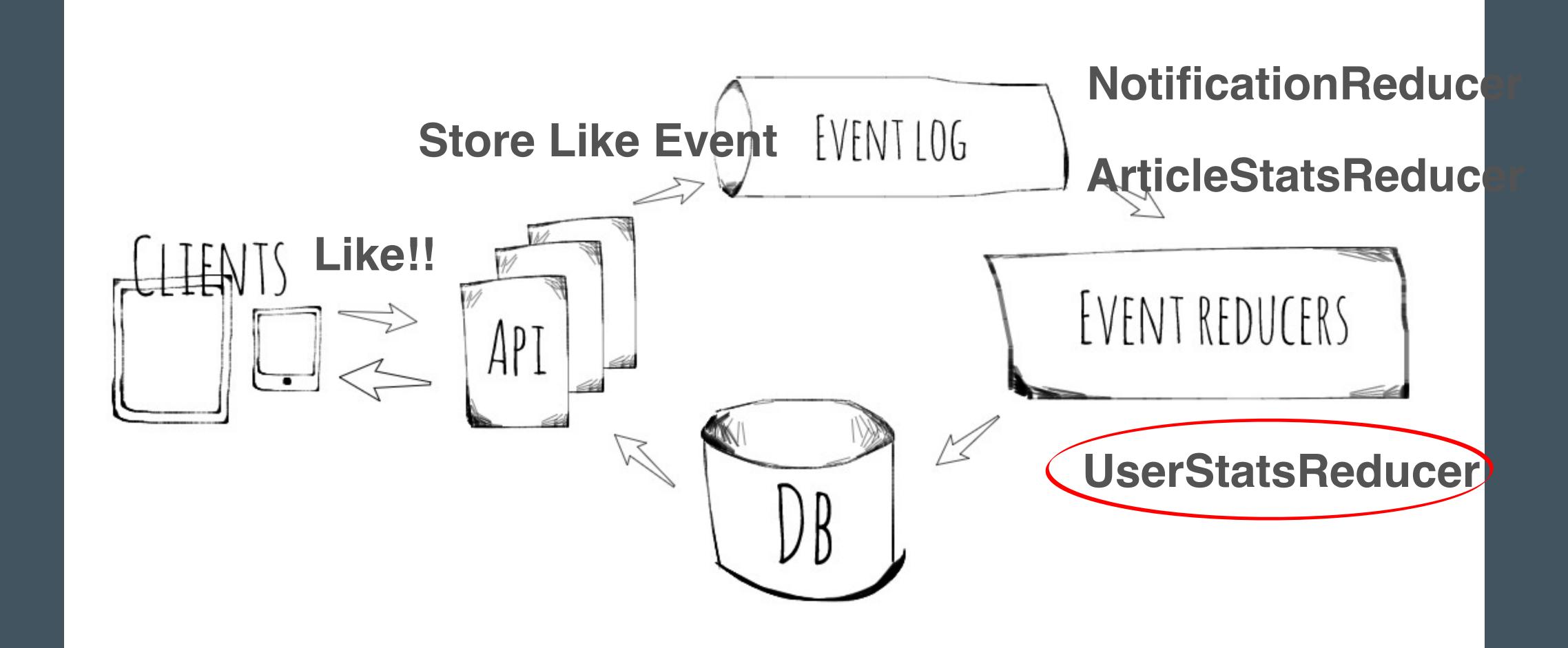




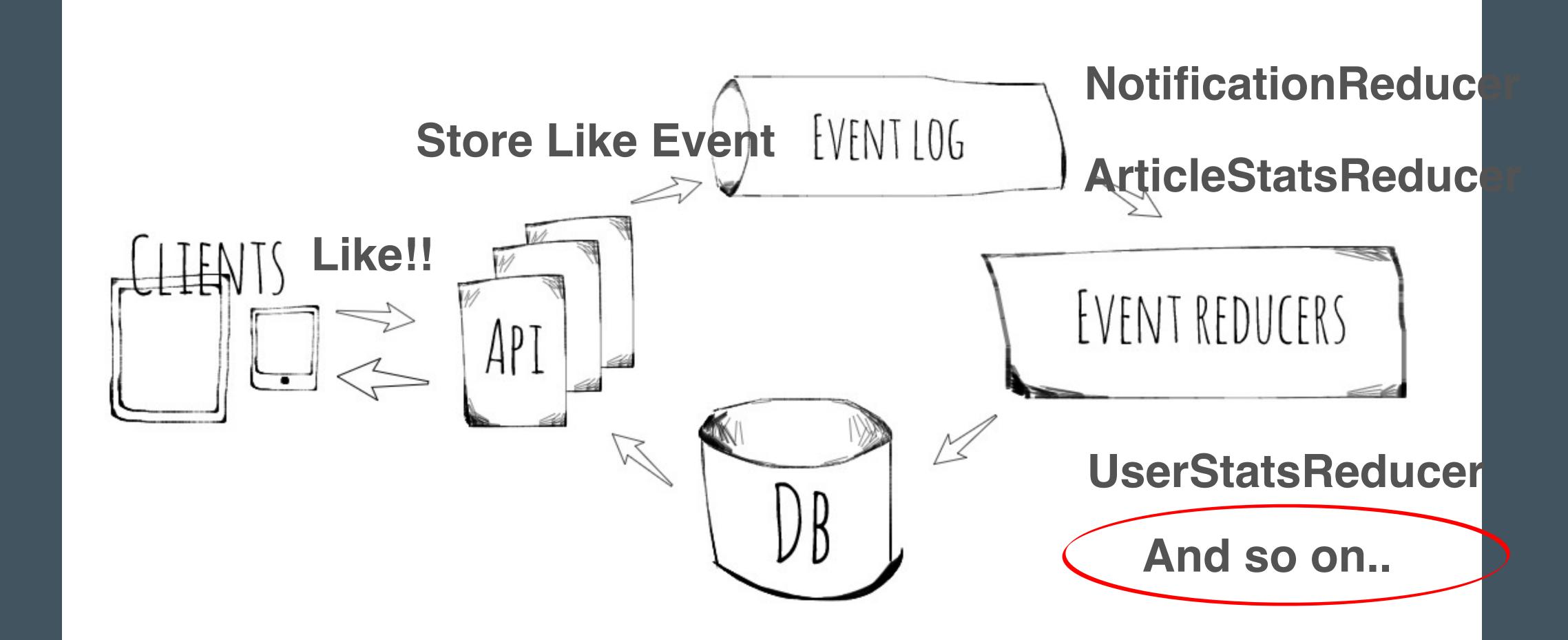




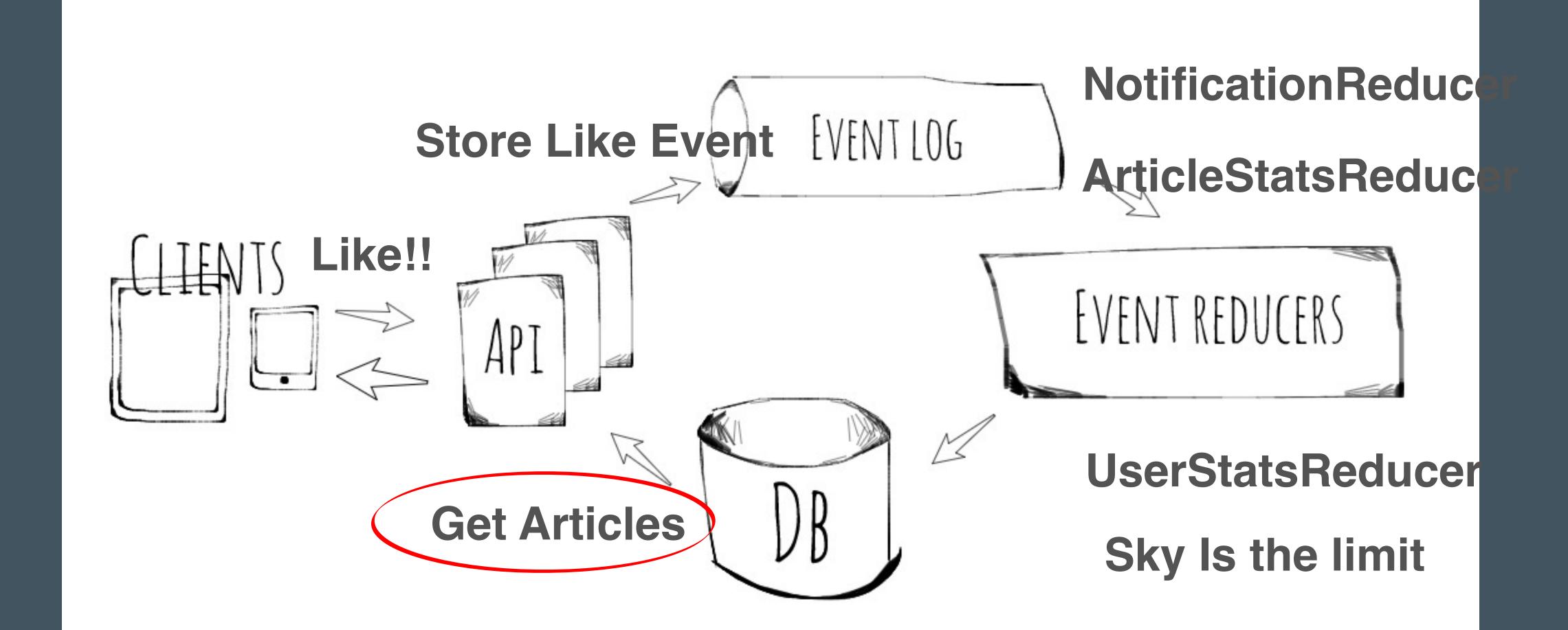








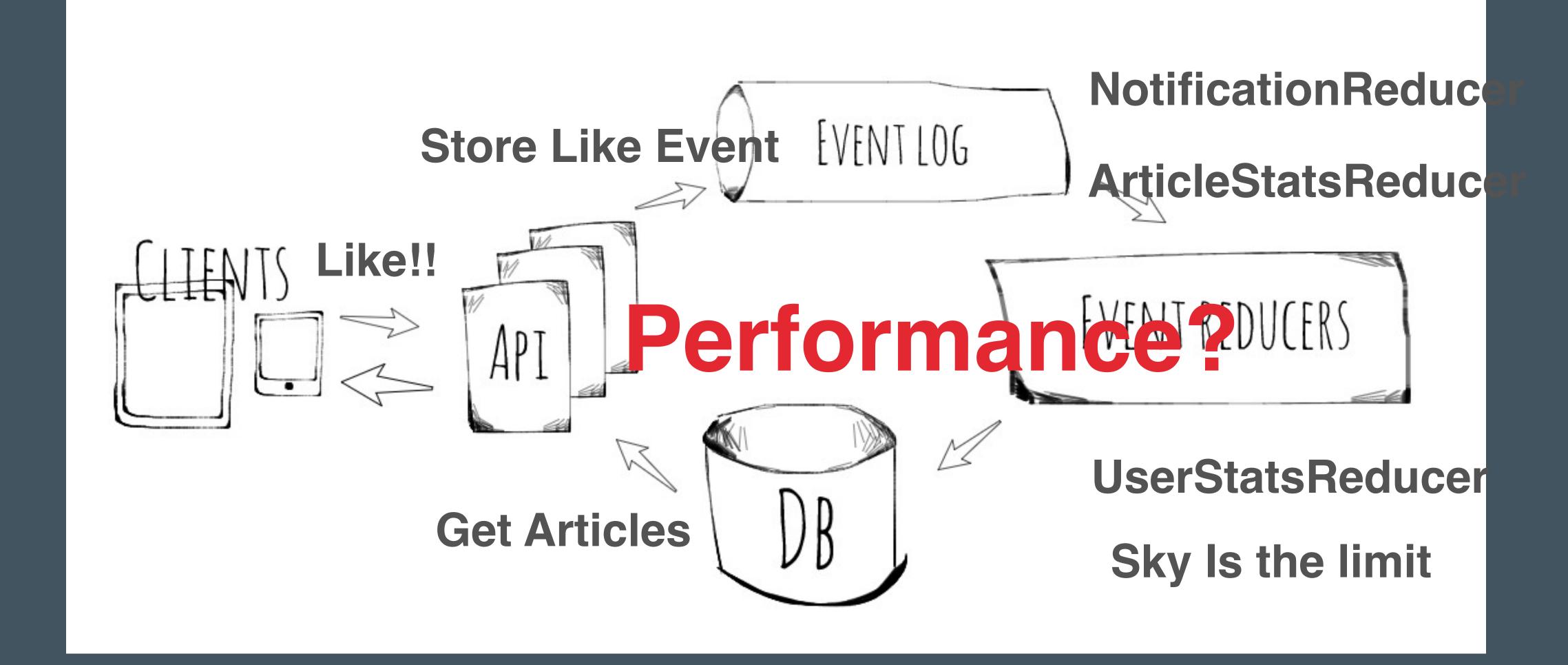




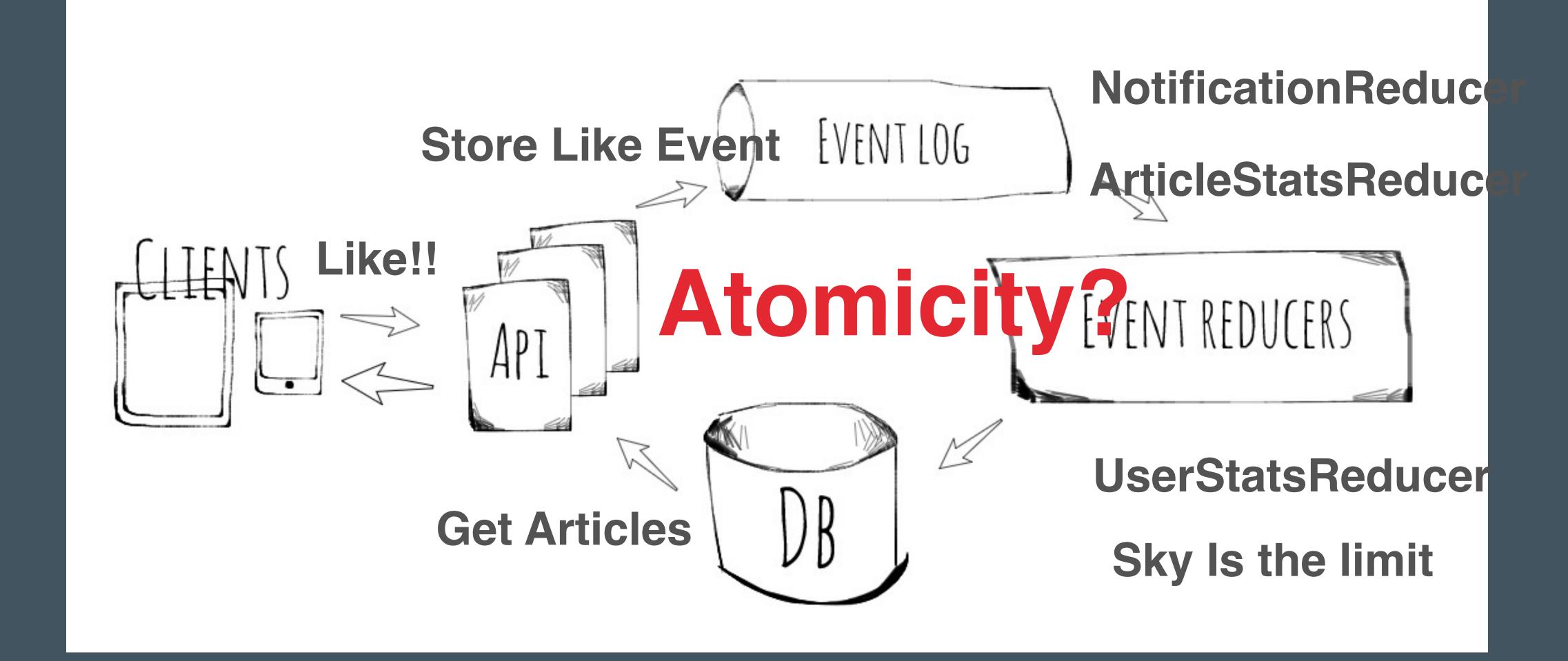




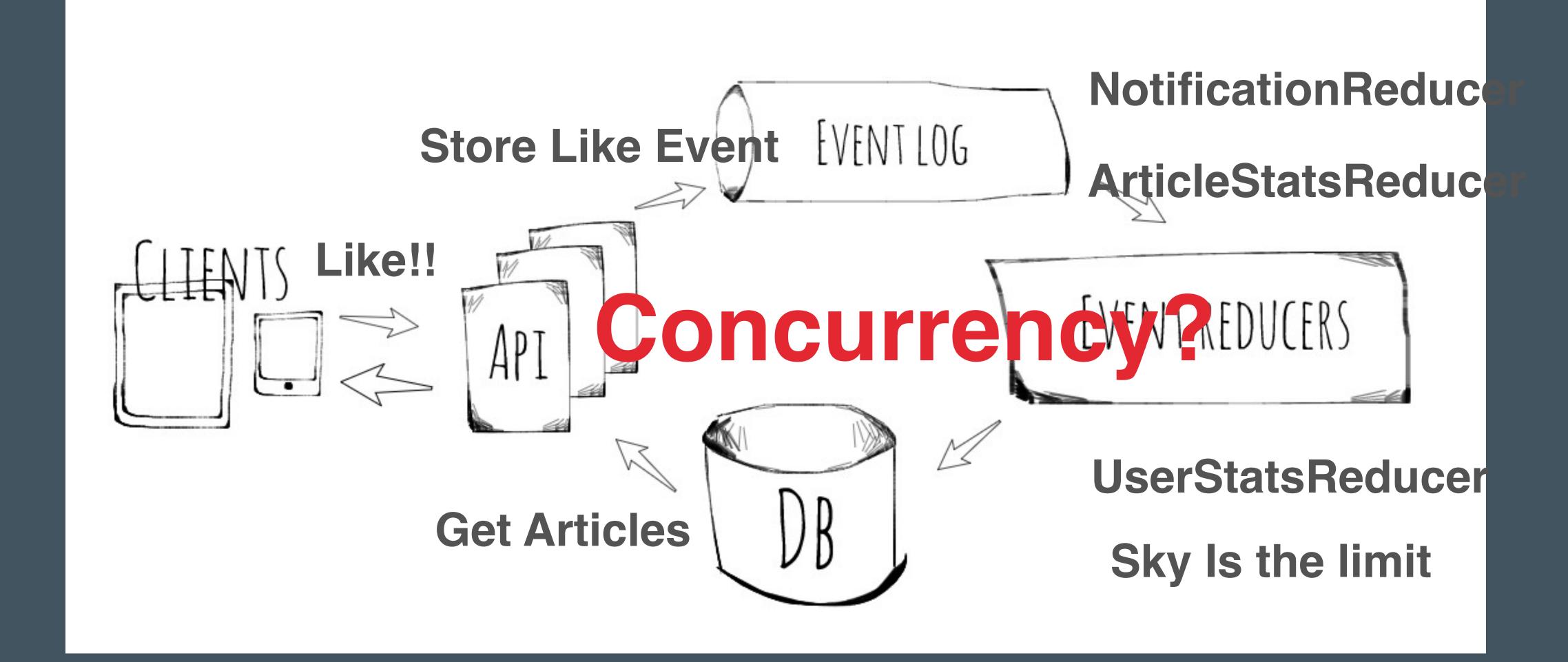














IMPLEMENTATION?

WE NEED A LOG



Distributed, fault-tolerant, durable and fast append-only lo Can scale the thousands of nodes, producers and consur Each business event type can be stored in its own topic.

APACHE KAFKA

နိုင်



afka

WE NEED A STREAM PROCESSOR



Scalable, performant, mature. Elegant high level APIs in Scala. Powerful low level APIs for advanced tuning. Multiple battle-tested integrations. Very nice and active community.

APACHE FLINK





WENEED DATASTORE



Horizontally scalable document store. Rich and expressive query language. Dispensable. Can be replaced.

