



# Java at Speed

# getting the most out of modern hardware

Gil Tene, CTO & co-Founder, Azul Systems



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# High level agenda

- Intro & Motivation
- Some hardware trends and new features
- Some compiler stuff
- A microbenchmark detour
- Some more compiler stuff
- Warmup, and what we can do about it
- Outting it all together (and maybe some bragging)

#### About me: Gil Tene

co-founder, CTO @Azul
 Systems

Have been working on "think different" GC and runtime approaches since 2002

A Long history building
 Virtual & Physical Machines,
 Operating Systems,
 Enterprise apps, etc...

At Azul we make JVMs that dramatically improve response time and latency behaviors

 I also depress people by demonstrating how terribly wrong their latency measurements are...



\* working on real-world trash compaction issues, circa 2004





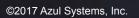
## What is it good for?







# Are you fast?



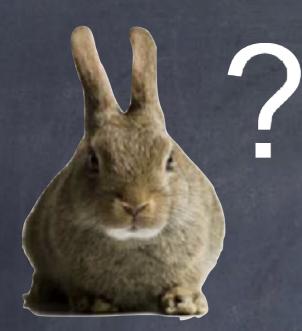


#### Are you fast when new code rolls out?





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## Are you fast when it matters?









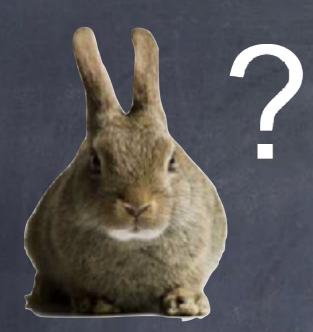
### Are you fast at Market Open?





### Are you fast when you actually trade?







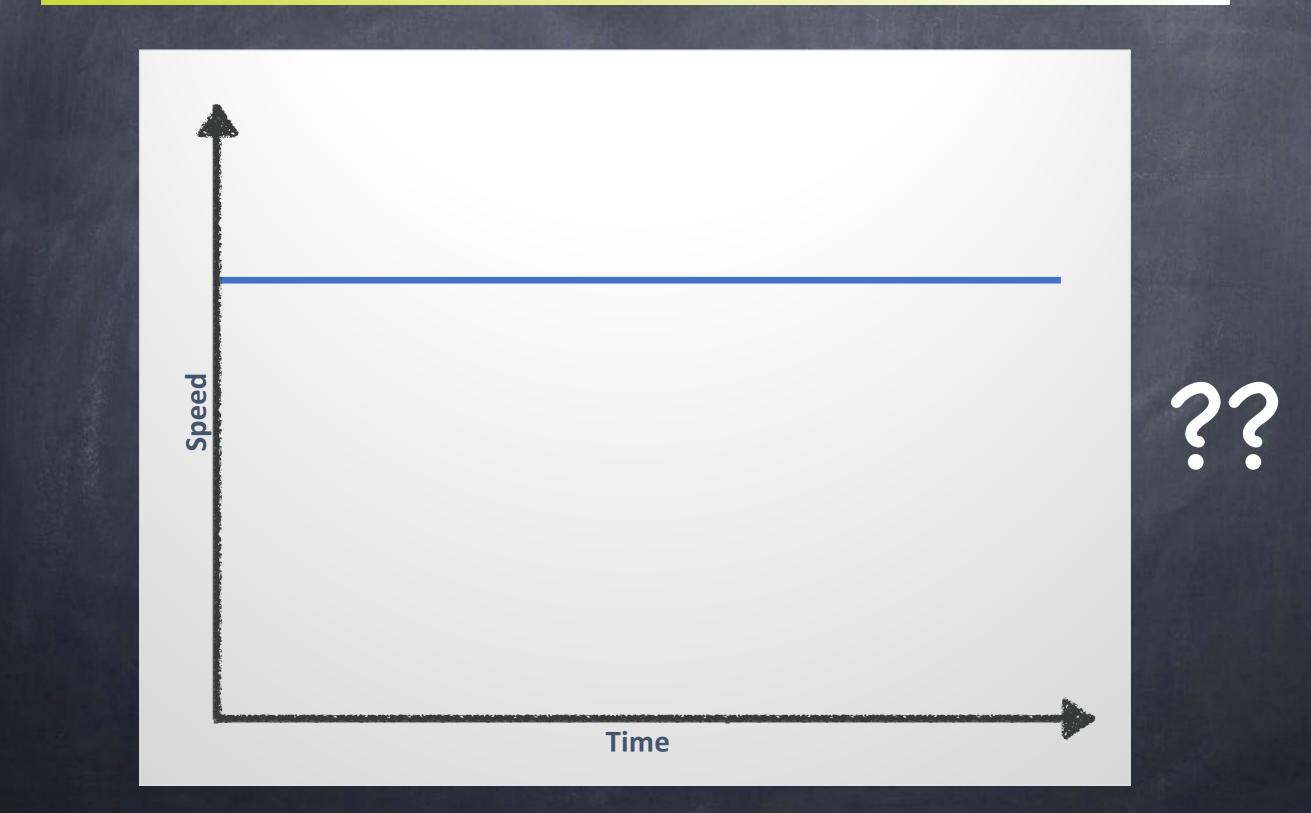
# Are you reliably fast?



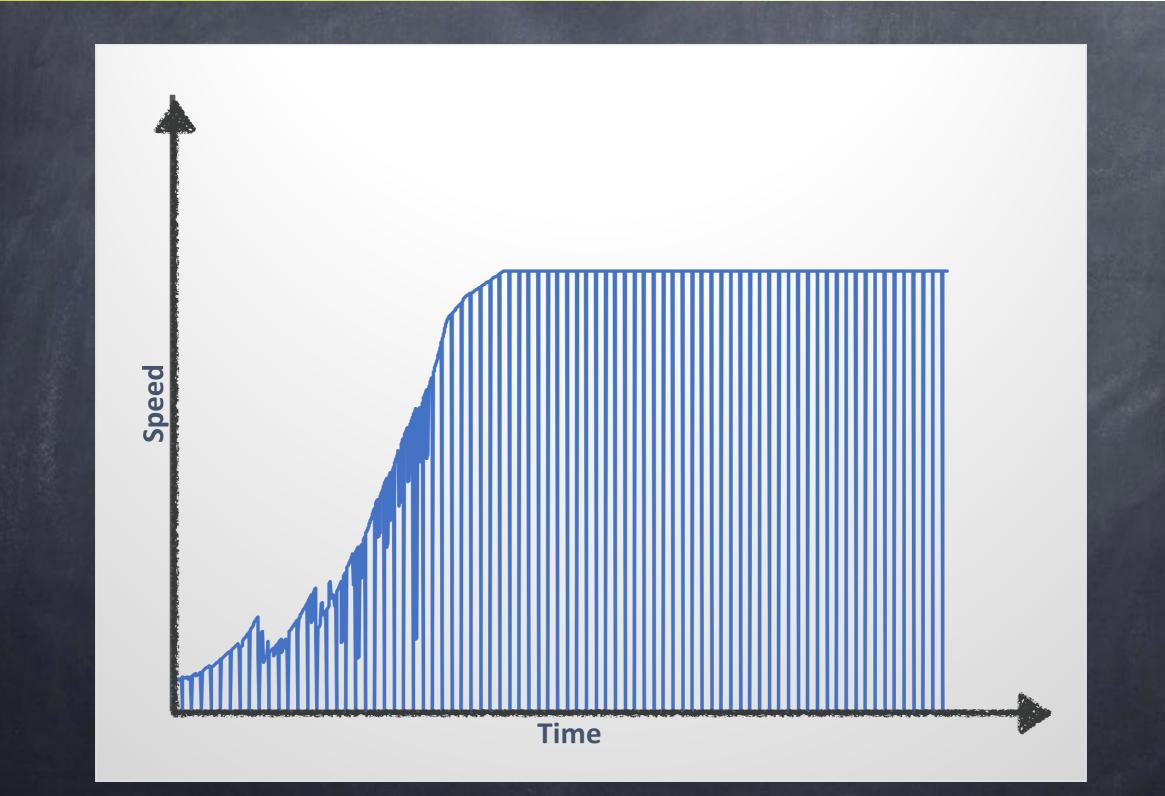


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#### What do you mean by "fast"?



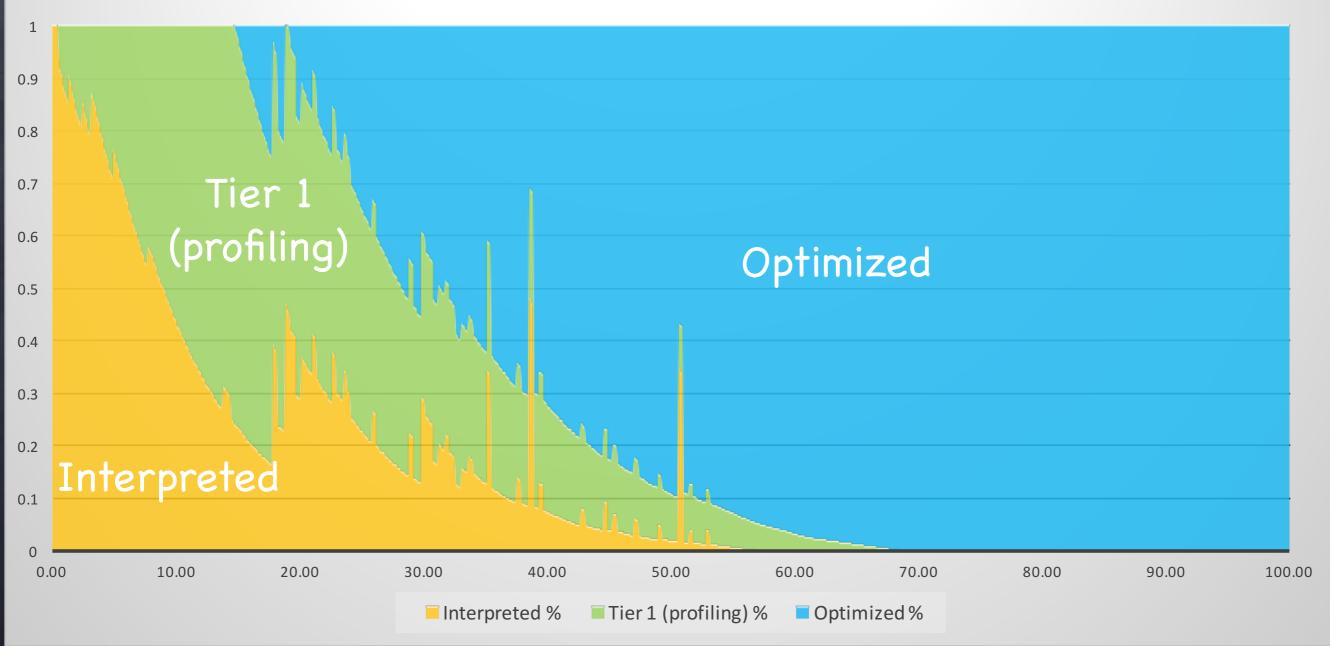
### What do you mean by "fast"?



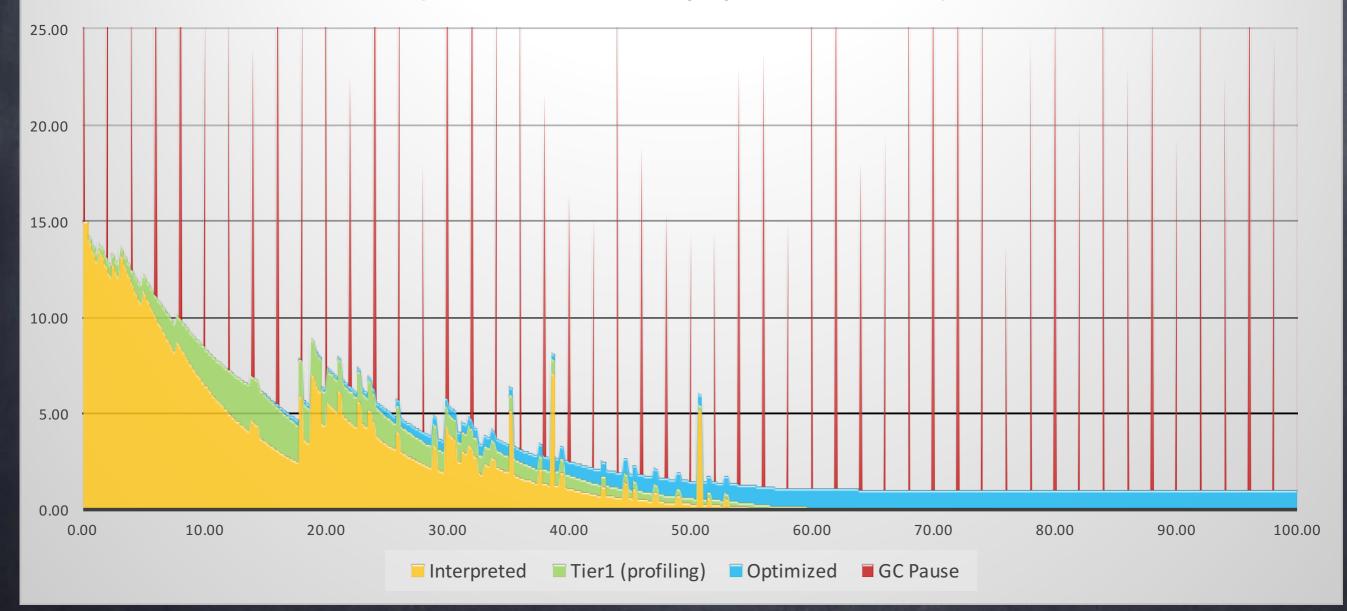
### Speed in the Java world...