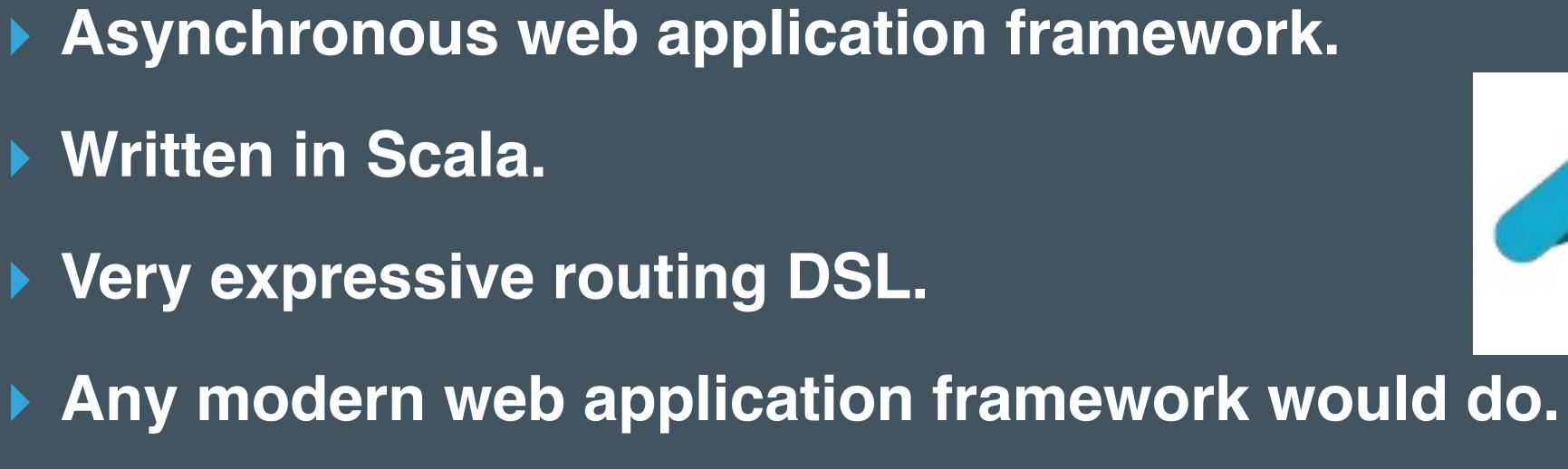
ELASTICSEARCH





WE NEED AN API



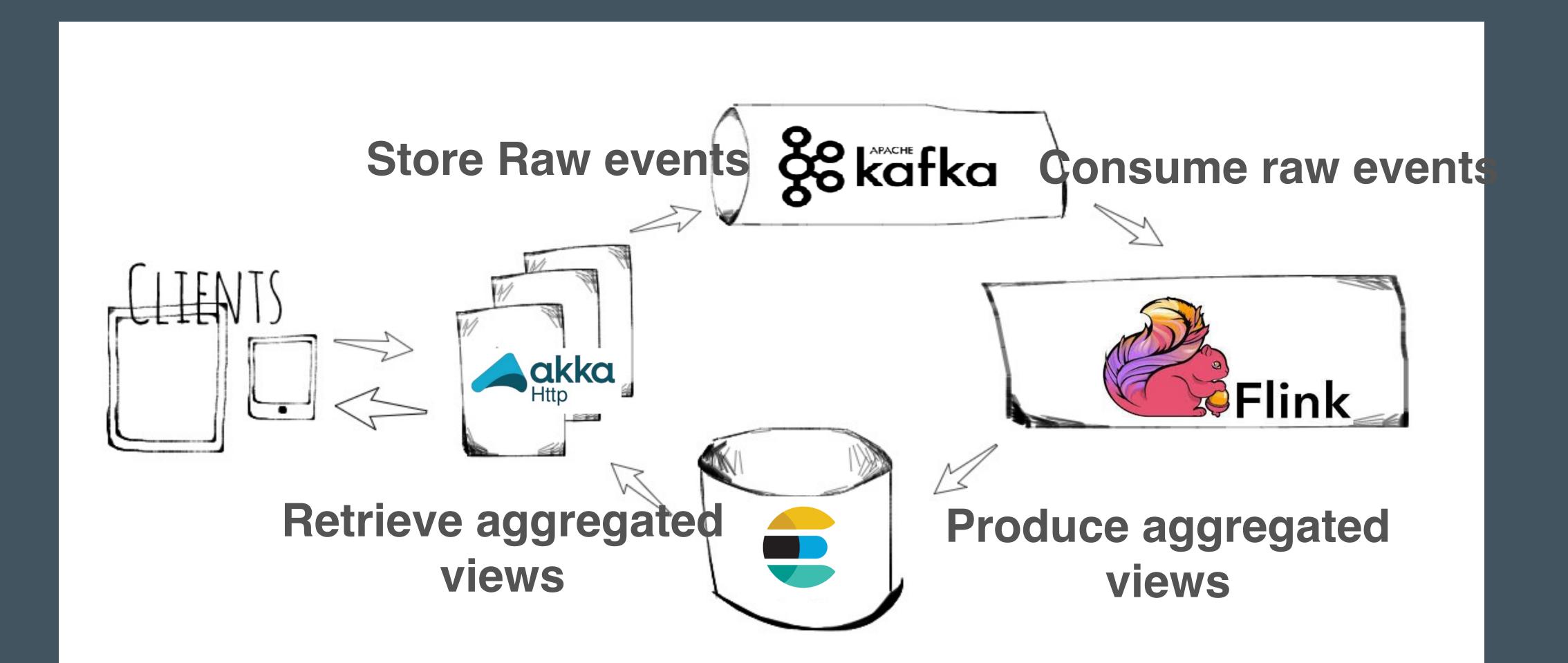


AKKA HTTP

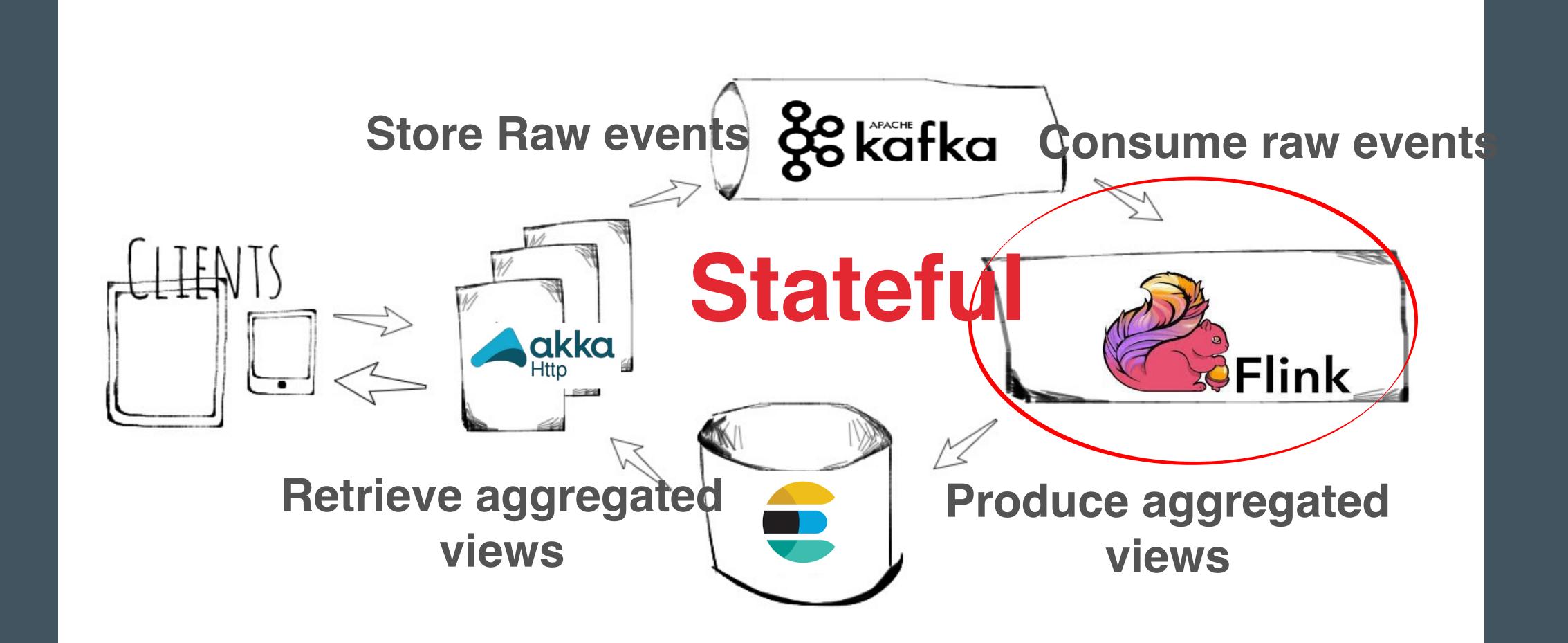