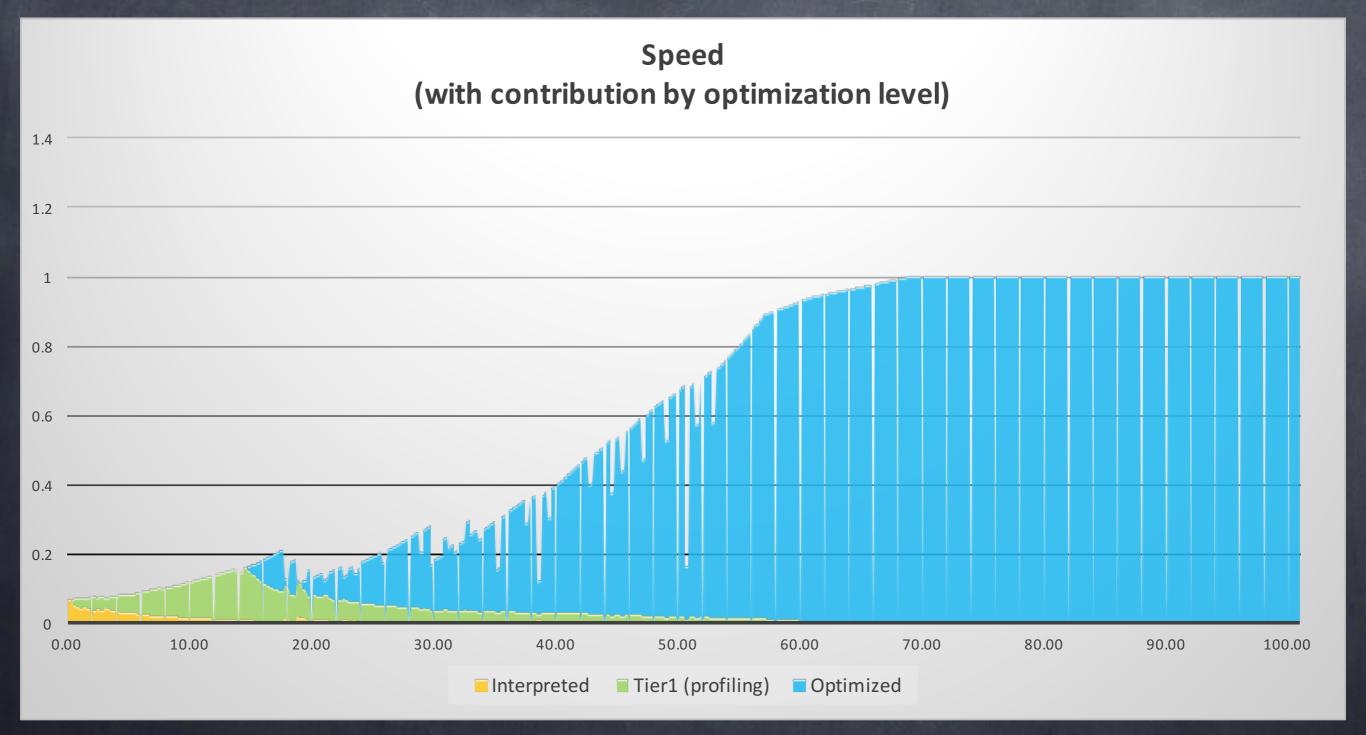






# Response time (with contribution by optimization level)





#### Some notes on modern servers



Code name	Model	Intro Date		cores/chip
Nehalem EP	Xeon 5500	March 2009		4
Westemere EP	Xeon 5600	June 2010		6
Sandy Bridge EP	E5-2600	March 2012	AVX	8
Ivy Bridge EP	E5-2600 V2	Sep. 2013		12
Haswell EP	E5-2600 V3	Sep. 2014	AVX2, BMI, BMI2	18
Broadwell EP	E5-2600 V4	March 2016	TSX, HLE	22
Skylake SP	Silver/Gold/	July 2017	AVX512	32



#### **Instruction Window Keeps Increasing**

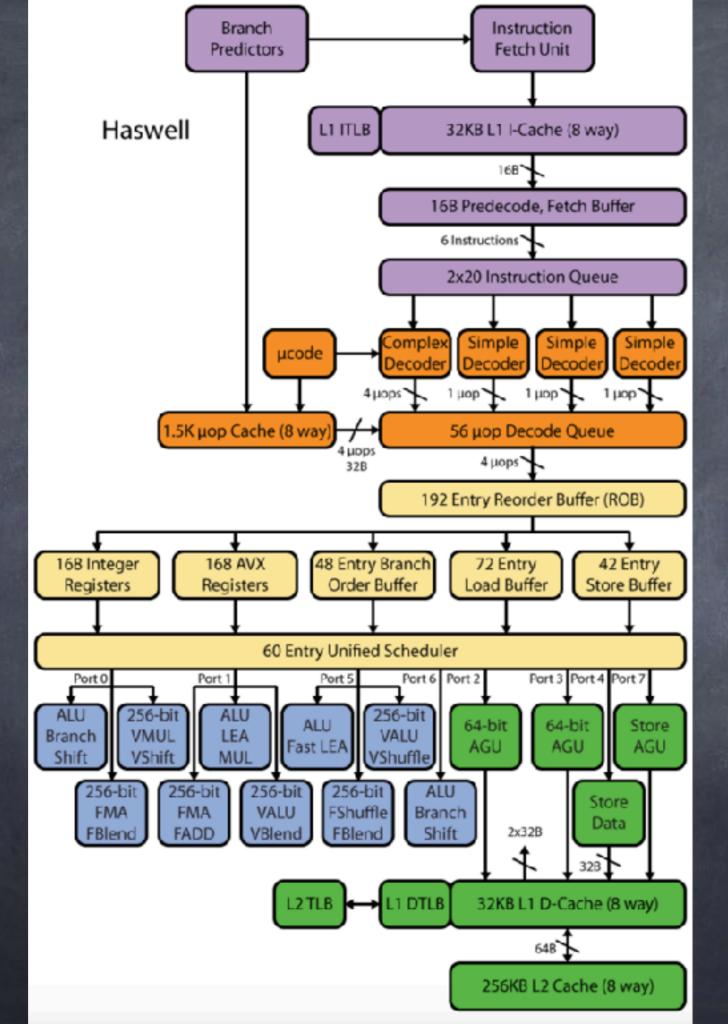
	Sandy Bridge	Haswell	SkyLake
Out-of-order Window	168	192	224 👚
In-flight Loads	64	72	72
In-flight Stores	36	42	56 👚
Scheduler Entries	54	60	97 👚
Integer Register File	160	168	180 👚
FP Register File	144	168	168
Allocation Queue	28/thread	56	64/thread 👚

#### Extract more parallelism in every generation

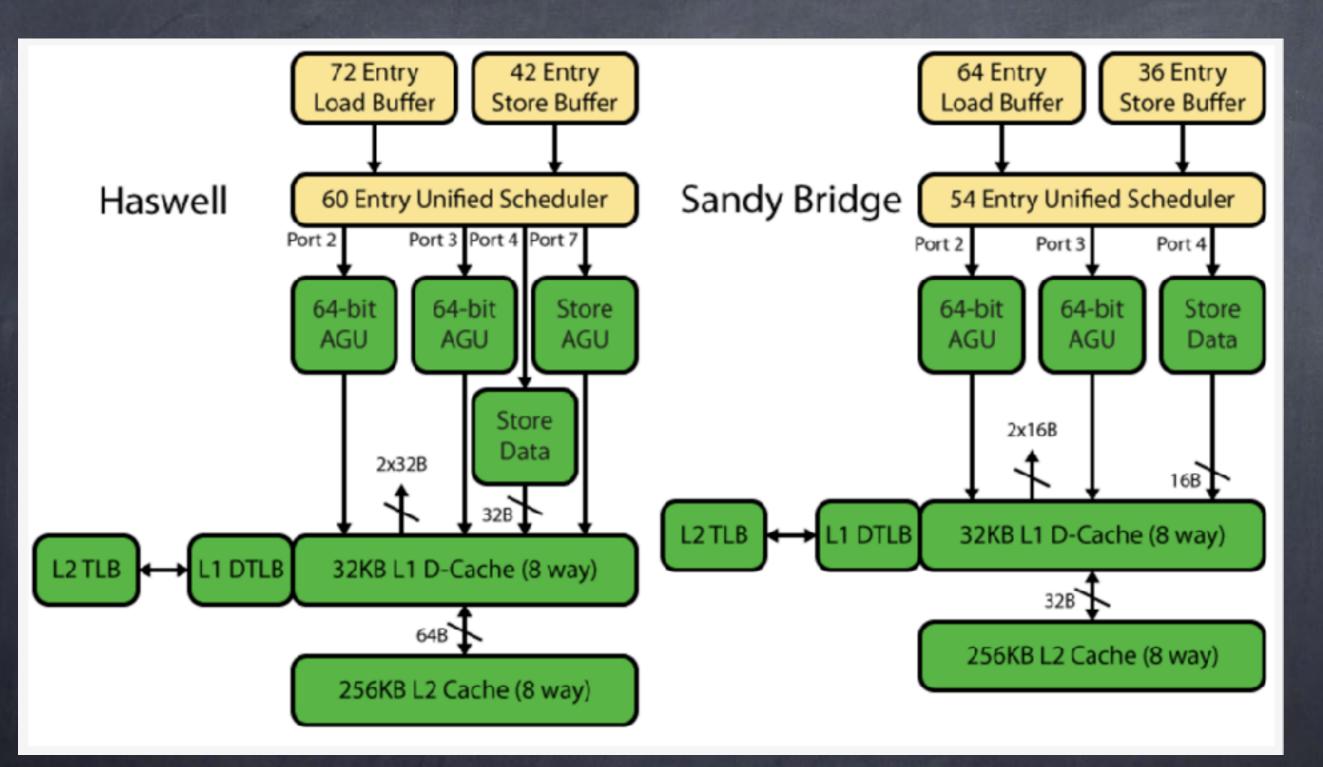


Intel Next Generation Microarchitecture Code Name Skylake

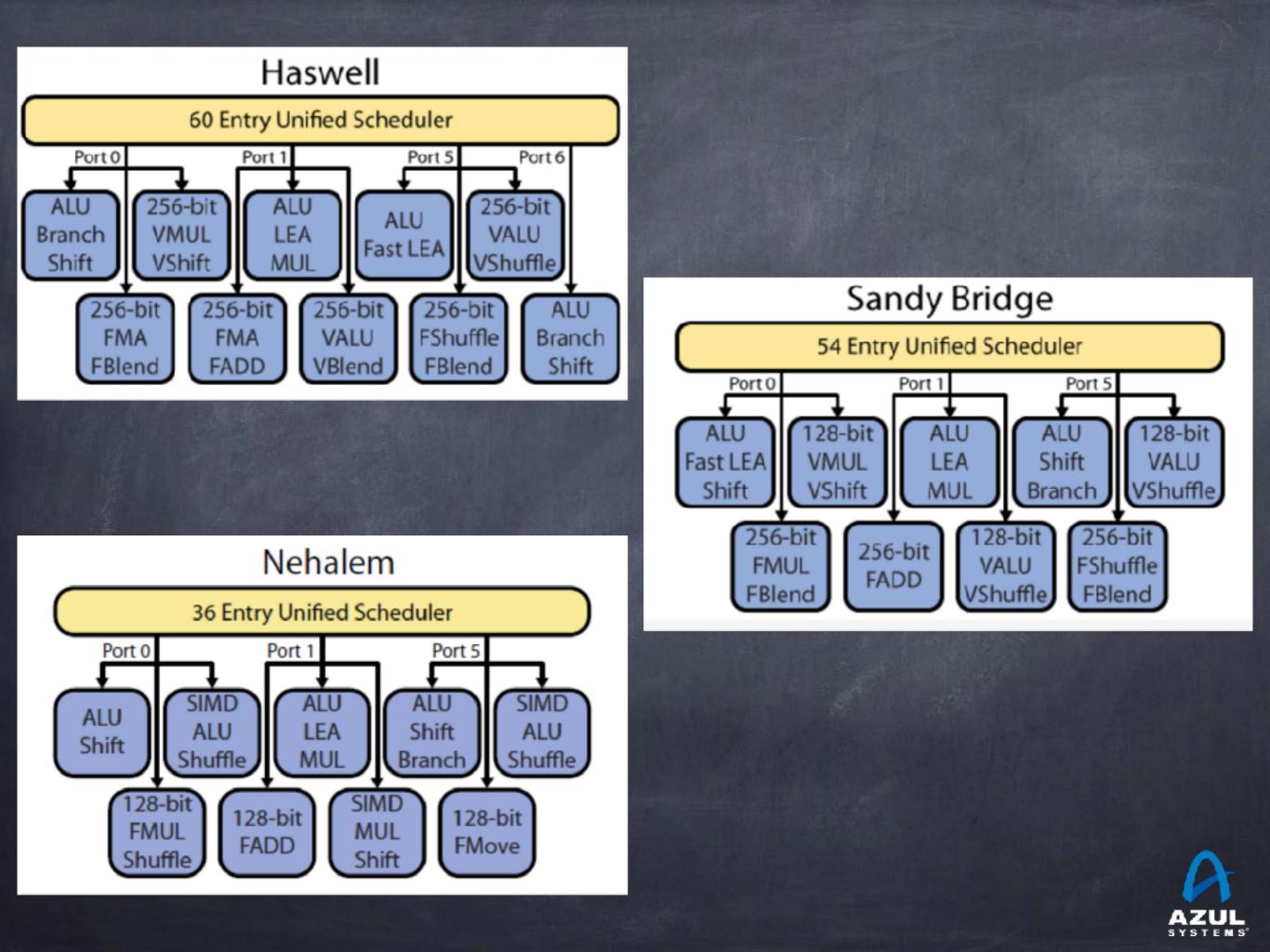












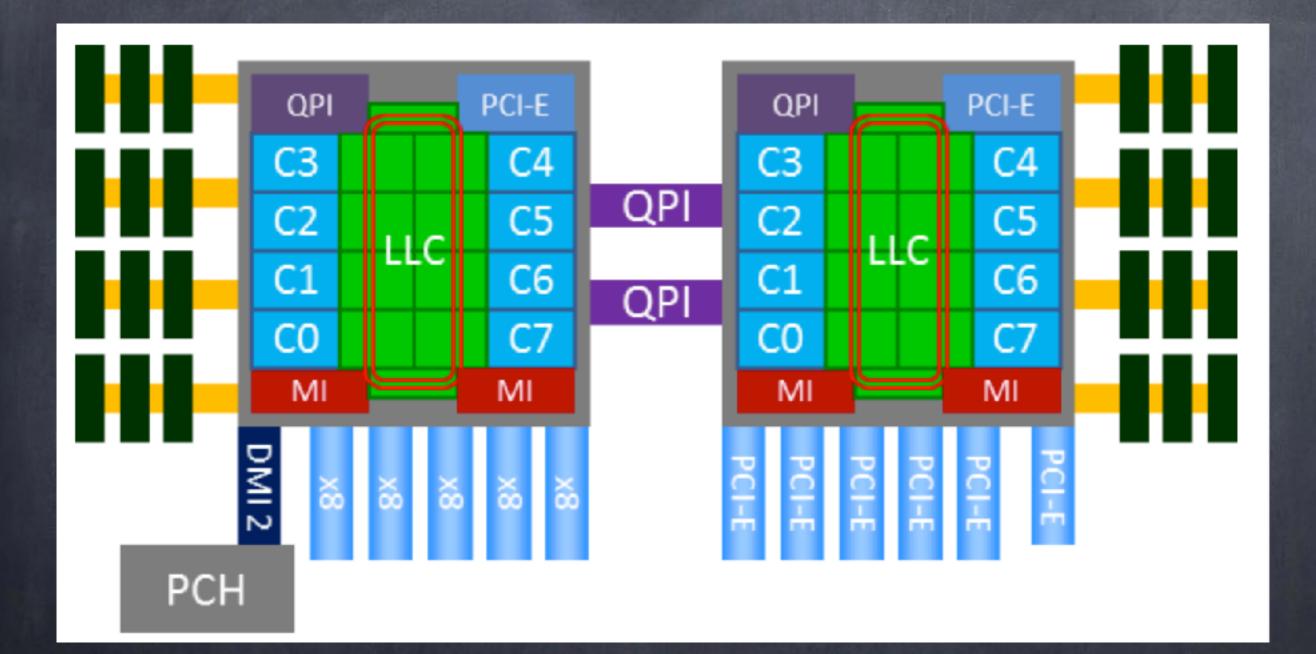
Caches

#### 

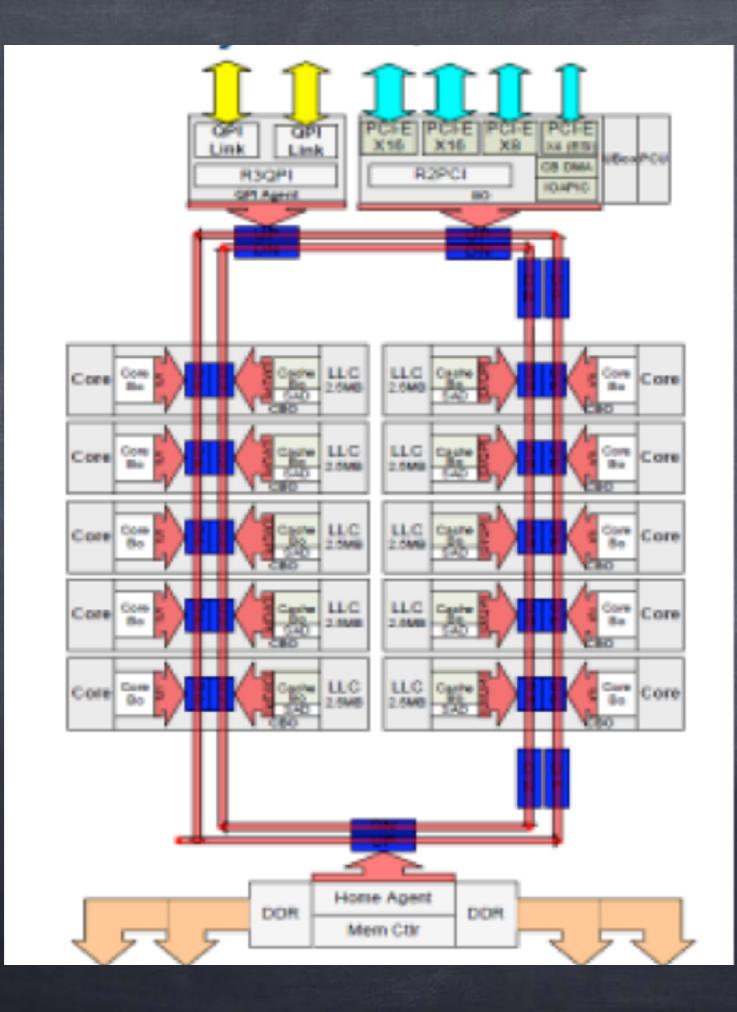
L3: ~2.5MB/core <u>shared</u> L3. (up to 55MB per socket)
TLBs: 2MB page support improved with E5-v3 & E5-v4
L2 DTLB used to be 4KB page only. Now 4KB/2MB
1024 entry 4KB/2MB L2 in E5-v3. 1536 entry in E5-v4