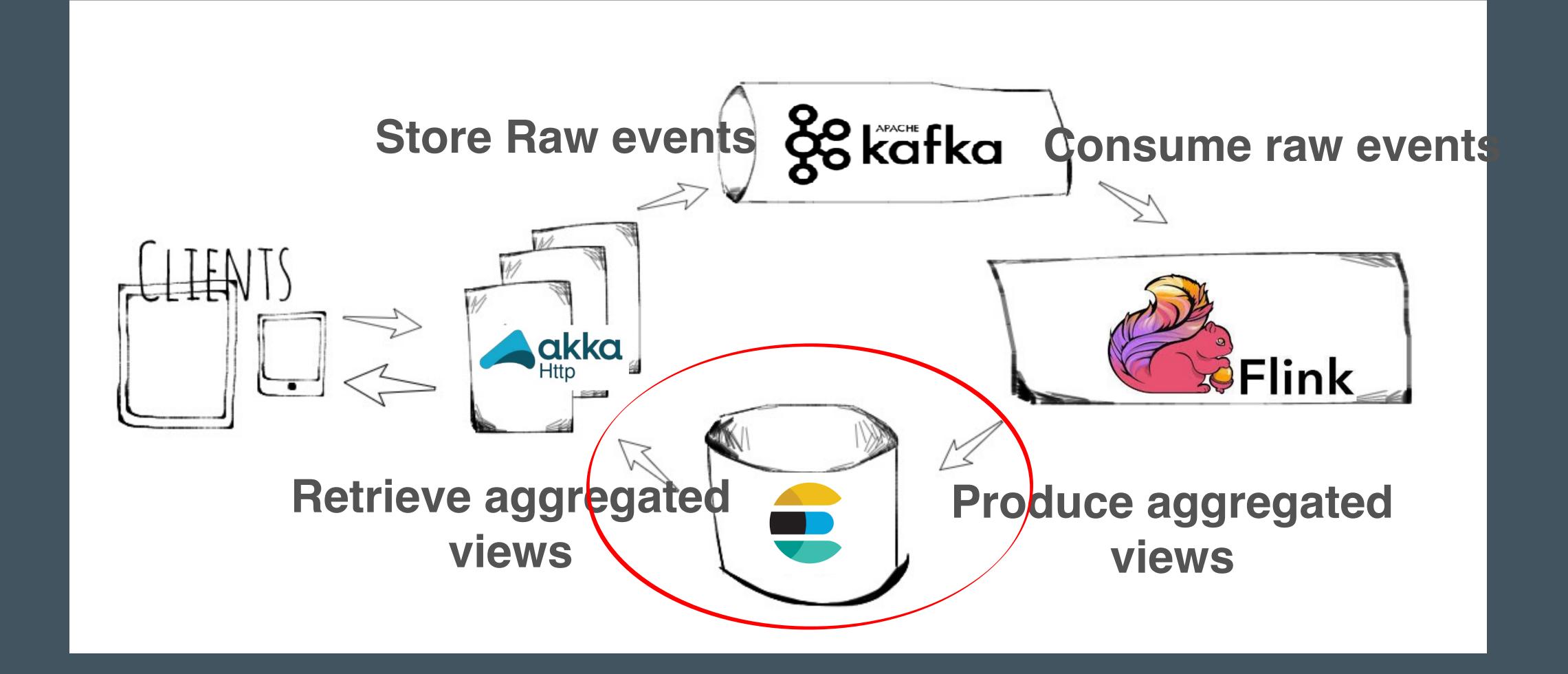








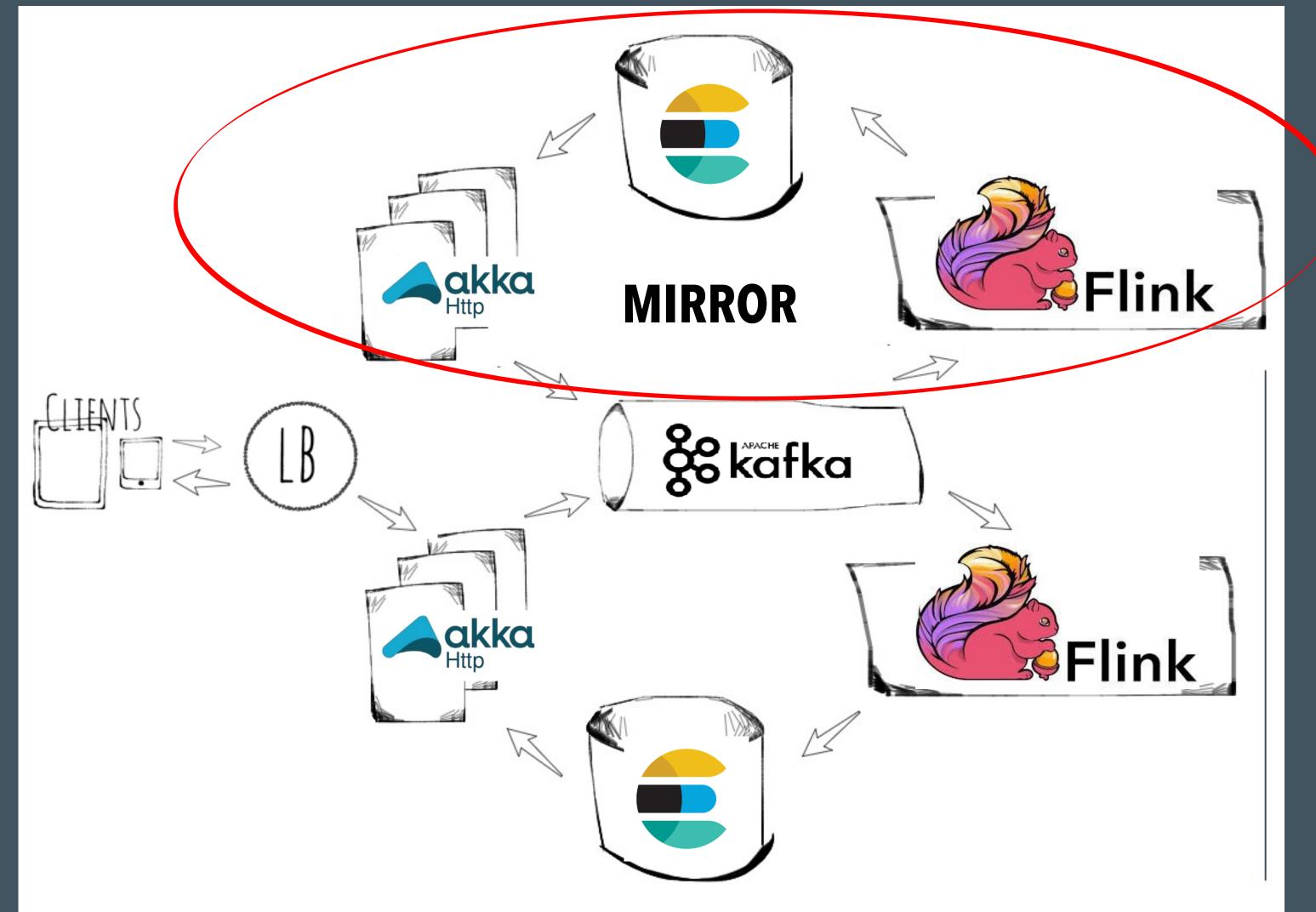








BLUE/GREEN APPROACH



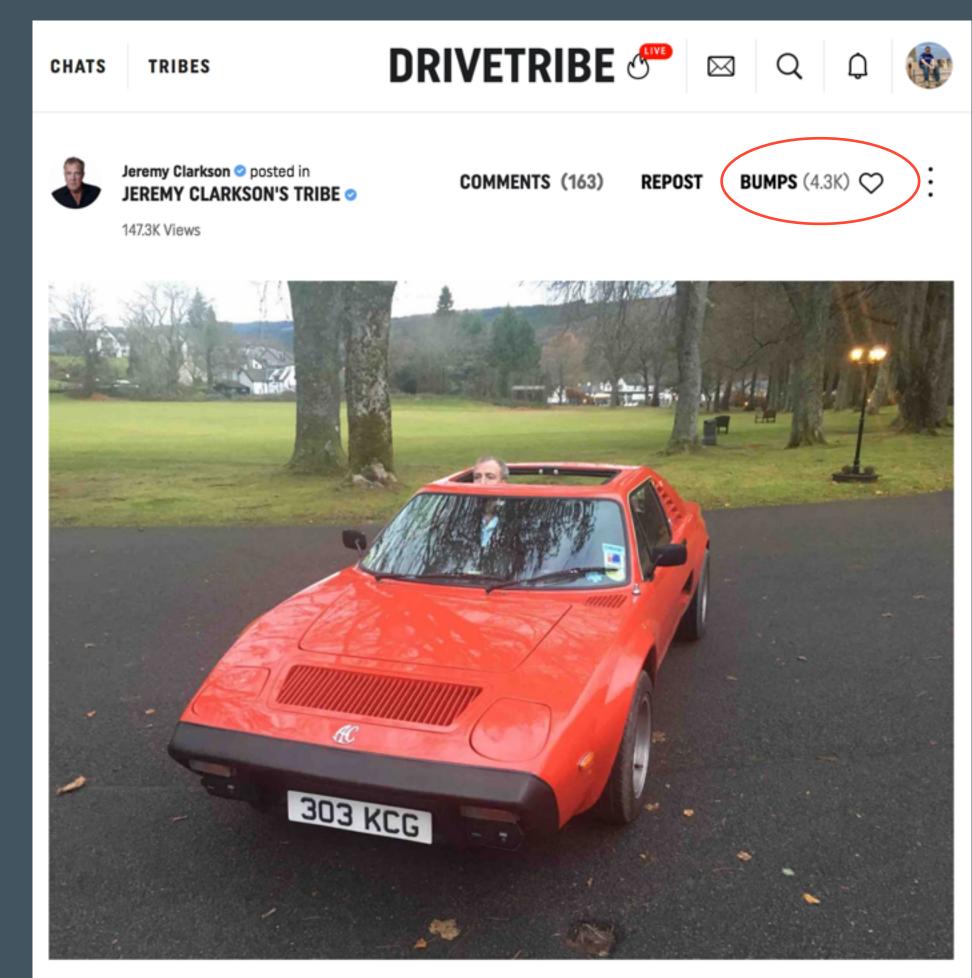


A REAL WORLD EXAMPLE