# System Topology: 2 sockets





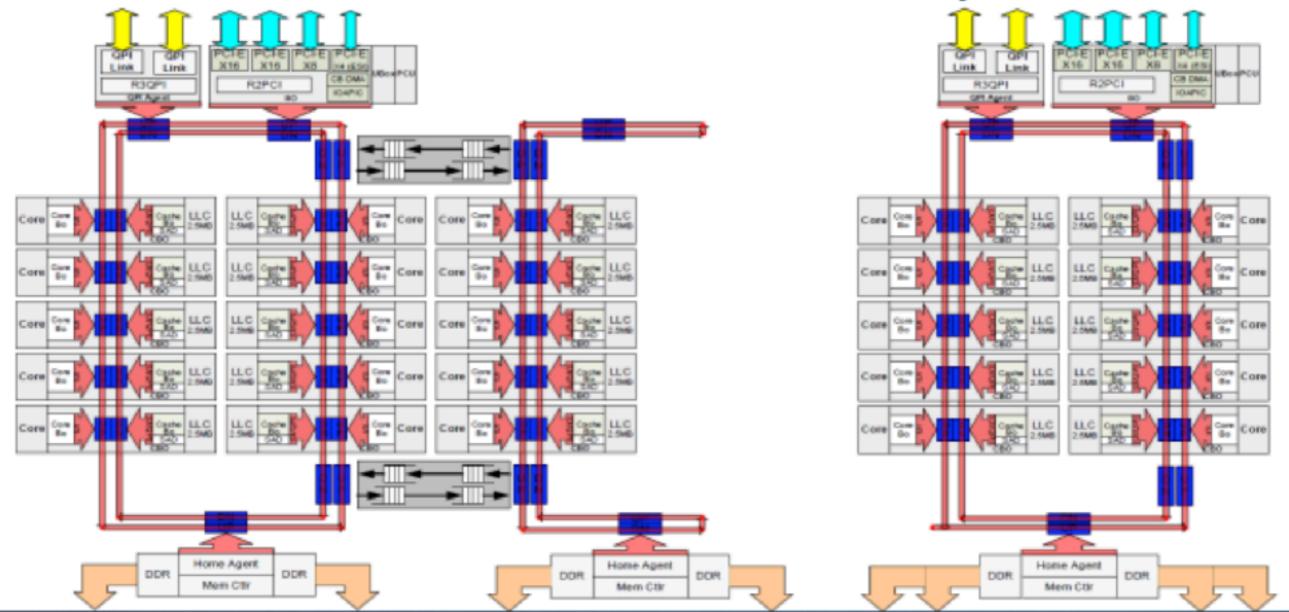




#### Intel<sup>®</sup> Xeon<sup>®</sup> Processor E5 v4 Product Family HCC Link Link 64. (BED) ROOP :## Bo Digen of Core 2.08448 2.0541 2.05/1 LLC LUC Bo Core LLC Core Core LLC Bo Core Bo Core Core BO Core B0 Core Core BO Core Core BO Core Core Bo Bo Be Core Core Home Agent Home Agent DDR DOR DOR DDR Mem Clir Mem Ctir

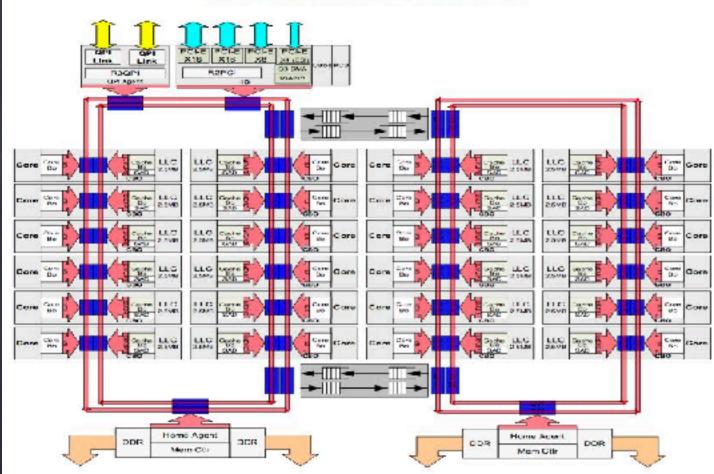


#### Intel® Xeon® Processor E5 v4 Product Family MCC/LCC





#### Broadwell EX 24-core die



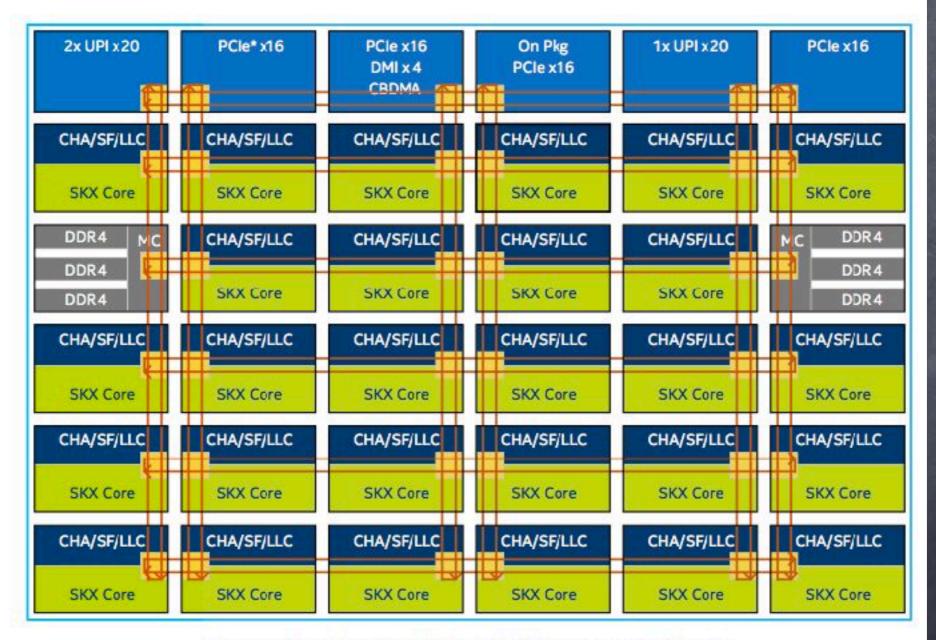
#### Skylake-SP 28-core die

2xUPLx20	PCle*x16	PCIe x16 DMI x 4 CBDMA	On Pkg PCIex16	1x UPI x 20	PClex16
CHA/SF/LLC	снаубрыс	CHA/SF/LLC	CHAYSEALLC	CHW/SF/LLC	CHA/SF/LLC
SKX Core	SKX Core	SKX Core	SKX Core	SKX Core	SiQL Core
DOR4 NC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	VC DDR4
DDR4	SKX Core	SKX Core	SKX Core	SKX Core	DDR4
CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC
SKX Core	SKX Core	SKX Core	SKX Core	SKX Core	SKX Core
CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC
SIX Core	SKX Core	SKX Core	SKX Core	SKX Core	SIX Core
CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC	CHA/SF/LLC
SKX Core	SKX Core	SKX Core	SKX Core	SKX Core	SKX Core

CHA – Caching and Home Agent ; SF – Snoop Filter; LLC – Last Level Cache; SKX Core – Skylake Server Core; UPI – Intel® UltraPath Interconnect

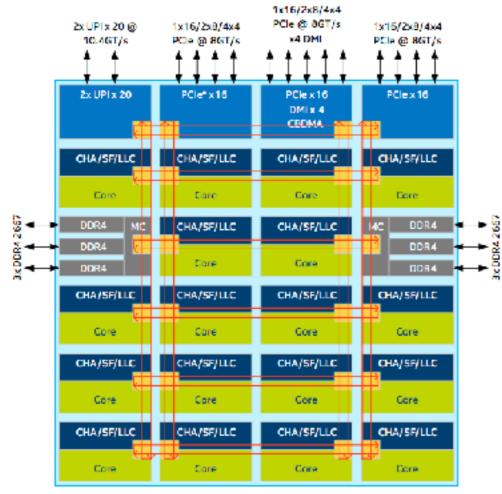


#### Skylake-SP 28-core die

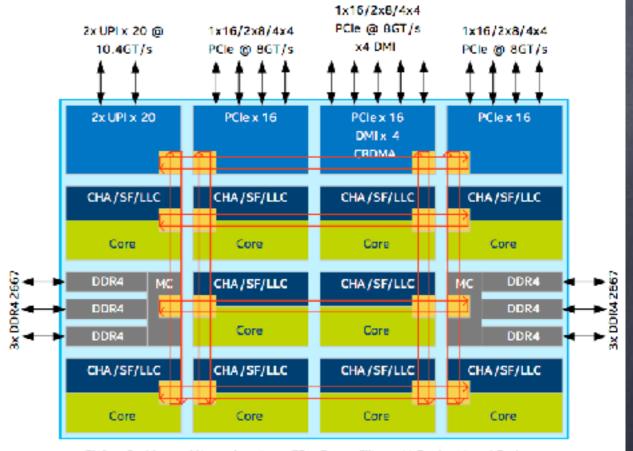


CHA – Caching and Home Agent ; SF – Snoop Filter; LLC – Last Level Cache; SKX Core – Skylake Server Core; UPI – Intel\* UltraPath Interconnect





CHA - Caching and Home Agent ; SF- Shoop Filter ; LLC- Last Level Cache ; Core - Skylake-SP Core; UPI - Intel® UltraPath Interconnect LCC (up to 10 Cores)



CHA – Caching and Home Agent ; SF – Snoop Filter ; LLC – Last Level Cache ; Core – Skylake -SP Core ; UPI – Intel® UltraPath Interconnect



HCC (up to 18 cores)

#### Some machine code zoom-in



AZUL	Azul ZVIsion	User: gil Host: cperf2/47781	Version: 1.8.0-zing_99.99.99.99.dev-b1686 Uptime: 00:21:51
Overview   Azul Suppo	ort I Threads CPU Memory I Compilers	Applications	
Tick Profile I Metatick	s Event Tracker		
Pause The Tick Colle	ection Reset Tick Profile Stop Saving Ticks T	o Disk	
Timer Tick Pro	filo		

#### TIMEL TICK FROME

Cutoff:	1

Threads (comma separated list):

o<sup>None</sup> ⊖<sup>JVM</sup> ⊖<sup>All</sup> Submit

SYSTEMS

Note: Functions that could not be resolved to a name string are displayed as an address followed by "<-" and the first calling function (in their stack) that can.

Percent	Ticks	Source
		bench.CodeGenExampleBench.truelfMaskMatched (codeblob)
18.9%	5,169	bench.CodeGenExampleBench.addArraysIfEven (codeblob)
11.2%	3,063	bench.CodeGenExampleBench.addXtoArray (codeblob)
8.3%	2,267	bench.CodeGenExampleBench.sumShifted (codeblob)
3.0%	2,190	bench.CodeGenExampleBench.sumIfEvenLoop (codeblob)
5.1%	1,405	bench.CodeGenExampleBench.sumLoop (codeblob)
0.5%	128	ElfSymbolTable::lookup(unsigned char*, int*, int*, int*, unsigned long*) (PC ref)
0.3%	98	_dl_addr (PC ref)
0.2%	64	_int_free (PC ref)
0.2%	59	libLLVM-4.0svn.so (PC ref)
0.2%	58	_int_mailoc (PC ref)
0.2%	57	libc_malloc (PC ref)
0.1%	35	syscall (PC ref)
0.1%	32	llvm::PMDataManager::findAnalysisPass(void const*, bool) (PC ref)
0.1%	16	computeKnownBits(Ilvm::Value const*, Ilvm::APInt&, Ilvm::APInt&, unsigned int, (anonymous namespace)::Query const&) (PC ref)
0.1%	15	llvm::PMTopLevelManager::setLastUser(llvm::ArrayRef, llvm::Pass*) (PC ref)
0.0%	12	bench.CodeGenExampleBench.doAddArraysIfEven (codeblob)
0.0%	12	ioctl (PC ref)
0.0%	12	malloc_consolidate (PC ref)
0.0%	12	Ilvm::InstCombiner::run() (PC ref)
0.0%	11	Ivm::PMTopLevelManager::findAnalysisPass(void const*) (PC ref)
0.0%	11	Ivm::Use::getImpliedUser() const (PC ref)
0.0%	10	Ivm::ValueHandleBase::AddToUseList() (PC ref)
0.0%	10	mprotect (PC ref)
0.0%		malloc (PC ref)
0.0%	9	sysmalloc (PC ref)
0.0%	8	bench.CodeGenExampleBench.doSumShiftedLoop (codeblob)
0.0%	8	llvm:Value:getValueName() const (PC ref)