COUNTING BUMPS

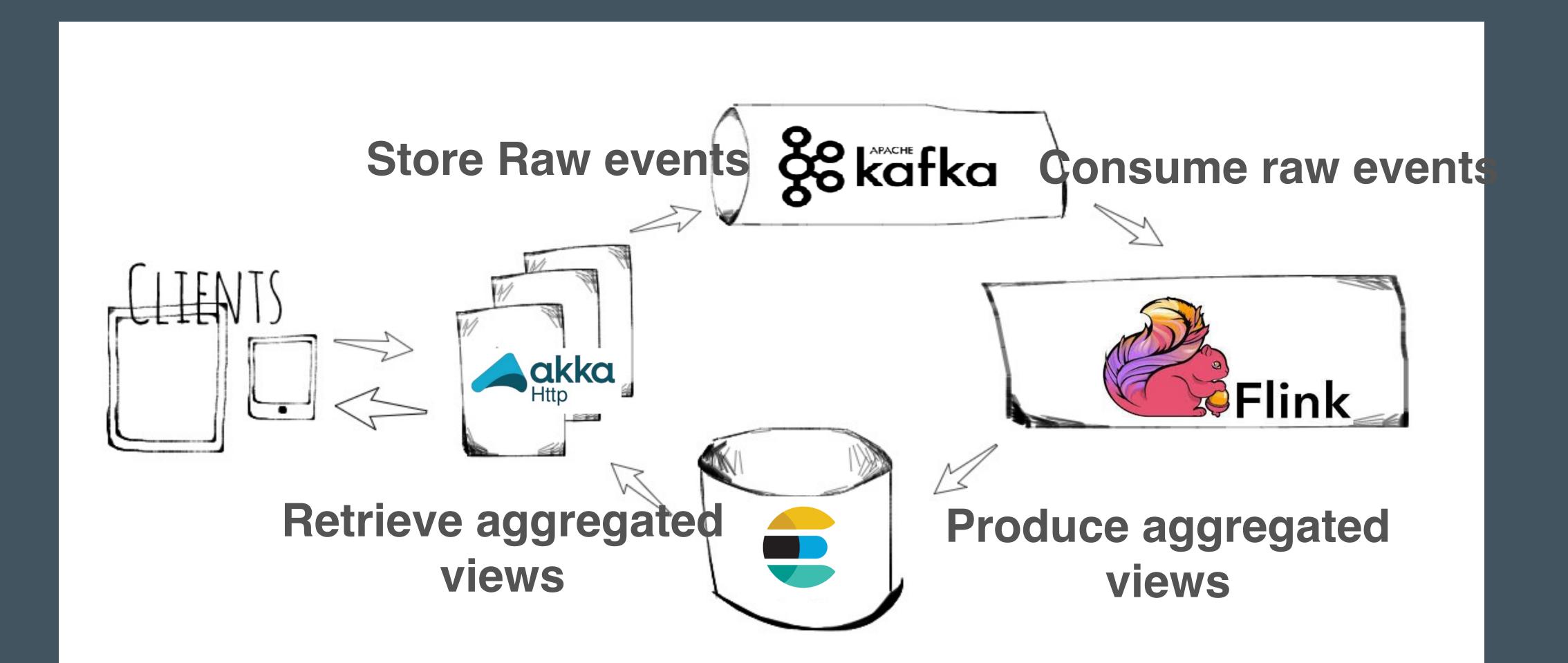
Thousands of people like the fact that Jeremy Clarkson is a really tall guy Users can "bump" a post if they like it

case class BumpEvent(id: Id[Bump], postId: Id[Post], userId: Id[User], bumpedOn: DateTime



Damn. Another dream dashed







COUNTING BUMPS

case class State(id: Id[Post], bumpCounter: ???)

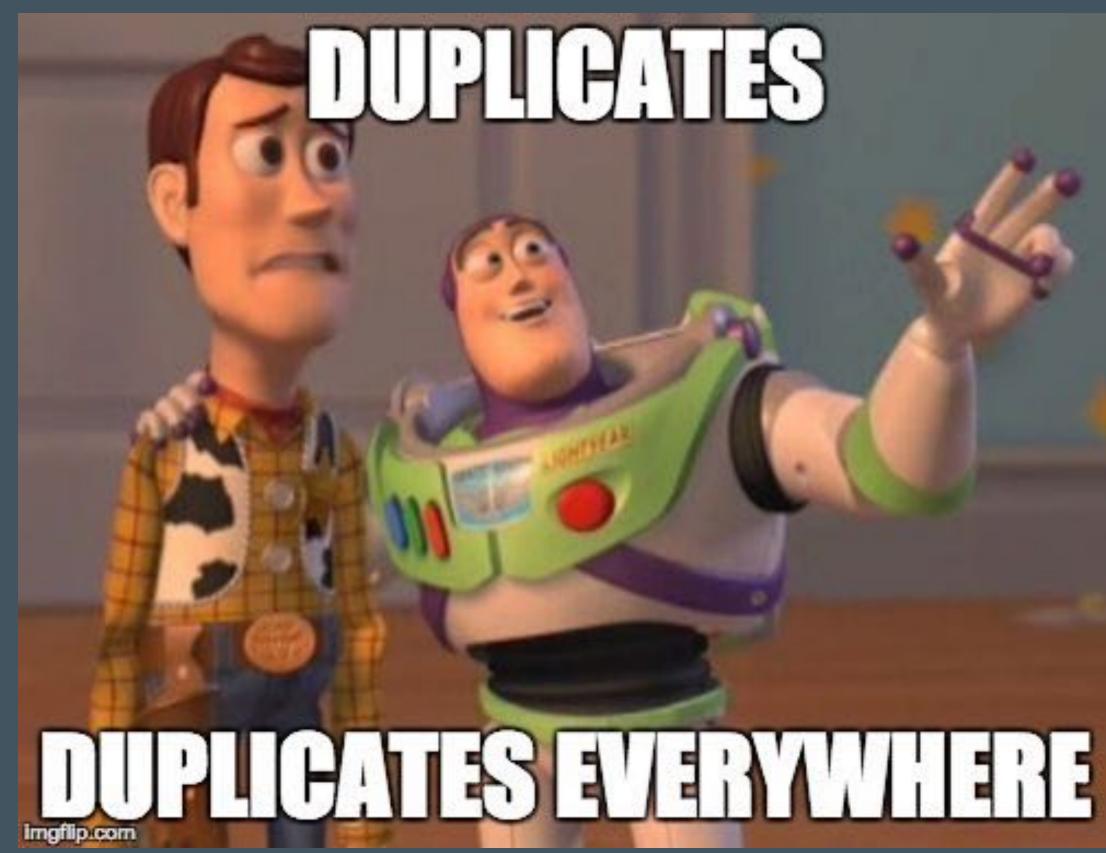
object PostStatsReducer { def apply(bumpEvents: DataStream[BumpEvent]): DataStream[State] = ??? }



```
case class State(id: Id[Post], bumpCounter: Long)
object State {
  def apply (bumpEvent: BumpEvent): State =
    State(bumpEvent.postId, 1L)
}
object PostStatsReducer {
  def apply(bumpEvents: DataStream[BumpEvent]): DataStream[State] =
    bumpEvents
      .map(State(_))
      .keyBy(_.id)
      . . .
      .combine
```

State(state1.id, state1.bumpCounter + state2.bumpCounter)