						Uptime: 00	:29:18
verview	Azul Su	pport Threads   CPU   Memory   Com	pilers I Applications				
		Web server I Perf data I Code cache		PC   Code profile	Stub code	e I <u>Old stats</u>	1
BA stats	Polli	ng Opportunities   Flush memory   M	lonitor   java.lang.Object				
					_	Reset Tick Pr	ofile
ench.Co	deGenE	xampleBench.sumLoop([i)I	Dx30	000f600		Heset Hok Ph	onne
Assembly	I <u>Call</u>	ee Caller					
Percent	Ticks A	ddress Code	Opcode	1			
		x30001650 pushq %rax	0111x0				
		x30001652 cmpl S0, %gs:104	0x65833c256800000000				
		3000f65b jne 127 ; ABS: 0x3000f8dc	0x757f				
		x3000f65d movi 8(%rsi), %ecx // NPE->	0x8b4e08				
		0x30001615					
	0:	x30001660 testq %rcx, %rcx	0x4885c9				
	0:	x30001663 je 115 ; ABS: 0x3000f6d8	0x7473				
	0:	x3000f665 cmpl \$7, %ecx	0x83f907				
	0;	x3000f668 ja 20 ; ABS: 0x3000f67e	0x7714				
	0:	x3000166a xorl %edx, %edx	0x31d2				
	0:	x3000166c xorl %eax, %eax	0x31c0				
	0:	k3000f66e nop	0x6690				
	0:	x3000±670 addl 12(%rsi,%rdx,4), %eax	0x0344960c				
	0:	x30001674 incq %rdx	0x48ffc2				
	0:	x30001677 cmpq %rcx, %rdx	0x4839ca				
	0:	x3000±67a jl -12 ; ABS: 0x3000/670	0x7cf4				
	0:	k3000±67e popq %rcx	0x59				
0.05%	1 0:	x3000167d retq	0xc3				
	0:	x3000167e movi %ecx, %r8d	0x4189c8				
	0:	x3000±691 and \$7, %r8d	0x4183c007				
	0:	x3000±685 movq %rcx, %rdx	0x4889ca				
	0:	x30001688 subq %r8, %rdx	0x4c29c2				
0.05%	1 0:	x3000168b je -35 ; ABS: 0x3000f66a	0x74dd				
	0:	x3000±68d leag 28(%rsi), %rax	0x488d461c				
	0:	x3000±691 pxor %xmm0, %xmm0	0x660fefc0				
	0:	x30001695 movq %rdx, %rdi	0x4889d7				
	0:	x30001698 pxor %xmm1, %xmm1	0x660fefc9				
	0:	x3000±69e nopl (%rax)	0x0f1f4000				
21.90%	418 0:	x3000±6a0 movdqu -16(%rax), %xmm2	0xf30f6f50f0				
5.08%	97 0:	x3000f6a5 movdqu (%rax), %xmm3	0xf30f6f18				
40.86%	780 0:	x300016a9 paddd %xmm2, %xmm0	0x660ffec2				
2.72%		x300016ad paddd %xmm3, %xmm1	0x660ffecb				
29.23%	558 0:	к3000f6b1 addq \$32, %rax	0x4883c020				
		x3000f6b5 addq S-8, %rdi	0x4883c7f8				
	0:	x300016b9 jne -27 ; ABS: 0x3000f6a0	0x75e5				
		x300016bb paddd %xmmD, %xmm1	0x660ffec8				
		x3000f6bf pshuld \$78, %xmm1, %xmm0	0x660f70c14e				
		x3000f6c4 paddd %xmm1, %xmm0	0x660ffec1				
		x300016c8 phaddd %xmm0, %xmm0	0x660f3802c0				
0.05%		x300016cd movd %xmm0, %eax	0x660f7ec0				
0.05%		x3000f6d1 testl %r8d, %r8d	0x4585c0				
		x3000f6d4 jne -102 ; ABS: 0x30001870	0x759a				
		x30001646 jmp -92 ; ABS: 0x3000f67c	0xeba4				
		x300016d8 xorl %eax, %eax	0x31c0				
		k3000f6da popq %rcx	0x59	1			
		x3000f6db retg	0xc3				

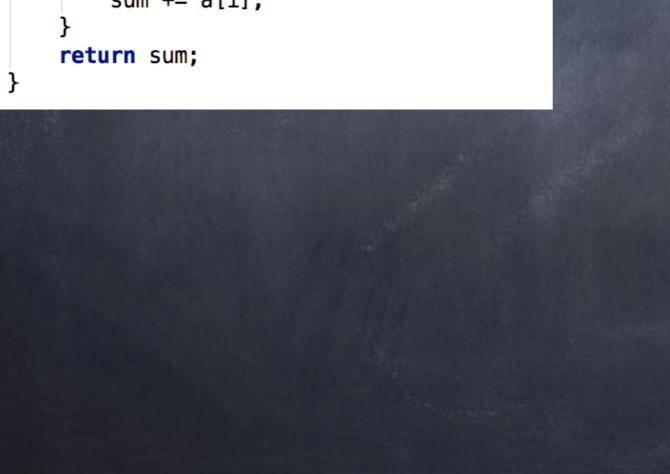


#### A simple array summing loop

```
private int sumLoop(int[] a) {
    int sum = 0;
    for (int i = 0; i < a.length; i++) {
        sum += a[i];
    }
    return sum;
}</pre>
```



```
private int sumLoop(int[] a) {
    int sum = 0;
    for (int i = 0; i < a.length; i++) {
        sum += a[i];
    }
    return sum;</pre>
```



Address	Code	Opcode
0x3000fe50	pushq %rax	0xfff0
0x3000fe52	empl \$0, %gs:104	0x65833c25680000000
0x3000fe5b	ne 127 ; ABS: 0x3000fedc	0x757f
	movl 8(%rsi), %ecx // NPE-> 0x3000fef5	0x8b4e08
0x3000fe60	testq %rcx, %rcx	0x4885c9
0x3000fe63	je 115 ; ABS: 0x3000fed8	0x7473
0x3000fe65	cmpl \$7, %ecx	0x83f907
0x3000fe68	ja 20 ; ABS: 0x3000fe7e	0x7714
0x3000fe6a	xorl %edx, %edx	0x31d2
0x3000fe6c	xorl %eax, %eax	0x31c0
0x3000fe6e		0x6690
	addl 12(%rsi,%rdx,4), %eax	0x0344960c
0x3000fe74		0x48ffc2
	empq %rcx, %rdx	0x4839ca
	J -12 ; ABS: 0x3000fe70	0x7cf4
0x3000fe7c		0x59
0x3000fe7d		0xc3
	movl %ecx, %r8d	0x4189c8
	andi \$7, %r8d	0x4183e007
	movq %rcx, %rdx	0x4889ca
0x3000fe88	subq %r8, %rdx	0x4c29c2
	je -35 ; ABS: 0x3000fe6a	0x74dd
0x3000fe8d	leaq 28(%rsi), %rax	0x488d461c
0x3000fe91	pxor %xmm0, %xmm0	0x660fefc0
0x3000fe95	movq %rdx, %rdi	0x4889d7
0x3000fe98	pxor %xmm1, %xmm1	0x660fefc9
0x3000fe9c		0x0f1f4000
	movdqu -16(%rax), %xmm2	0xf30f6f50f0
	movdqu (%rax), %xmm3	0xf30f6f18
	paddd %xmm2, %xmm0	0x660ffec2
	paddd %xmm3, %xmm1	0x660ffecb
	addq \$32, %rax	0x4883c020
	addq \$-8, %rdi	0x4883c7f8
	jne -27 ; ABS: 0x3000fea0	0x75e5
	paddd %xmm0, %xmm1	0x660ffec8
0x3000febf	pshufd \$78, %xmm1, %xmm0	0x660f70c14e
	paddd %xmm1, %xmm0	0x660ffec1
	phaddd %xmm0, %xmm0	0x660f3802c0
	movd %xmm0, %eax	0x660f7ec0
	testi %r8d, %r8d	0x4585c0
	jne -102 ; ABS: 0x3000fe70	0x759a
	mp -92 ; ABS: 0x3000fe7c	0xeba4
	10rl %eax, %eax	0x31c0
0x3000feda	popq %rcx	0x59
0x3000fedb		0xc3
0x3000fedc	movq %rsi, (%rsp)	0x48893424
	movabsq \$805334400, %rax	0x48b8806d00300000000
0x3000feea		0biix0
	movq (%rsp), %rsi	0x488b3424
	jmp -152 ; ABS: 0x3000fe5d	0xe968ffffff
	movabsq \$805319872, %rax	0x48b8c03400300000000
	movl \$7, %edi	0xbf07000000
0x3000ff04		0xffd0
	addq \$-8, %rsp	0x4883c4f8
0x3000ff0a	mp -50575 ; ABS: 0x30003980 = StubRoutines::deoptimize	0xe9713affff
0x3000ff0f	ht3	0xcc



```
private int sumLoop(int[] a) {
    int sum = 0;
    for (int i = 0; i < a.length; i++) {
        sum += a[i];
    }
    return sum;
}</pre>
```

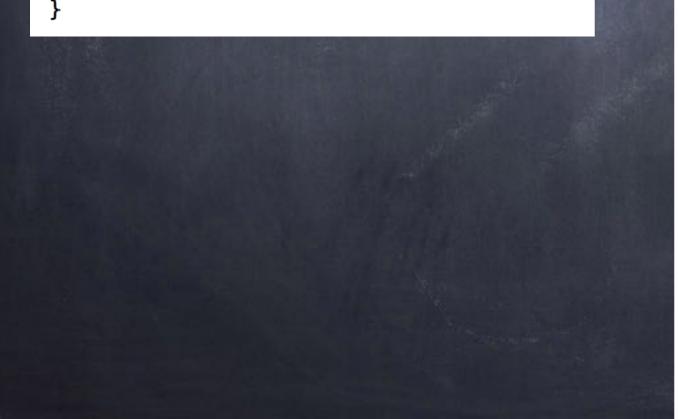




### This is on X5690 (Westmere)

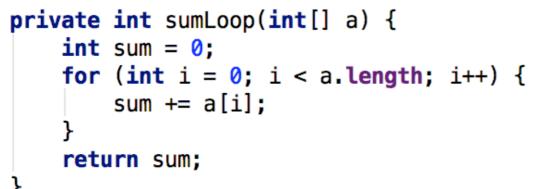
### Uses SSE (128bit)

```
private int sumLoop(int[] a) {
    int sum = 0;
    for (int i = 0; i < a.length; i++) {
        sum += a[i];
    }
    return sum;
}</pre>
```



		0x30001e81	andi \$7, %r80	
		0x3000fe85	movq %rcx, %rdx	
		0x3000fe88	subq %r8, %rdx	
		0x3000fe8b	je -35 ; ABS: 0x3000fe6a	
		0x3000fe8d	leaq 28(%rsi), %rax	
		0x3000fe91	pxor %xmm0, %xmm0	
		0x3000fe95	movq %rdx, %rdi	
		0x3000fe98	pxor %xmm1, %xmm1	
		0x3000fe9c	nopl (%rax)	N
20.64%	327	0x3000fea0	movdqu -16(%rax), %xmm2	
35.42%	561	0x3000fea5	movdqu (%rax), %xmm3	
11.30%	179	0x3000fe	paddd %xmm2, %xmm0	
13.70%	217	0x3000fead	paddd %xmm3, %xmm1	
18.69%	296	0x3000feb1	addq \$32, %rax	
		0x3000feb5	addq \$-8, %rdi	
		0x3000feb9	jne -27 ; ABS: 0x3000fea0	
		0x3000febb	paddd %xmm0, %xmm1	
		0x3000febf	pshufd \$78, %xmm1, %xmm0	
		0x3000fec4	paddd %xmm1, %xmm0	
		0x3000fec8	phaddd %xmm0, %xmm0	
0.13%	2	0x3000fecd	movd %xmm0, %eax	
		0x3000fed1	testl %r8d, %r8d	
		0x3000fed4	jne -102 ; ABS: 0x3000fe70	
		0x3000fed6	jmp -92 ; ABS: 0x3000fe7c	
		0x3000fed8	xorl %eax, %eax	
		0x3000feda	popq %rcx	
		0x3000fedb	retq	
		0x3000fedc	movq %rsi, (%rsp)	
		0x3000fee0	movabsq \$805334400, %rax	
		0x3000feea	callq *%rax	
		0x3000feec	movq (%rsp), %rsi	





### This is on E5-2690 v4 (Broadwell)

Uses AVX2 (256bit)

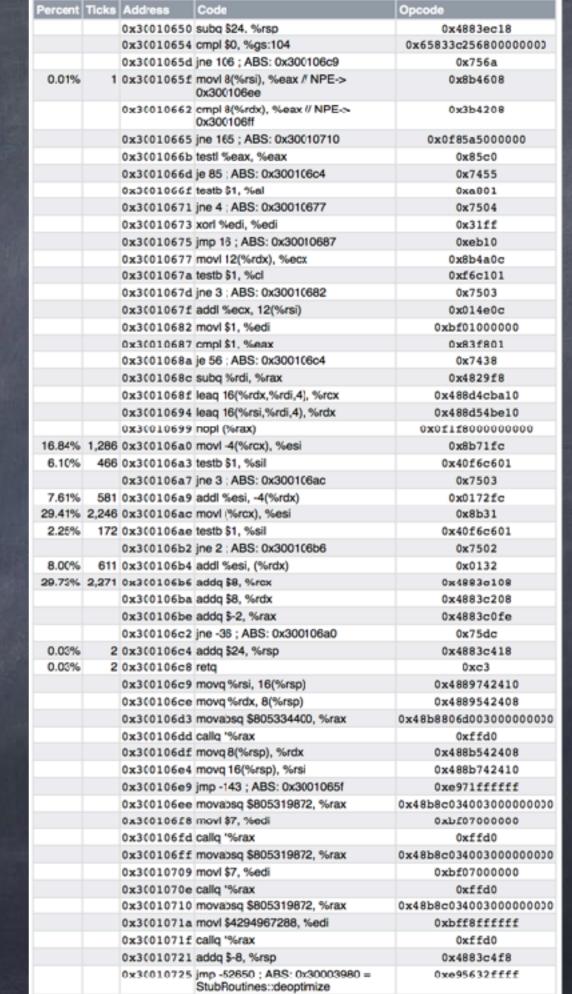
		0X3001269C	subd wird, wirdx	
		0x3001269f	je -47 ; ABS: 0x30012672	
		0x300126a1	leaq 108(%rsi), %rax	
		0x300126a5	vpxor %ymm0, %ymm0, %ymm0	
		0x300126a9	movq %rdx, %rdi	
		0x300126ac	vpxor %ymm1, %ymm1, %ymm1	
		0x300126b0	vpxor %ymm2, %ymm2, %ymm2	
		0x300126b4	vpxor %ymm3, %ymm3, %ymm3	
		0x30012658	nopl (%rax %rax)	0 x
0.06%	1	0x300126c0	vpaddd -96(%rax), %ymm0, %ymm0	
37.08%	603	0x300126c5	vpaddd -64(%rax), %ymm1, %ymm1	1
0.92%	15	0x300126ca	vpaddd -32(%rax), %ymm2, %ymm2	
54.80%	891	0x300126cf	vpaddd (%rax), %ymm3, %ymm3	
0.55%	9	0x300126d3	subq \$-128, %rax	
5. <b>84%</b>	95	0x300126d7	addq \$-32, %rdi	
		0x300126db	jne -29 ; ABS: 0x300126c0	
		0x300126dd	vp_dd %ymm0, %ymm1, %ymm0	
0.12%	2	0x300126e1	ypaddd %ymm2, %ymm0, %ymm0	
		0x300126e5	vpaddd %ymm3, %ymm0, %ymm0	
		0x300126e9	vextracti128 \$1, %ymm0, %xmm1	
0.25%	4	0x309126ef	vpaddd %ymm0, %ymm1, %ymm0	
		0x300126f3	vpshufd \$78, %xmm0, %xmm1	
0. <b>06%</b>	1	0x300126f8	vpaddd %ymm0, %ymm1, %ymm0	
0.12%	2	0x300126fc	vphaddd %ymm0, %ymm0, %ymm0	
		0x30012701	vmovd %xmm0, %eax	
		0x30012705	testl %r8d, %r8d	
		0x30012708	jne -142 ; ABS: 0x30012680	
		0x3001270e	jmp -134 ; ABS: 0x3001268d	
		0x30012713	xorl %eax, %eax	
		0x30012715	popq %rcx	
		0x30012716	retq	
		0x30012717	movq %rsi, (%rsp)	
		0x3001271b	movabsq \$805344640, %rax	<b>0x4</b> 81
		0x30012725	callq *%rax	
		0x30012727	movq (%rsp), %rsi	
		0x3001272b	jmp -207 ; ABS: 0x30012661	

#### A conditional array cell addition loop

```
private void addArraysIfEven(int a[], int b[]) {
    if (a.length != b.length) {
        throw new RuntimeException("length mismatch");
    }
    for (int i = 0; i < a.length; i++) {
        if ((b[i] & 0x1) == 0) {
            a[i] += b[i];
        }
    }
}</pre>
```



```
private void addArraysIfEven(int a[], int b[]) {
    if (a.length != b.length) {
        throw new RuntimeException("length mismatch");
    }
    for (int i = 0; i < a.length; i++) {
        if ((b[i] & 0x1) == 0) {
            a[i] += b[i];
        }
    }
}</pre>
```



0x3001072a int3

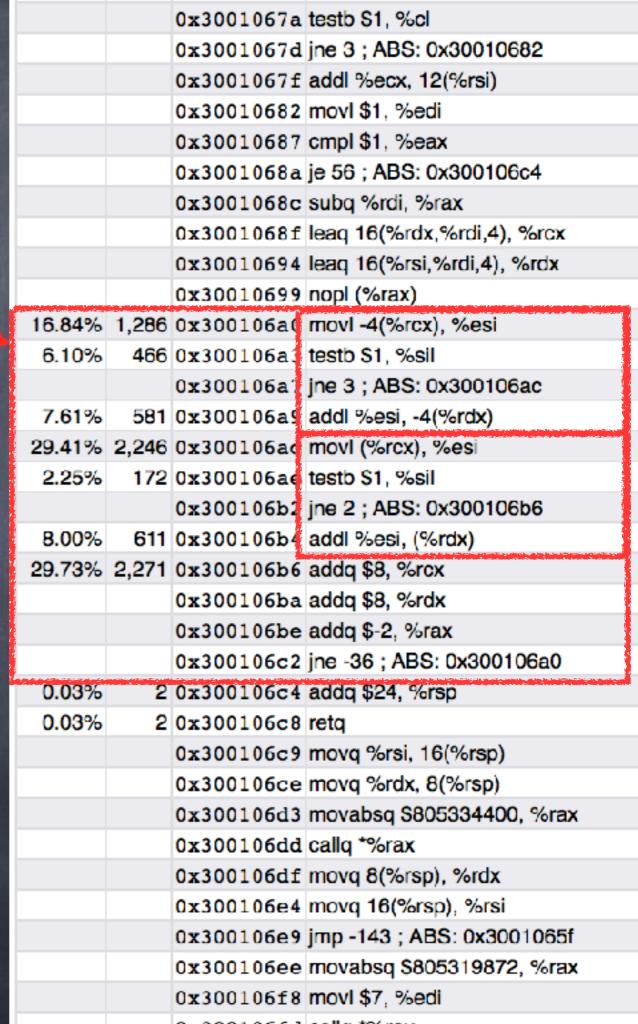


0xcc

#### Traditional JVM JITs

# per-element jumps,2 elements per iteration

```
private void addArraysIfEven(int a[], int b[]) {
    if (a.length != b.length) {
        throw new RuntimeException("length mismatch");
    }
    for (int i = 0; i < a.length; i++) {
        if ((b[i] & 0x1) == 0) {
            a[i] += b[i];
        }
    }
}</pre>
```





```
private void addArraysIfEven(int a[], int b[]) {
    if (a.length != b.length) {
        throw new RuntimeException("length mismatch");
    }
    for (int i = 0; i < a.length; i++) {
        if ((b[i] & 0x1) == 0) {
            a[i] += b[i];
        }
    }
}</pre>
```

### This is on E5-2690 v4 (Broadwell)

Vectorized with AVX2 32 elements per iteration

			$0 \times 30014548$	movabsq \$805389904, %rbx
			0x30014552	vpbroadcastd (%rbx), %ymm0
			0x30014557	vpxor %ymm1, %ymm1, %ymm1
			0x3001455b	movq %rdi, %rbx
			0x3001455e	nop
	0.15%	4	0x30014560	vmovdqu -96(%r11), %ymm2
	12.31%	320	0x30014566	vmovdqu -64(%r11), %ymm3
	0.50%	13	0x3001456c	vmovdqu -32(%r11), %ymm4
	2.04%	53	0x30014572	vmovdqu (%r11), %ymm5
	0.31%	8	0x30014577	vpand %ymm0, %ymm2, %ymm6
	4.54%	118	0x3001457b	vpand %ymm0, %ymm3, %ymm7
	0.69%	18	0x3001457f	vpand %ymm0, %ymm4, %ymm8
	1.35%	35	0x30014583	vpand %vmm0, %vmm5, %vmm9
	0.42%	11	0x30014587	vpcmpeqd %ymm1, %ymm6, %ymm6
	2.58%	67	0x3001458b	vpmaskmovd -96(%rcx), %ymm6, %ymm10
	3.58%	93	0x30014591	vpcmpeqd %ymm1, %ymm7, %ymm7
	2.12%	55	0x30014595	vpmaskmovd -64(%rcx), %ymm7, %ymm11
	12.12%	315	0x3001459b	vpcmpeqd %ymm1, %ymm8, %ymm8
	1.50%	39	0x3001459f	vpmaskmovd -32(%rcx), %ymm8, %ymm12
	3.69%	96	0 <b>x300145a</b> 5	vpcmpeqd %ymm1, %ymm9, %ymm9
	1.81%	47	0x300145a9	vpmaskmovd (%rcx), %ymm9, %ymm13
	12.27%	319	0x300145ae	vpaddd %ymm2, %ymm10, %ymm2
	0.58%	15	0x300145b2	vpaddd %ymm3, %ymm11, %ymm3
	0.19%	5	0x300145b6	vpaddd %ymm4, %ymm12, %ymm4
Staff and	0.58%	15	0x300145ba	vpaddd %ymm5, %ymm13, %ymm5
	3.27%	85	0x300145be	vpmaskmovd %ymm2, %ymm6, -96(%rcx)
	7.15%	186	0x300145c4	vpmaskmovd %ymm3, %ymm7, -64(%rcx)
	13.65%	355	0x300145ca	vpmaskmovd %ymm4, %ymm8, -32(%rcx)
	4.58%	119	0x300145d	vpmaskmovd %ymm5, %ymm9, (%rcx)
	6.81%	177	0x300145d5	subq \$-128, %r11
	0.69%	18	0x300145d9	subq \$-128, %rcx
	0.31%	8	0x300145dd	addq \$-32, %rbx
			0x300145e1	jne -135 ; ABS: 0x30014560
			0x300145e7	testl %r9d, %r9d
			0x300145ea	jne -356 ; ABS: 0x3001448c
			0x300145f0	jmp -233 ; ABS: 0x3001450c



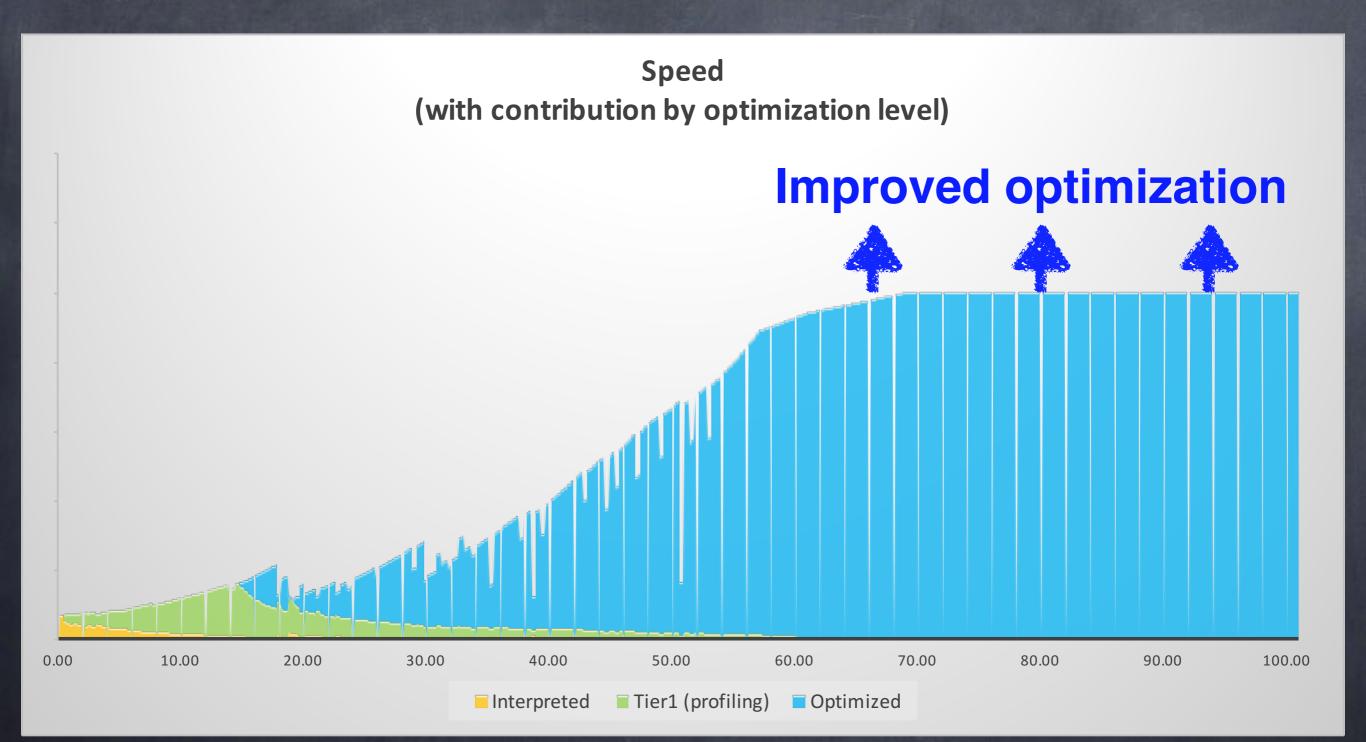
```
private void addArraysIfEven(int a[], int b[]) {
    if (a.length != b.length) {
        throw new RuntimeException("length mismatch");
    }
    for (int i = 0; i < a.length; i++) {
        if ((b[i] & 0x1) == 0) {
            a[i] += b[i];
        }
    }
}</pre>
```

#### This is on Skylake SP

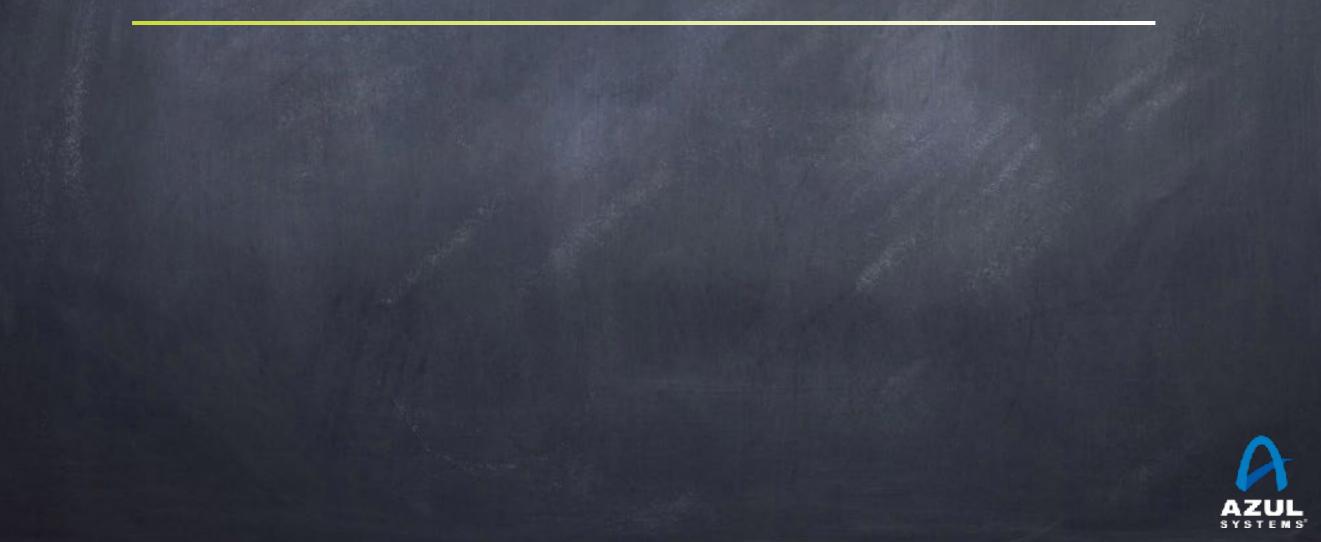
Vectorized with AVX512 64 elements per iteration

			0x3001ab4a	jb -206 ; ABS: 0x3001aa82
			0x3001ab50	movl %eax, %ecx
			0x3001ab52	andl \$2147483584, %ecx
	0.02%	1	0x3001ab58	leaq 204(%rdx), %r8
			0x3001ab5f	leaq 204(%rsi), %rdi
			0x3001ab66	movabsq \$805416064, %rbx
			0x3001ab70	vpbroadcastd (%rbx), %zmm0
	0.04%	2	0x3001ab76	movq %rcx, %r9
			0x3001ab79	nopl (%rax)
	0.46%	24	0x3001ab80	vmovdqu32 -192(%r8), %zmm1
	3.46%	179	0x3001ab87	vmovdqu32 -128(%r8), %zmm2
	3.52%	182	0x3001ab8e	vmovdqu32 -64(%r8), %zmm3
	4.95%	256	0x3001ab95	vmovdqu32 (%r8), %zmm4
	12.73%	658	0x3001ab9b	vptestnmd %zmm0, %zmm1, %k1
	1.99%	103	0x3001aba1	vptestnmd %zmm0, %zmm2, %k2
	1.53%	79	0x3001aba7	vptestnmd %zmm0, %zmm3, %k3
	2.69%	139	0x3001abad	vmovdqu32 -192(%rdi), %zmm5 {%k1} {z}
	8.38%	433	0x3001abb4	vmovdqu32 -128(%rdi), %zmm6 {%k2} {z}
	11.13%	575	0x3001abbb	vmovdqu32 -64(%rdi), %zmm7 {%k3} {z}
2	10.51%	543	0x3001abc2	vptestnmd %zmm0, %zmm4, %k4
	1.76%	91	0x3001abc8	vmovdqu32 (%rdi), %zmm8 {%k4} {z}
	18.25%	943	0x3001abce	vpaddd %zmm1, %zmm5, %zmm1
	0.97%	50	0x3001abd4	vpaddd %zmm2, %zmm6, %zmm2
	0.91%	47	0x3001abda	vpaddd %zmm3, %zmm7, %zmm3
Ĩ.	0.46%	24	0x3001abe0	vmovdqu32 %zmm1, -192(%rdi) {%k1}
	2.26%	117	0x3001a_7	vmovdqu32 %zmm2, -128(%rdi) {%k2}
	0.66%	34	0x3001abee	vmovdqu32 %zmm3, -64(%rdi) {%k3}
	1.10%	57	0x3001abf5	vpaddd %zmm4, %zmm8, %zmm1
	0.64%	33	0x3001abfb	vmovdqu32 %zmm1, (%rdi) {%k4}
	9.85%	514	0x3001ac01	addq \$256, %r8
	0.64%	33	0x3001ac08	addq \$256, %rdi
	0.85%	44	0x3001ac0f	addq \$-64, %r9
			0x3001ac13	jne -153 ; ABS: 0x3001ab80
			0x3001ac19	cmpq %rax, %rcx
			0x3001ac1c	jne -414 ; ABS: 0x3001aa84
			0x2001ac22	imp -262 · ABS· 0v3001ab21