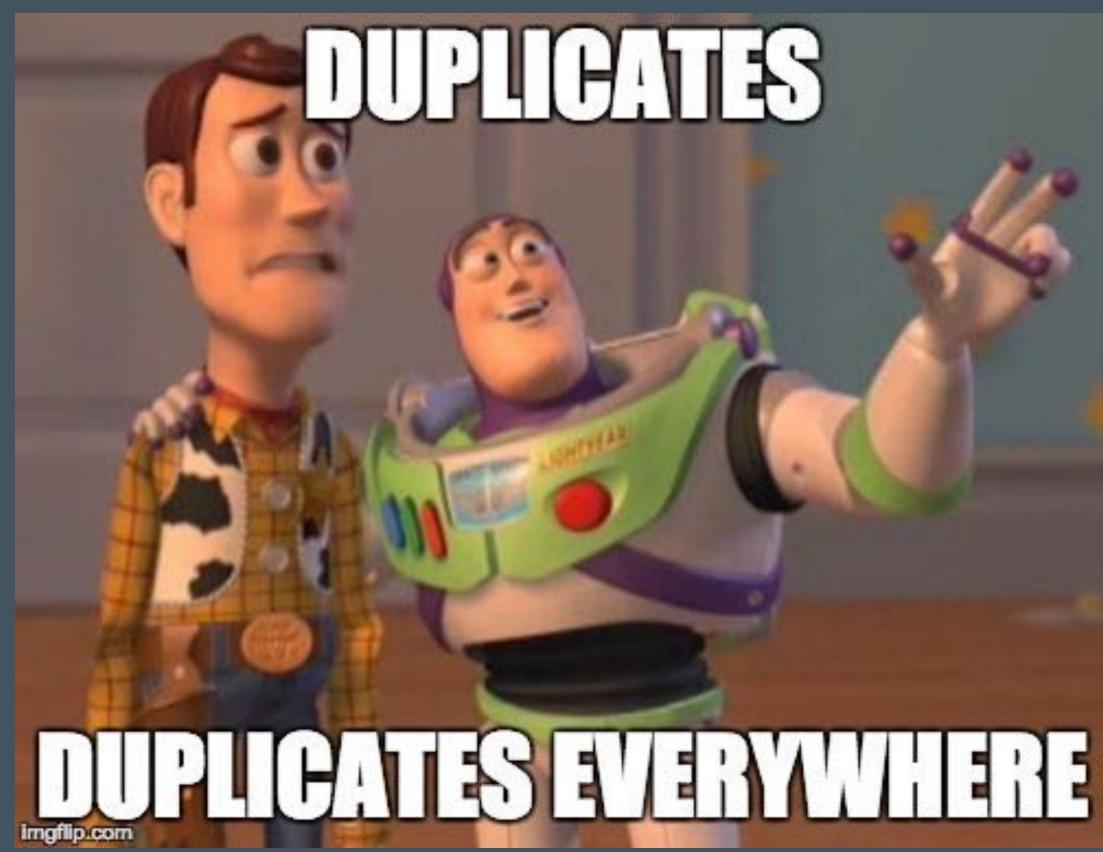








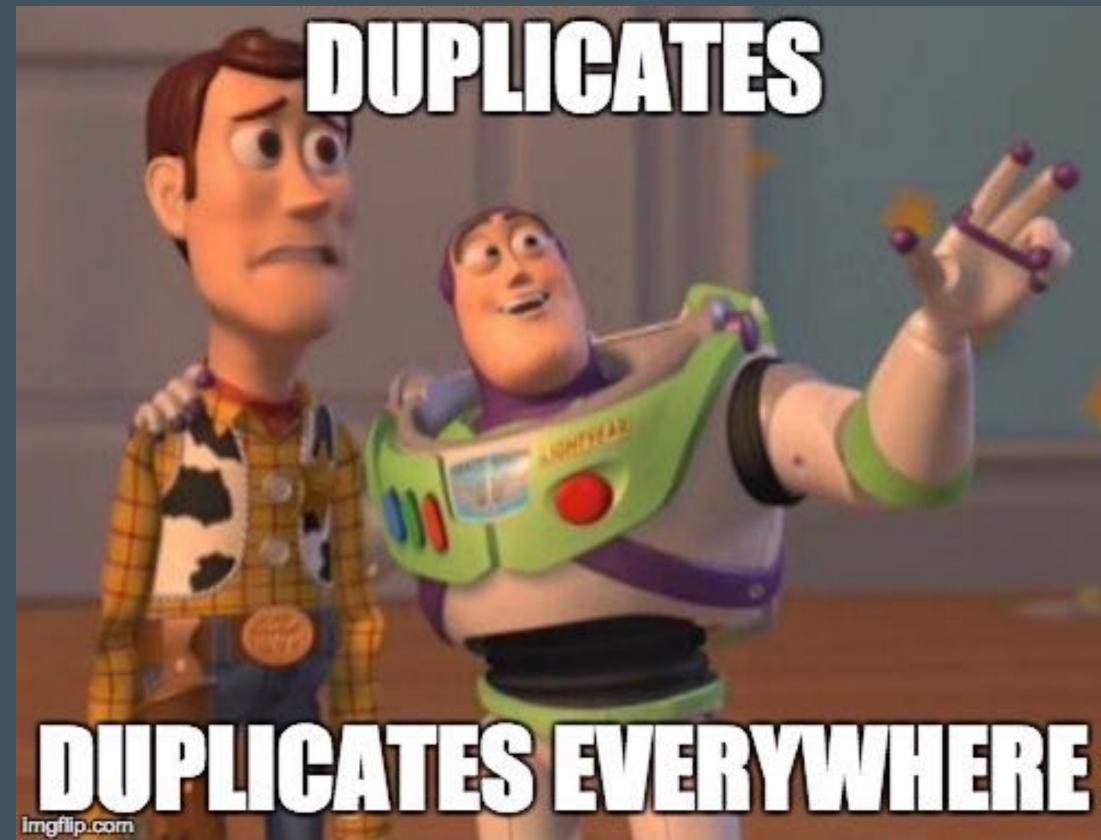
Use Flink with at least once semantics







- Use Flink with at least once semantics
- Our system is eventually consistent







WHAT DO WE KNOW ABOUT OUT COUNTER?

we know we're doing some kind of combine operation over States



WHAT DO WE KNOW ABOUT OUT COUNTER?

we know we're doing some kind of combine operation over States
we want our counter to be idempotent: a I+I a === a



WHAT DO WE KNOW ABOUT OUT COUNTER?

- we know we're doing some kind of combine operation over States
- \triangleright we want our counter to be idempotent: a |+| a == a
- we also want our counter to be associative: a + (b + c) = = (a + b) + c



BAND ALGEBRA

trait Band[T] { def combine(t1: T, t2: T): T }

- Closed
- Idempotent
- Associative



COUNTING BUMPS - SECOND ATTEMPT

```
case class State[T: Band](id: Id[Post], bumpCounter: T)
object State {
 def apply (bumpEvent: BumpEvent): State = ???
 implicit val stateBand: Band[State] = ???
object PostStatsReducer {
   bumpEvents
      .map(State(_))
      .keyBy(_.id)
      . . .
        band.combine(state1, state2)
      .combine
```

def apply(bumpEvents: DataStream[BumpEvent])(implicit band: Band[State]): DataStream[State] =



COUNTING BUMPS - SECOND ATTEMPT

```
implicit def setBand[U] = new Band[Set[U]] {
 def combine(s1: Set[U], s2: Set[U]): Set[U] =
   s1 ++ s2
case class State(id: Id[Post], bumps: Set[Bump])
object State {
  implicit val stateBand = new Band[State] {
   def combine(s1: State, s2: State): State =
      State(s1.id, s1.bumps |+| s2.bumps)
```

if all components of a case class have a band then so does the case class

would normally bring in Set implementation from a library

In normally that library would have a law testing module







COUNTING BUMPS - PUTTING IT TOGETHER

```
case class State(id: Id[Post], bumps: Set[BumpEvent])
object State {
 def apply (bumpEvent: BumpEvent): State =
   State(bumpEvent.postId, Set(bumpEvent))
 implicit val stateBand = new Band[State] {
   def combine(s1: State, s2: State): State =
     State(s1.id, s1.bumps |+| s2.bumps)
   }
object PostStatsReducer {
 def apply(bumpEvents: DataStream[BumpEvent])(implicit band: Band[State]): DataStream[State] =
   bumpEvents
      .map(State(_))
      .keyBy(_.id)
      band.combine(state1, state2)
      . . .
     .combine
```



OTHER ALGEBRAS WE USE

Adding events: Semigroup/Monoid Duplicate events: Band Out of order and duplicate events: Semilattice