### Better JIT'ing is basically about speed



### Compiler Stuff



### Some simple compiler tricks



### Code can be reordered...

int doMath(int x, int y, int z) {
 int a = x + y;
 int b = x - y;
 int c = z + x;
 return a + b;

#### Can be reordered to:

}

int doMath(int x, int y, int z) {
 int c = z + x;
 int b = x - y;
 int a = x + y;
 return a + b;



### Dead code can be removed

int doMath(int x, int y, int z) {
 int a = x + y;
 int b = x - y;
 int c = z + x;
 return a + b;
}

#### Can be reduced to:

}

int doMath(int x, int y, int z) {
 int a = x + y;
 int b = x - y;
 return a + b;



## Values can be propagated

int doMath(int x, int y, int z) {

int a = x + y; int b = x - y; int c = z + x; return a + b; }

Can be reduced to:

int doMath(int x, int y, int z) {
 return x + y + x - y;
}



## Math can be simplified

int doMath(int x, int y, int z) {
 int a = x + y;
 int b = x - y;
 int c = z + x;
 return a + b;
}

#### Can be reduced to:

int doMath(int x, int y, int z) {
 return x + x;
}



### Some more compiler tricks



```
propagation can affect flow
  constants can be propagated to pre-compute results:
          int computeBias() {
           int bias, val = 5;
           if (val > 10) {
               bias = computeComplicatedBias(val);
           else {
               bias = 1;
           return bias;
  Can be reduced to:
```

int computeBias() {
 return 1;



### Reads can be cached

class Point { int x, y; }
int distanceRatio(Point a) {
 int distanceTo = a.x - start;
 int distanceAfter = end - a.x;
 return distanceTo/distanceAfter;

Is (semantically) the same as
 int distanceRatio(Point a) {
 int x = a.x;
 int distanceTo = x - start;
 int distanceAfter = end - x;
 return distanceTo/distanceAfter;

}



### Reads can be cached

class Trigger { boolean flag; }
void loopUntilFlagSet(Tigger a) {
 while (!a.flag) {
 loopcount++;
 }
}

#### Is the same as:

}

}

void loopUntilFlagSet(Object a) {
 boolean flagIsSet = a.flag;
 while (!flagIsSet) {
 loopcount++;



That's what volatile is for...

### Writes can be eliminated

Intermediate values might never be visible

void updateDistance(Point a) {
 int distance = 100;

a.x = distance; a.x = distance \* 2; a.x = distance \* 3;

#### Is the same as

void updateDistance(Point a) {
 a.x = 300;



### Writes can be eliminated

Intermediate values might never be visible

void updateDistance(SomeObject a) {
 a.visibleValue = 0;

for (int i = 0; i < 1000000; i++) { a.internalValue = i;

a.visibleValue = a.internalValue;

Is the same as

void updateDistance(SomeObject a) {
 a.internalValue = 1000000;
 a.visibleValue = 1000000;



## Inlining...

public class Thing {
 private int x;
 public final int getX() { return x };
}

myX = thing.getX();

Is the same as

...

Class Thing { int x; }





### Inlining is very powerful inlining exposes other optimizations int computeBias(int val) { int bias; if (val > 10) { bias = computeComplicatedBias(val); else { bias = 1; } return bias; } ... myBias = computeBias(5);

Can be reduced to:

myBias = 1;

### A uBenchmark sidetrack



#### A simple loop uBenchmark (0)

Turns out this is "really fast" As in: when count = 1,000,000 we complete ~500,000,000 calls per second (for 5,000,000,000,000,000 iterations/sec)

#### A simple loop uBenchmark (1)

```
public void loopUbench1(int count) {
    long sum = 0;
    for (int i = 0; i < count; i++) {
        sum += i;
     }
}</pre>
```

Still "impossibly fast"

It's all "provably dead code". Compiler translates the method to a no-op



### A simple loop uBenchmark (2)

```
public long loopUbench2(int count) {
    long sum = 0;
    for (int i = 0; i < count; i++) {
        sum++;
     }
    return sum;
}</pre>
```

Better? No. Still "impossibly fast". Compiler returns count. <u>No loop.</u>



#### A simple loop uBenchmark (3)

```
public long loopUbench3(int count) {
    long sum = 0;
    for (int i = 0; i < count; i++) {
        sum += i;
    }
    return sum;
}</pre>
```

#### Better?

Depends. On HotSpot and Zing C2, yes. But Zing's new Falcon compiler is smart enough to recognize arithmetic series

### A simple loop uBenchmark (4)

```
public long loopUbench4(int count) {
    long sum = 0;
    for (int i = 0; i < count; i++) {
        sum *= i;
    }
    return sum;
}</pre>
```

How about this?

Zing's Falcon will even figure out this one. (it returns zero)



### A simple loop uBenchmark (5)

```
public long loopUbench5(int count) {
    long sum = 0;
    for (int i = 0; i < count; i++) {
        sum += i + ((i - 3) & 0x7);
     }
    return sum;
}</pre>
```

### Seems to be complicated enough to defeat \*current\* compilers...



## uBenchmarking Takeaways

- uBenchmarking is "hard". As in "very tricky"
- You may not be measuring what you think you are
- Trickiness" can change over time, between versions
- Sanity check EVERYTHING
- Ø Use jmh
- Ø Use jmh
- Ø Use jmh
- And even then, suspect everything



### Back to compiler stuff



Speculative compiler tricks

JIT compilers can do things that static compilers can have a hard time with...



### Untaken path example

"Never taken" paths can be optimized away with benefits:

```
int computeMagnitude(int val) {
    if (val > 10) {
        bias = computeBias(val);
    else {
        bias = 1;
    }
    return Math.log10(bias + 99);
```

When all values so far were <= 10, could be compiled to: int computeMagnitude(int val) { if (val > 10) uncommonTrap(); return 2;

## Implicit Null Check example

All field and array access in Java is null checked

x = foo.x;

is (in equivalent required machine code):

if (foo == null)
 throw new NullPointerException();
x = foo.x;

But compiler can "hope" for non-nulls, and handle SEGV <Point where later SEGV will appear to throw> x = foo.x;

This is faster \*IF\* no nulls are encountered...



## Class Hierarchy Analysis (CHA)

Can perform global analysis on currently loaded code Deduce stuff about inheritance, method overrides, etc.
 Can make optimization decisions based on assumptions
 Re-evaluate assumptions when loading new classes Throw away code that conflicts with assumptions before class loading makes them invalid



# Inlining works without "final"

public class Animal {
 private int color;
 public int getColor() { return color };
}

myColor = animal.getColor();

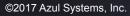
#### Is the same as

...

Class Animal { int color; } \*THIS\* (CHA) is why Java field accessors are free & clean

myColor = animal.color;

As long as only one implementer of getColor() exists



# Inlining monomorphic sites

public class Animal {
 private int color;
 public int getColor() { return color };
}

myColor = animal.getColor();

Can be converted to:

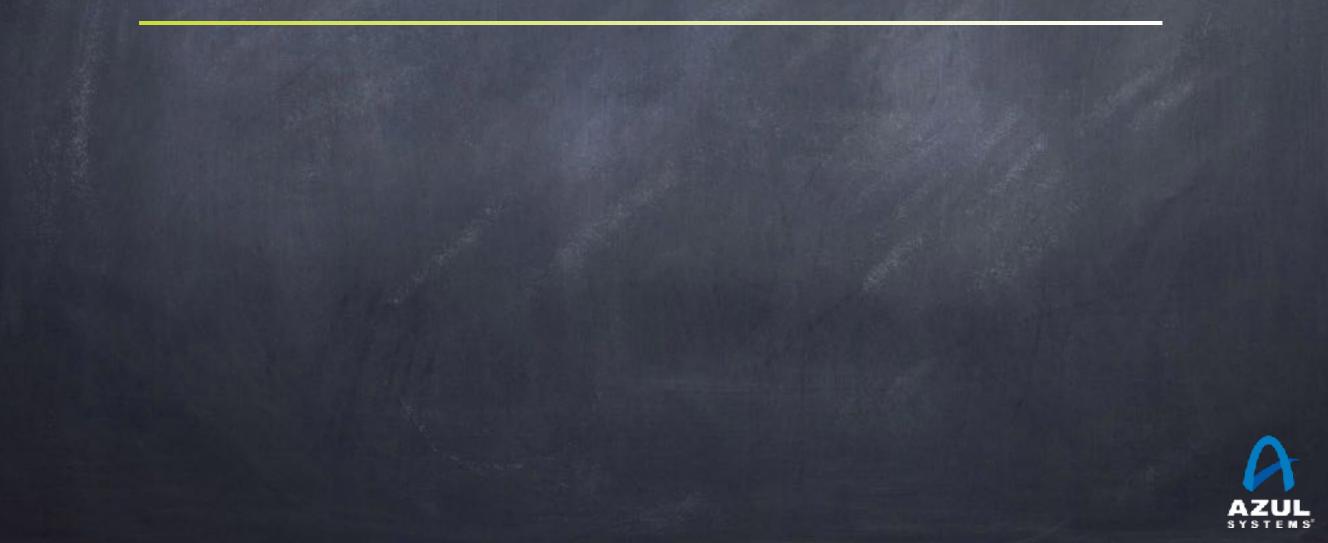
...

if (animal.type != Dog) uncommonTrap(); myColor = animal.color;

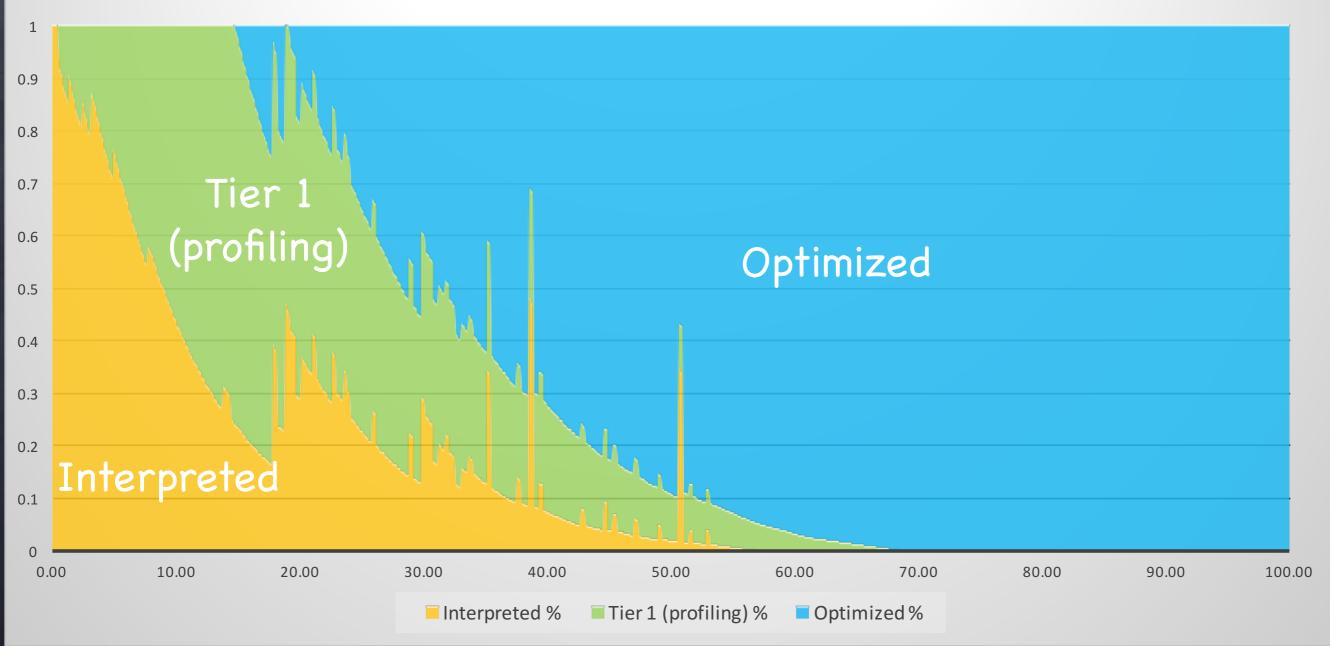
#### Even if we have multiple conflicting implementors...



## Deoptimization







# Deoptimization: Adaptive compilation is... adaptive Micro-benchmarking is a black art So is the art of the Warmup Running code long enough to compile is just the start... Deoptimizations can occur at any time often occur after you \*think\* the code is warmed up. Many potential causes



## Warmup often doesn't cut it...

#### Common Example:

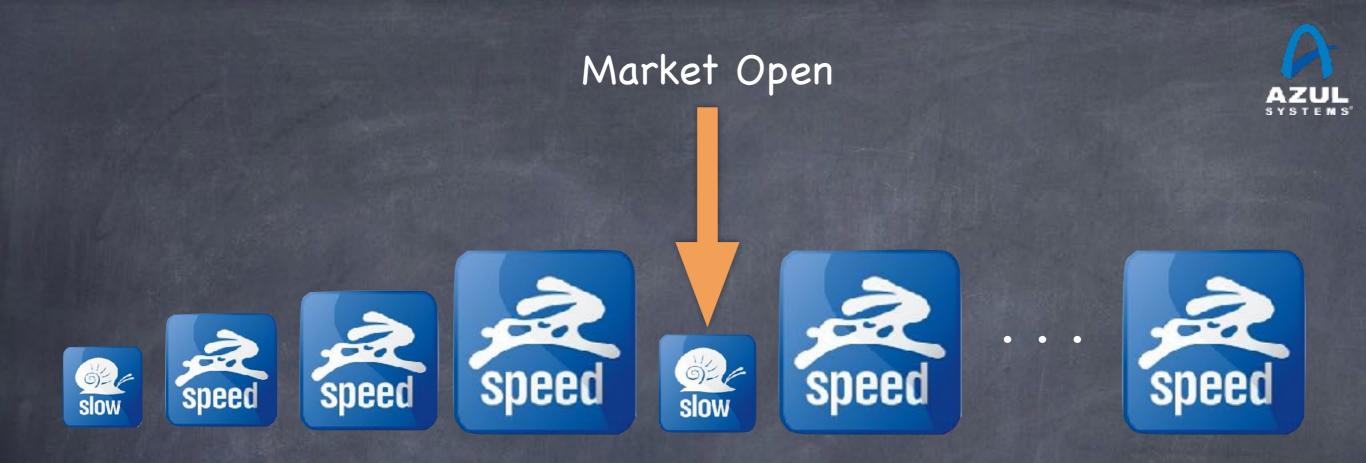
Trading system wants to have the first trade be fast
So run 20,000 "fake" messages through the system to warm up
let JIT compilers optimize, learn, and deopt before actual trades

@ But...

Code is written to do different things "if this is a fake message"
e.g. "Don't send to the exchange if this is a fake message"

#### What really happens

JITs optimize for fake path, including speculatively assuming "fake"
 First real message through causes a deopt...



### Java at Market Open



# Java's "Just In Time" Reality













Warmup

Starts slow, learns fast

Lazy loading & initialization

 Aggressively optimized for the common case

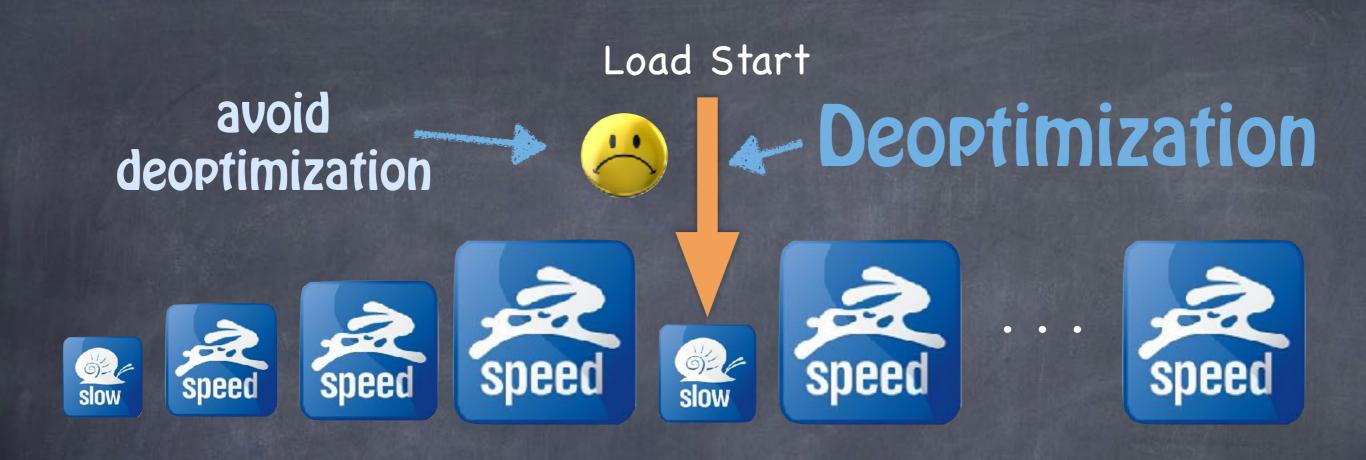
(temporarily) Reverts to
 slower execution to adapt



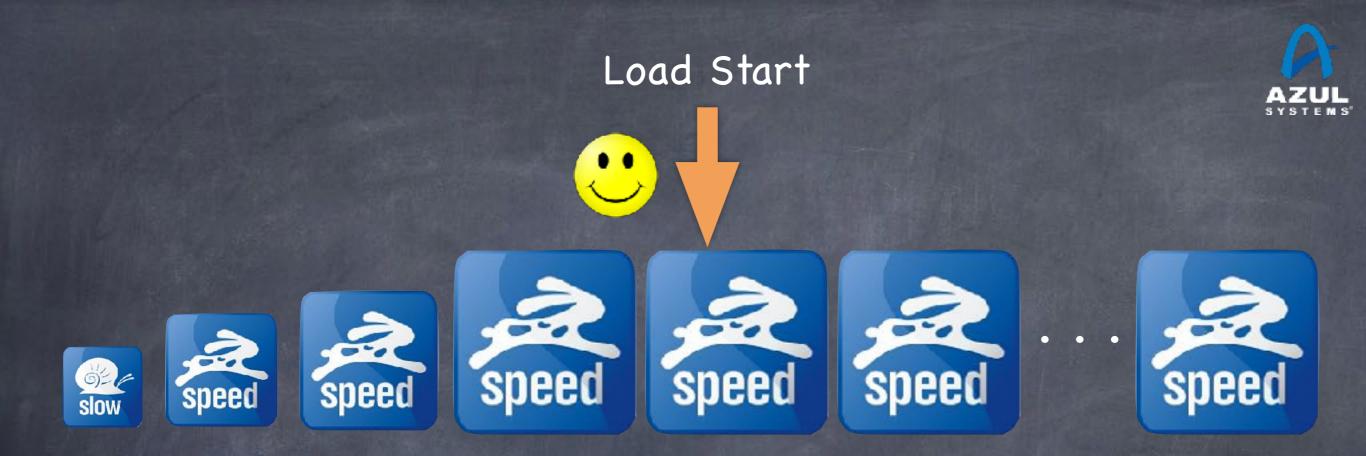
## Logging and "replaying" optimizations

#### Log optimization information

Record ongoing optimization decisions and stats Record optimization dependencies Stablish "stable optimization state" at end of previous run Read prior logs at startup Prime JVM with knowledge of prior stable optimizations Apply optimizations as their dependencies get resolved Build workflow to promote confidence Let you know if/when all optimizations have been applied If some optimization haven't been applied, let you know why...

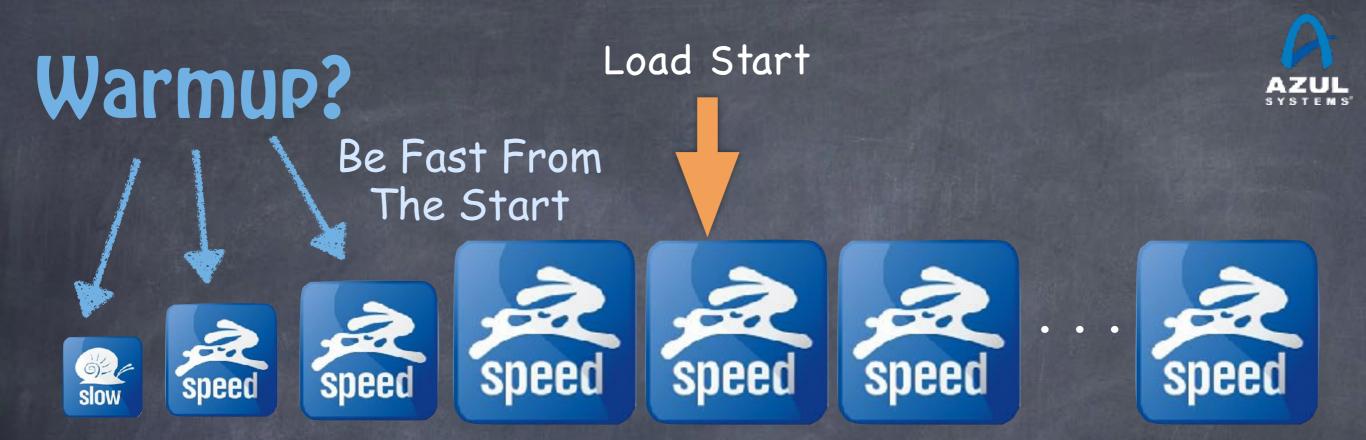




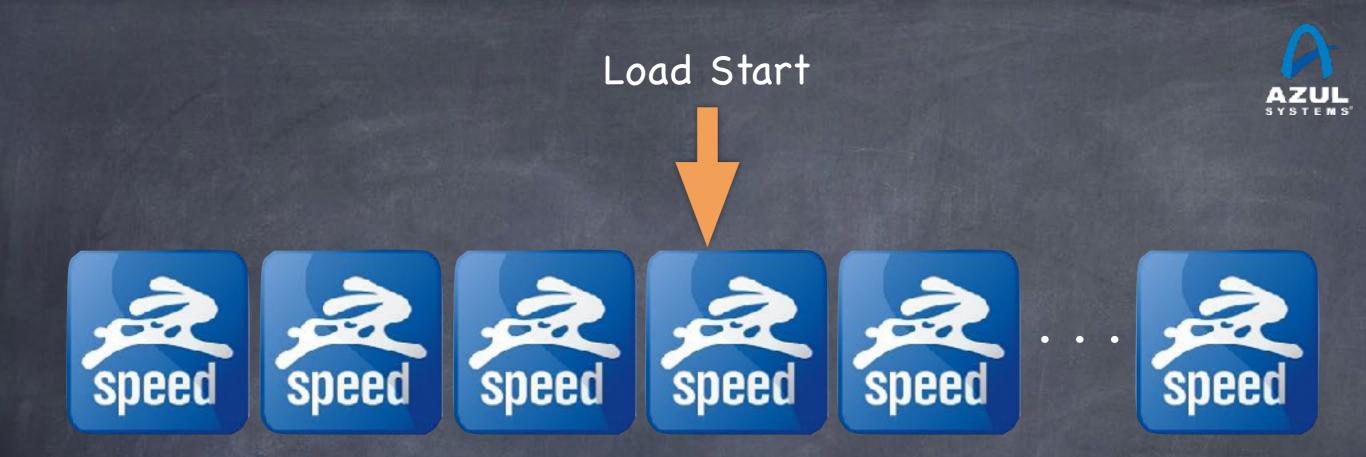


## With de-optimization avoided

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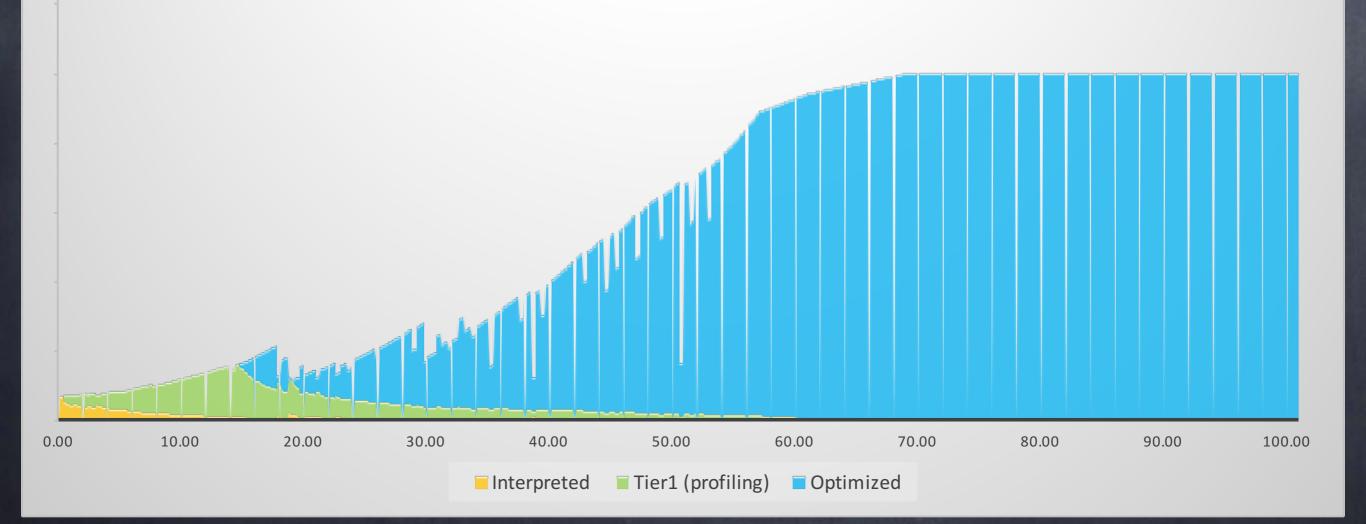
# With pre-loading of prior optimizations

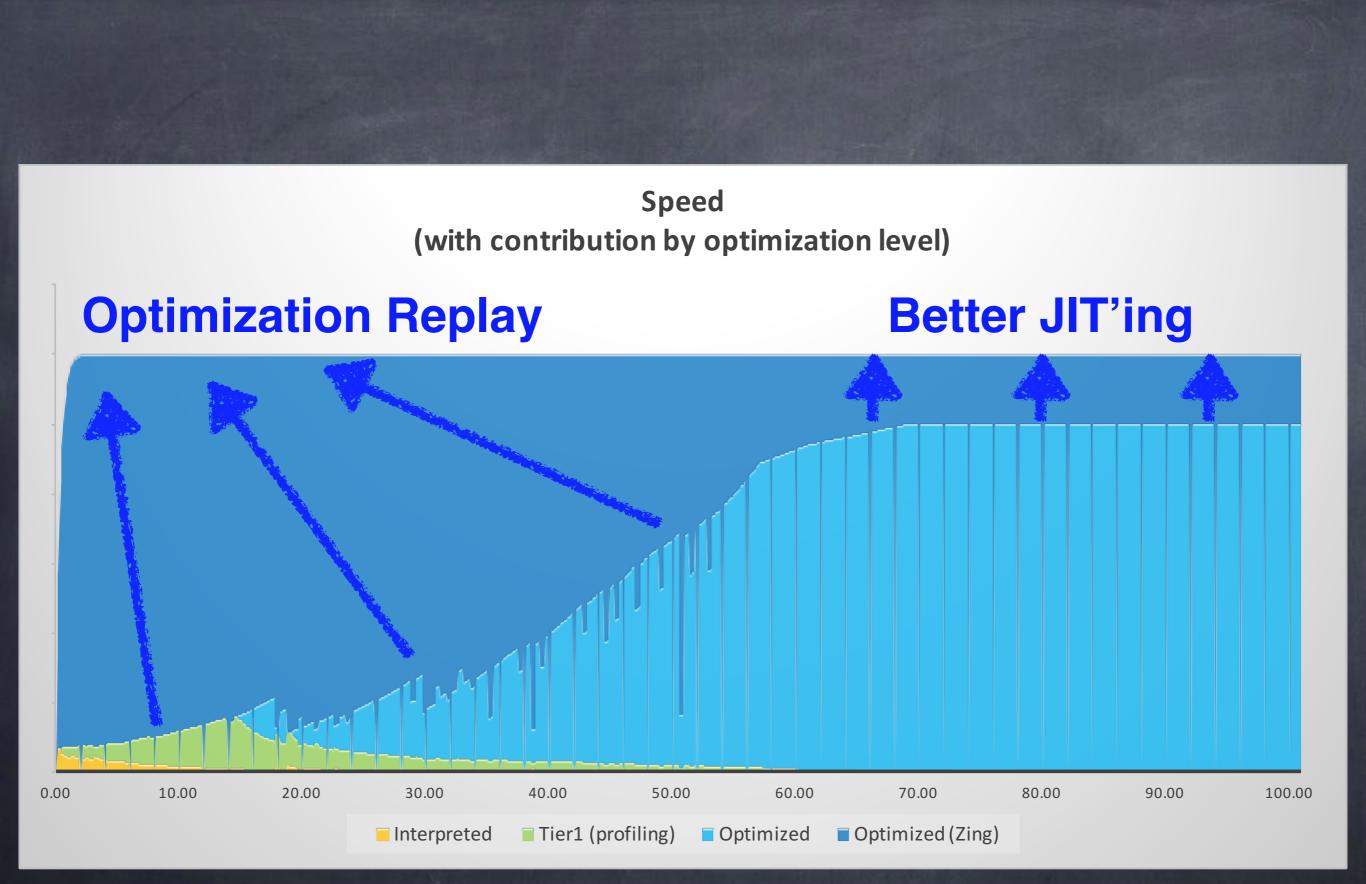


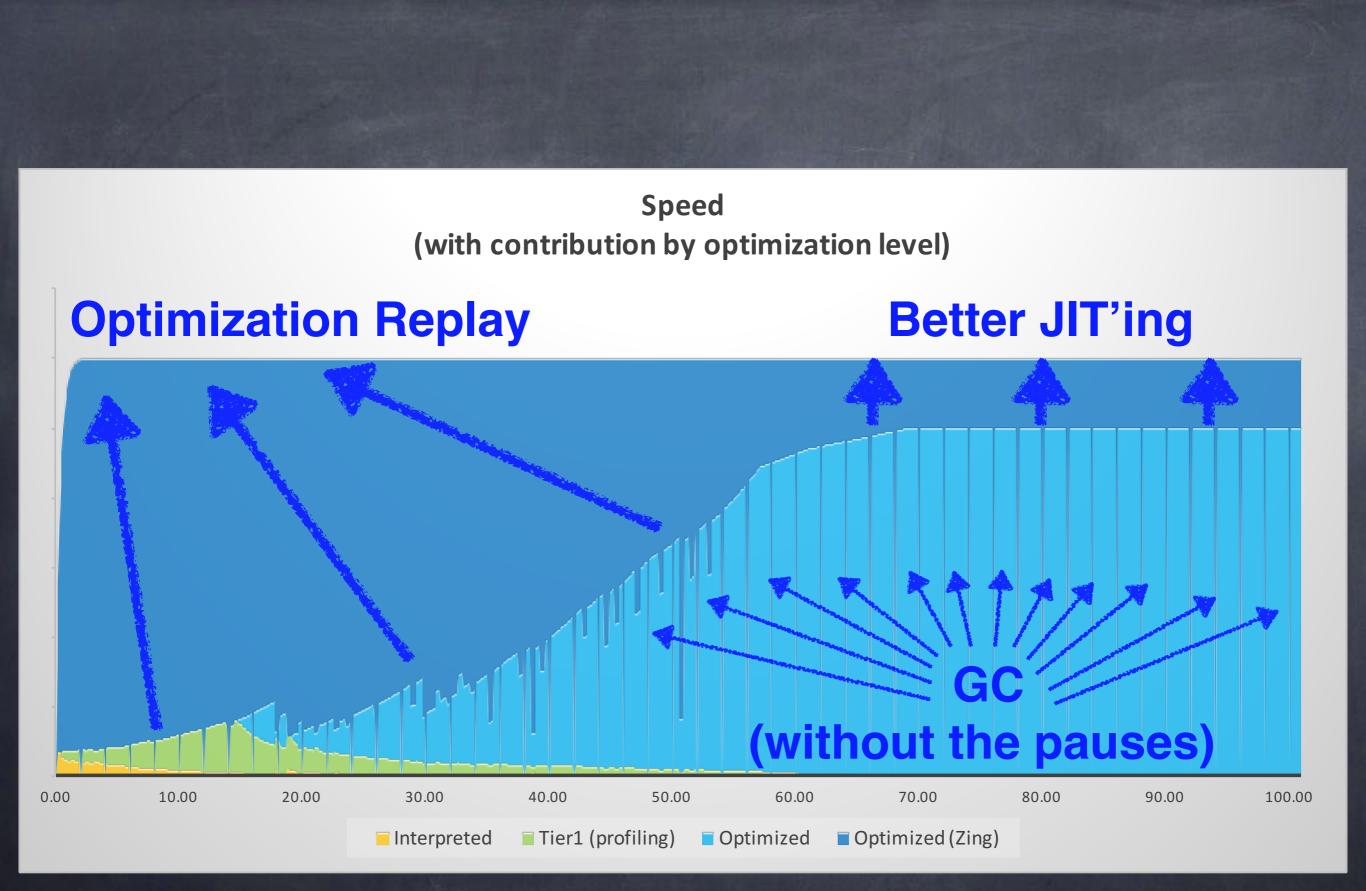


## Speed improvements

Speed (with contribution by optimization level)







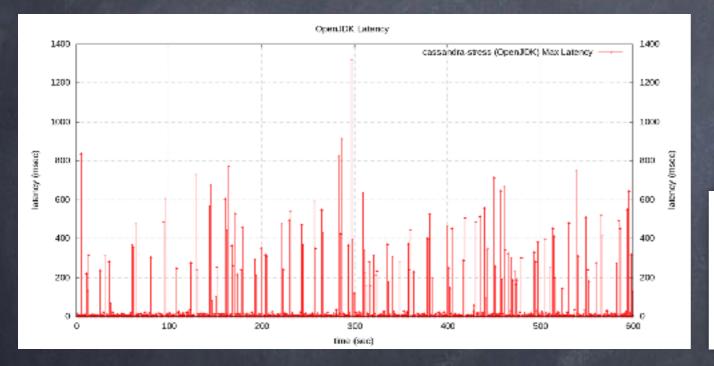
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### C4 Garbage Collector

# **ELIMINATES** Garbage Collection as a concern for enterprise applications



#### This is <Your App> on HotSpot



#### This is <Your App> on Zing

## Any Questions?



## GC Tuning

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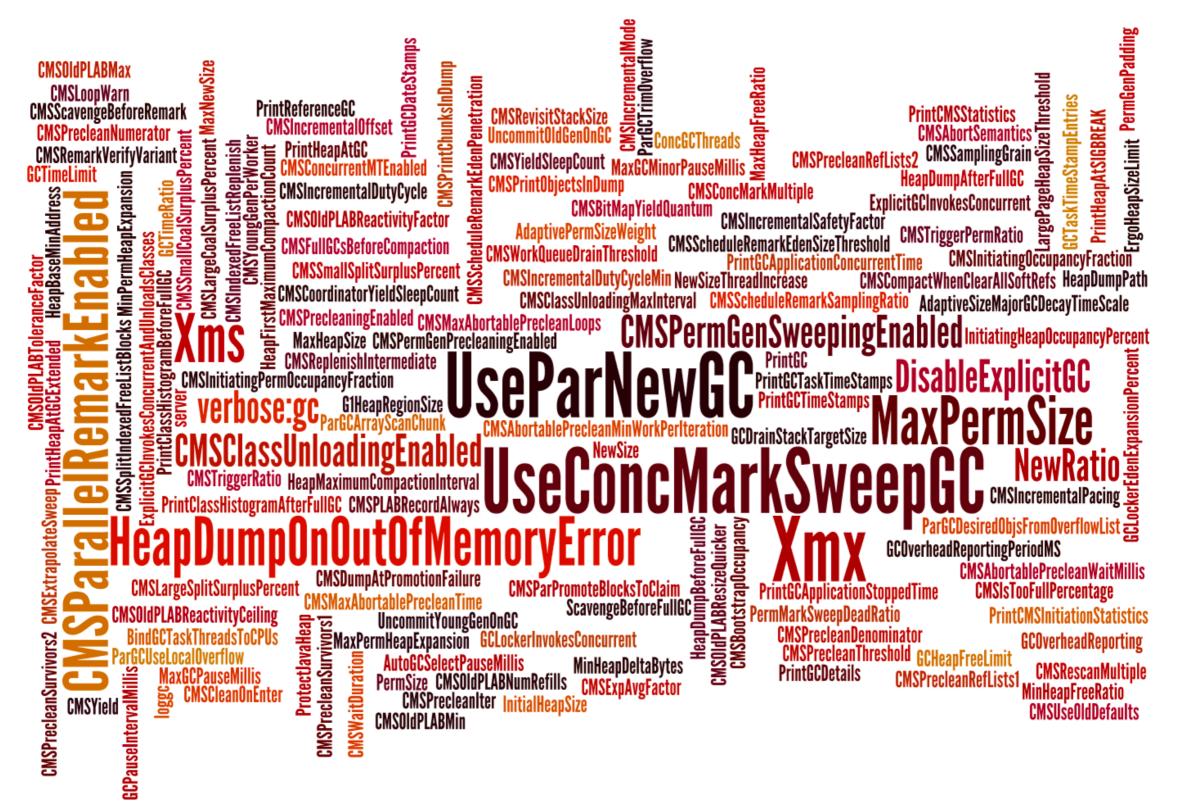
# Java GC tuning is "hard"...

Examples of actual command line GC tuning parameters:

Java -Xmx12g -XX:MaxPermSize=64M -XX:PermSize=32M -XX:MaxNewSize=2g -XX:NewSize=1, -XX:SurvivorRatio=128 -XX:+UseParNewGC -XX:+UseConcMarkSweepoc -XX:maxTenuringThreshold=0 -XX:CMSInitiatingOccupancyFraction=60 -XX:+CMSParallelRemarkEnabled -XX:+UseCMSInitiatingOccupancyOnly -XX:ParallelGCThreads=12 -XX:LargePageSizeInBytes=256m ...

Java -Xms8g -Xmx8g -Xmn2g -XX:PermSize=64M -XX:MaxPermSize=256M -XX:-OmitStackTraceInFastThrov -XX:SurvivorRatio=2 XX:-UseAdaptiveSizePolicy -XX:+UseConcMarkSweepGC -XX:+CMSConcurrentMTEnabled -XX:+CMSParallelRemarkEnabled -XX:+CMSParallelSurvivorRemarkEnabled -XX:CMSMaxAbortablePrecleanTime=10000 -XX:+UseCMSInitiatingOccupancyOnly -XX:CMSInitiatingOccupancyFraction=63 -XX:+UseParNewGC -Xnoclassgc ...

## A few more GC tuning flags



# The complete guide to modern GC tuning

## java -Xmx40g

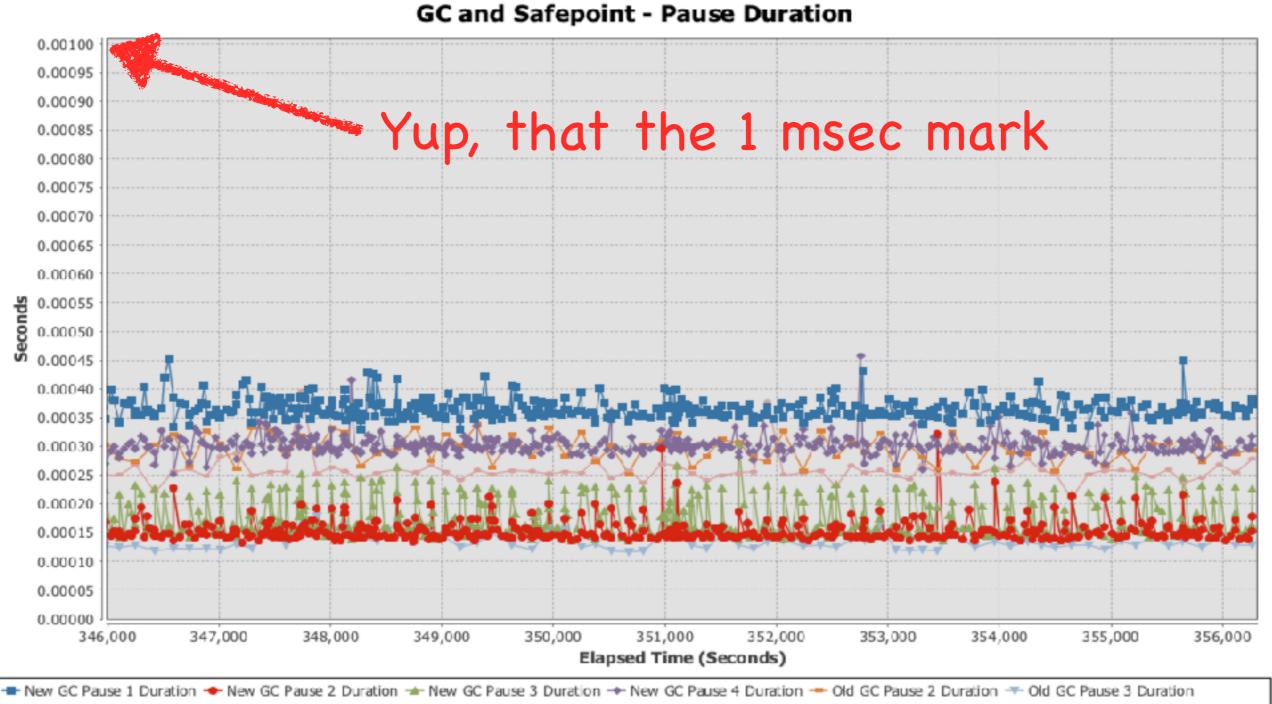
java -Xmx20g

java -Xmx10g

java -Xmx5g



#### Cassandra under heavy load, Intel E5-2690 v4 server



Old GC Pause 4 Duration 🔶 Deoptimize Pause 🕂 Force Safepoint Pause 🔫 Concurrent Deflation Pause 🚢 Thread Dump Pause 📌 Find Deadlocks Pause 📥 Other Safepoint Pause



### Warning: results may be too good



## A practical real-world example: Improve Digital (Video Advertising)

Cassandra cluster running on 6x AWS i3.2xlarge

- Approx. 80/20 write/read split
- Data read and written with quorum consistency
- 6 client machines sending requests collocated in the same AZs
- SLA requirements for read operations:
  - 20ms at 99.9%
  - 50ms at 99.99%
  - IOOms at 99.998% (not a typo, last 9 hard to maintain on AWS)

HotSpot w/G1: can maintain ~4K TPS before SLA breach

Zing: can maintain ~21K TPS before SLA breach





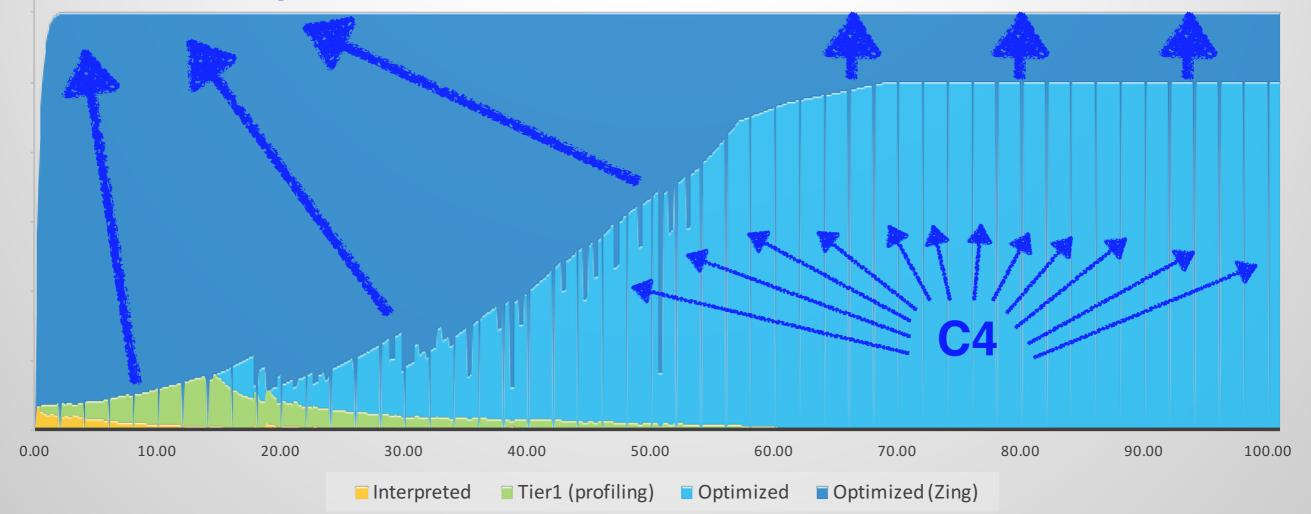
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## Start Fast, Go Fast, Stay Fast



#### **ReadyNow**

**Falcon** 



# 

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http://latencytipoftheday.blogspot.com

